

Intelligent Autonomous Agents and Cognitive Robotics

Rainer Marrone, Ralf Möller,
Hamburg University of Technology

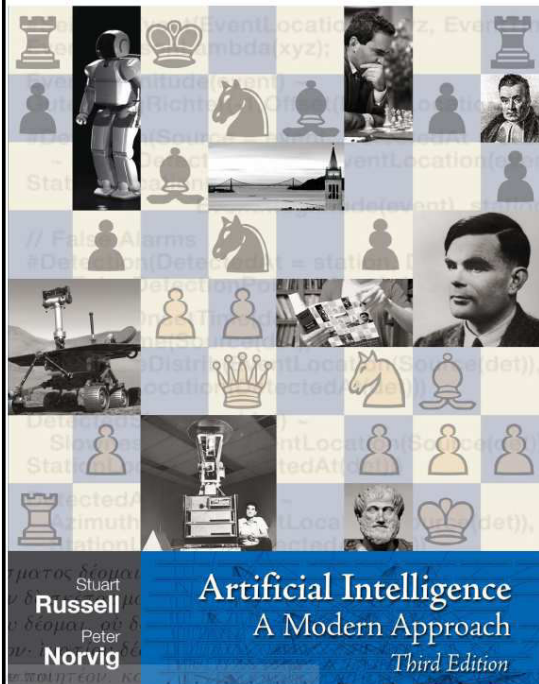
Lecture

- Lecture: Tuesday, 8:00 – 9:30, D-1.025
- pdf files of the lecture will be available on StudIP.

Exercise

- Thursday, 15:00-16:30, H-0.08
First exercise: 27.10
- I will upload exercise sheets every week, after the lecture.
- After the exercise, I will upload the solution as pdf.

Literature

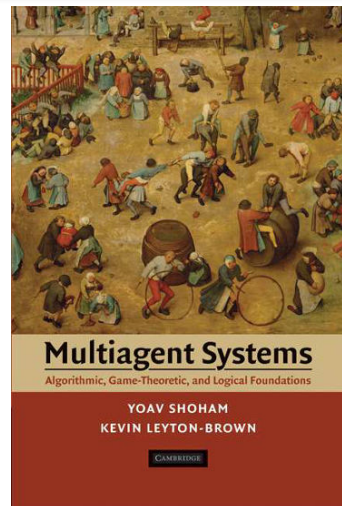


Chapters 2-5, 13 - 17

<http://aima.cs.berkeley.edu>

- with code repository
- further readings

Literature



Finally, we ask you not to link directly to the PDF or to distribute it electronically. Instead, we invite you to link to <http://www.masfoundations.org>. This will allow us to gauge the level of interest in the book and to update the PDF to keep it consistent with reprintings of the book.

Main Topics

- Solving Problems by Searching
- Adversarial Agents
- Constraint Satisfaction Problems
- Bayesian Networks
- Probabilistic Reasoning Over Time
- Decision Making
- Game Theory
- Mechanism Design

What is an Agent? (Wooldridge)

- Trivial (non-interesting) agents:
 - ♦ thermostat
 - ♦ UNIX daemon (e.g., xbiff)
- An intelligent agent is capable of *flexible autonomous action in some environment*
- By *flexible*, we mean:
 - ♦ *reactive*
 - ♦ *pro-active*
 - ♦ *social*

Reactivity

- A *reactive* system is one that maintains an ongoing interaction with its environment, and responds to changes that occur in it (in time for the response to be useful)
- The real world is more complicated: things change, information is incomplete. Many (most?) interesting environments are *dynamic*

Proactiveness

- Reacting to an environment is easy (e.g., stimulus → response rules)
- But we generally want agents to *do things for us*
- Hence *goal directed behavior*
- Pro-activeness = generating and attempting to achieve goals
 - ♦ Not driven solely by events
 - ♦ Taking the initiative

Balancing Reactive and Goal-Oriented Behavior

- We want our agents to be reactive, responding to changing conditions in an appropriate (timely) fashion
- We want our agents to systematically work towards long-term goals
- These two considerations can be at odds with one another
- Designing an agent that can balance the two remains an open research problem

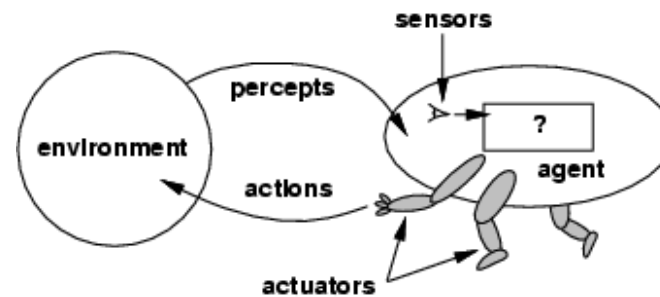
Social Ability

- The real world is a *multi*-agent environment: we cannot go around attempting to achieve goals without taking others into account
- Some goals can only be achieved with the cooperation of others
- *Social ability* in agents is the ability to interact with other agents (and possibly humans) via some kind of *agent-communication language*. Goal is to fulfill the design objectives commitments/cooperation.

Agents (Norvig, Russell)

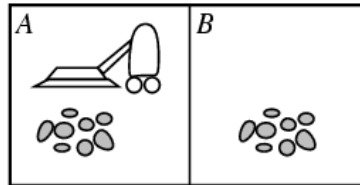
- An **agent** is anything that can be viewed as **perceiving** its **environment** through **sensors** and **acting** upon that environment through **actuators**
- *Human agent*: eyes, ears, and other organs for sensors; hands, legs, mouth, and other body parts for actuators
- *Robotic agent*: cameras and infrared range finders for sensors; various motors for actuators

Agents and environments



- The **agent function** maps from percept histories to actions .
 $f : P^* \rightarrow A$
- The **agent program** runs on the physical **architecture** to produce f
- agent = architecture + program

Vacuum-Cleaner World



- Percepts: location and contents, e.g., [A, Dirty]
- Actions: *Left*, *Right*, *Suck*, *NoOp*

A Vacuum-Cleaner Agent

Percept sequence	Action
<i>[A, Clean]</i>	<i>Right</i>
<i>[A, Dirty]</i>	<i>Suck</i>
<i>[B, Clean]</i>	<i>Left</i>
<i>[B, Dirty]</i>	<i>Suck</i>
<i>[A, Clean], [A, Clean]</i>	<i>Right</i>
<i>[A, Clean], [A, Dirty]</i>	<i>Suck</i>
<i>⋮</i>	<i>⋮</i>

Performance measure

- An agent should strive to "do the right thing", based on what it can perceive and the actions it can perform.
- Success to be measured w.r.t. an agent-local perspective of *environment states*.
- **Performance measure**: An objective criterion for success of an agent's behavior.
 - ♦ Performance measure of a vacuum-cleaner agent could be amount of dirt cleaned up, amount of time taken, amount of electricity consumed, amount of noise generated, etc.

Rational Agents

- **Rational Agent:** For each possible percept sequence, a rational agent
 - ◆ should select an action that is expected to maximize its *performance measure*,
 - ◆ given the evidence provided by the percept sequence and whatever built-in knowledge the agent has.
- **Rational = Intelligent**
- Rationality is distinct from omniscience (all-knowing with infinite knowledge)

Autonomous Agents

- Agents can perform actions in order to obtain useful information (information gathering, exploration)
- An agent is **autonomous** if its behavior is determined by its own experience (with ability to **learn** and **adapt**)

Applications

- Robotics: Drone, Explorer, Rescue BOT
- Web Agents: Personalized Search Engines
- Logistics: Tour planning
- Medicine: Diagnosis, Surgery, ...
- ...

First task in agent design: PEAS

Must first **specify** the setting/task environment for intelligent **agent design**.

- **P**erformance measure
- **E**nvironment
- **A**ctuators
- **S**ensors

PEAS

- Consider, e.g., the task of designing an automated taxi driver:
 - ◆ *Performance measure:*
 - Safe, fast, legal, comfortable trip, maximize profits, ...
 - ◆ *Environment:*
 - Roads, other traffic, pedestrians, customers, ...
 - ◆ *Actuators:*
 - Steering wheel, accelerator, brake, signal horn, ...
 - ◆ *Sensors:*
 - Cameras, sonar, speedometer, GPS, odometer, engine sensors, ...

PEAS

- Agent: Part-picking and sorting robot
 - ◆ *Performance measure:*
 - Percentage of parts in correct bins
 - ◆ *Environment:*
 - Conveyor belt with parts, bins
 - ◆ *Actuators:*
 - Jointed arm and hand, ...
 - ◆ *Sensors:*
 - Camera, joint angle sensors. ...

22

Environment Types

- **Fully observable** vs. **partially observable**: An agent's sensors give it access to the state of the environment at each point in time.
- **Deterministic** vs. **stochastic**: The next state of the environment is completely determined by the current state and the action executed by the agent. If the environment is deterministic except for the actions of other agents, then the environment is **strategic**.
- **Episodic** vs. **sequential**: The agent's experience is divided into atomic "episodes" (each episode consists of the agent perceiving and then performing a single action), and the choice of an action in each episode depends only on the episode itself.

Environment Types

- **Static** vs. **dynamic**: The environment is unchanged while an agent is deliberating. (The environment is **semidynamic** if the environment itself does not change with the passage of time but the agent's performance score does)
- **Discrete** vs. **continuous**: Discrete if there are a limited number of distinct, clearly defined percepts, states and actions.
- **Single agent** vs. **multiagent**: An agent operating by itself in an environment.

24

Environment Types

	Chess with a clock	Chess without a clock	Taxi driving
Fully observable	[]	
Deterministic			
Episodic			
Static			
Discrete			
Single agent			

- The environment type largely determines the agent design
- The real world is (of course) partially observable, stochastic, sequential, dynamic, continuous, multi-agent

Environment Types

	Chess with a clock	Chess without a clock	Taxi driving
Fully observable	Yes	Yes	No
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Environment Types

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Episodic	---	---	---
Static			
Discrete			
Single agent			

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Environment Types

	Chess with a clock	Chess without a clock	Taxi driving
Fully observable	Yes	Yes	No
Deterministic	Strategic	Strategic	No
Episodic	No	No	No
Static			
Discrete			
Single agent			

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Environment Types

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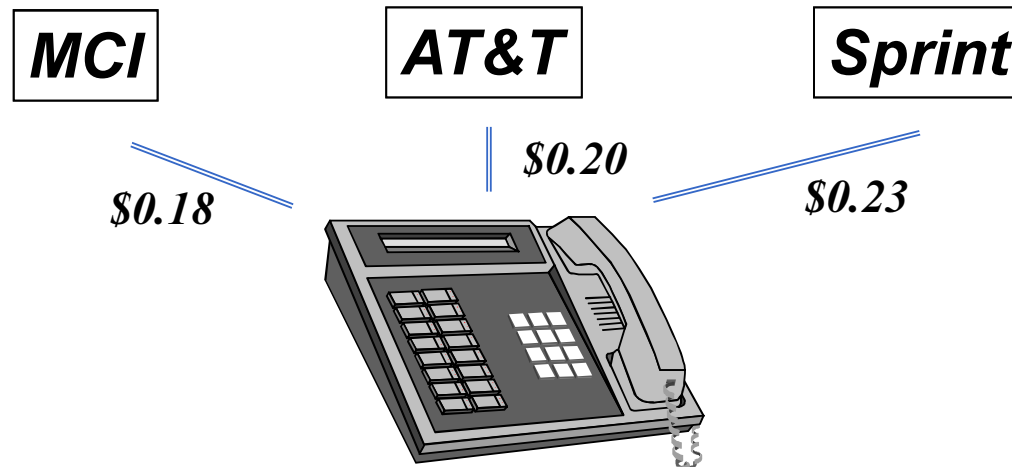
Environment Types

	Chess with a clock	Chess without a clock	Taxi driving
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Single agent	No	No	No

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Mechanisms for multi-agent environments

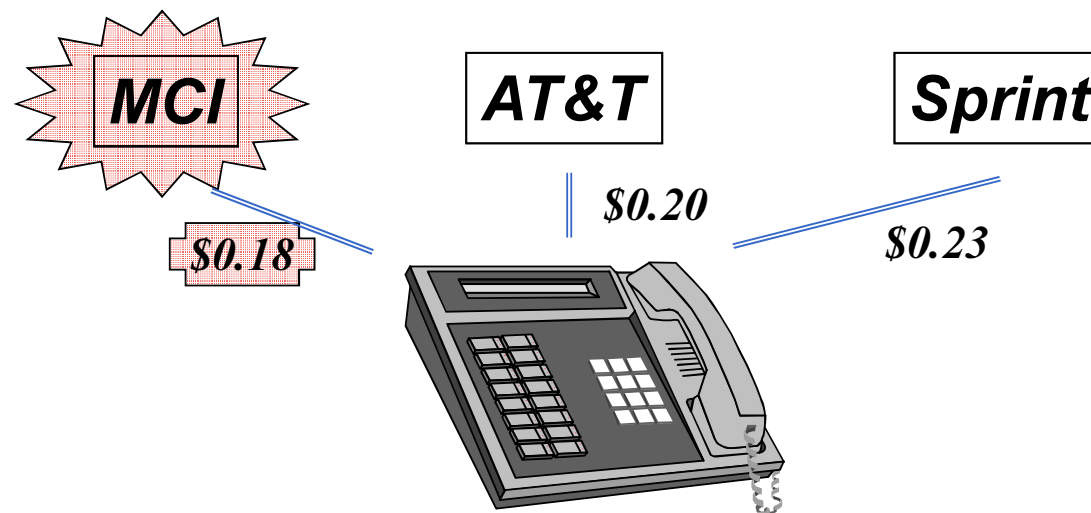
- Customer wishes to place long-distance call
- Carriers simultaneously bid, sending proposed prices
- Phone automatically chooses the carrier (dynamically)



32

Best Bid Wins

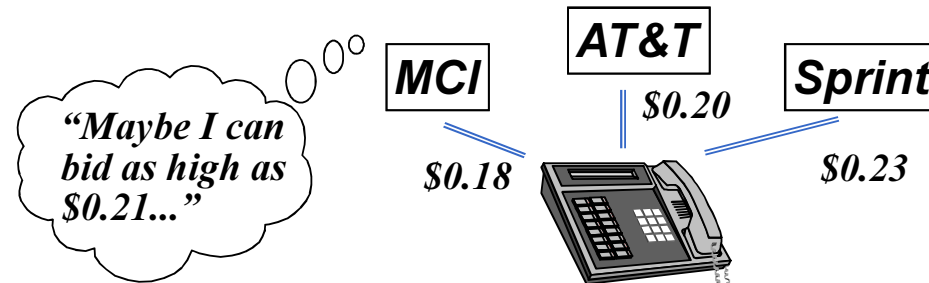
- Phone chooses carrier with lowest bid
- Carrier gets amount that it bid



Attributes of the Mechanism

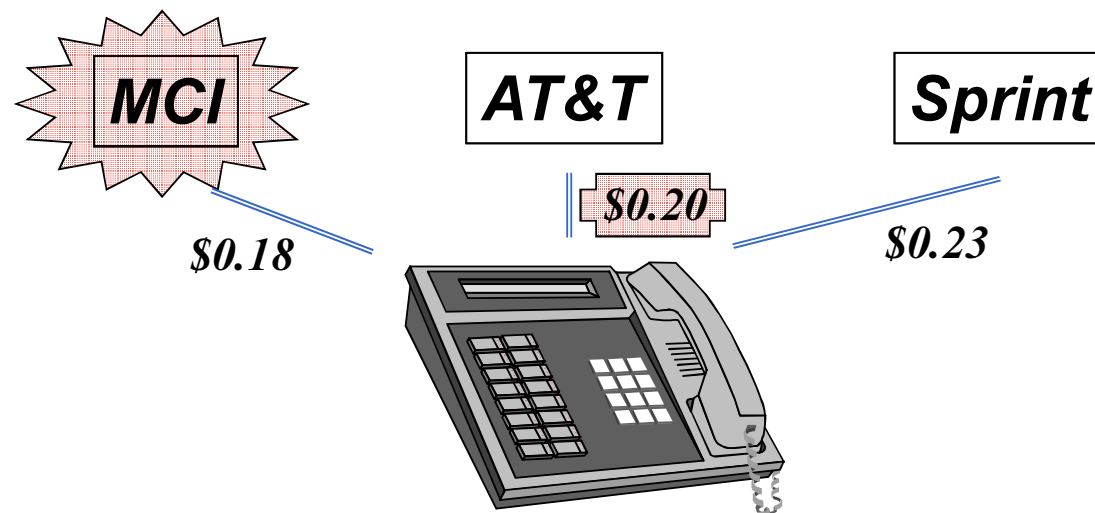
- ✓ *Distributed*
- ✓ *Symmetric*
- ✗ *Stable*
- ✗ *Simple*

Carriers have an incentive to invest effort in strategic behavior



Best Bid Wins, Gets Second Price (Vickrey Auction)

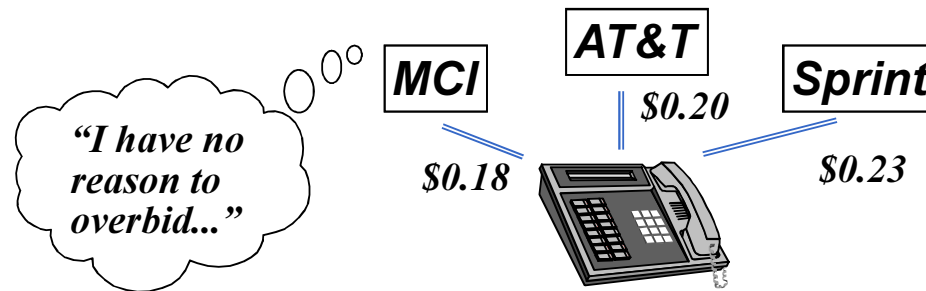
- Phone chooses carrier with lowest bid
- Carrier gets amount of second-best price



Attributes of the Vickrey Mechanism

- ✓ *Distributed*
- ✓ *Symmetric*
- ✓ *Stable*
- ✓ *Simple*

**Carriers have no
incentive to
invest effort in
strategic
behavior**

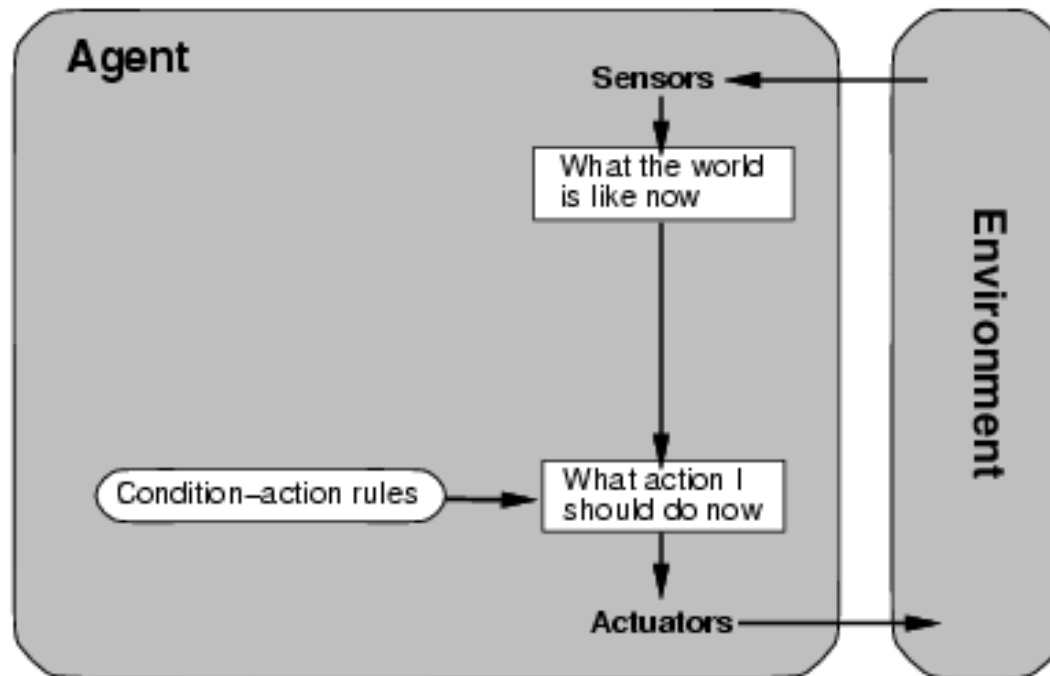


Agent Types

- Five basic types in order of increasing generality:
 - ◆ Simple reflex agents
 - ◆ Model-based reflex agents
 - ◆ Goal-based agents
 - ◆ Utility-based agents
 - ◆ Learning agents
see lecture *Machine Learning*

add features
↓

Simple Reflex Agents

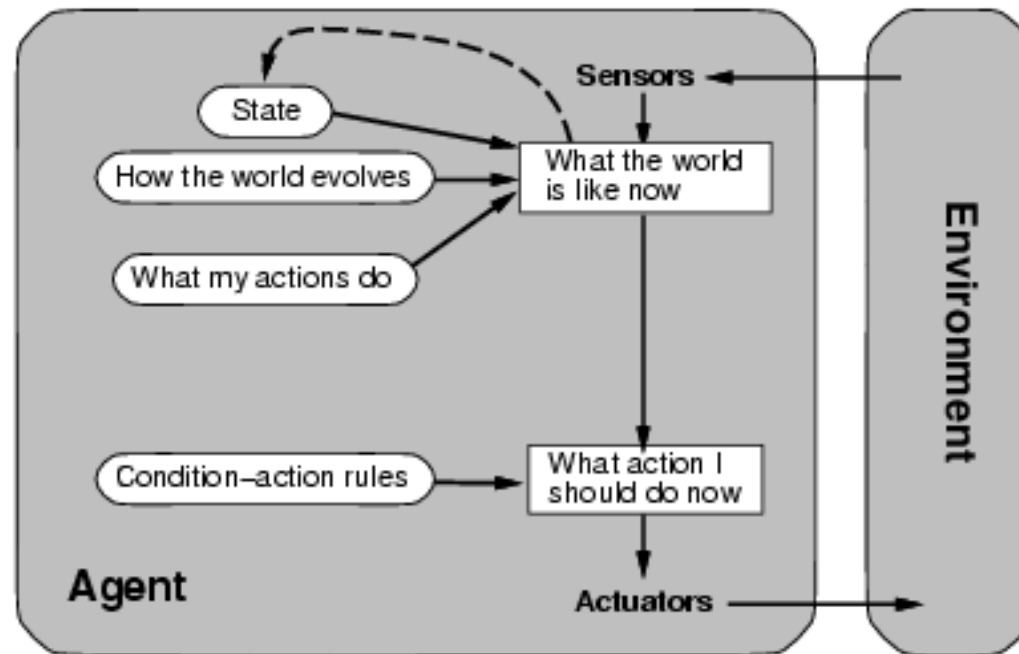


38

Simple Reflex Agent

- Drawbacks:
 - ◆ No autonomy
 - ◆ Decision depends on current percepts.
 - ◆ Sensitive to sensor fault

Model-Based Reflex Agents

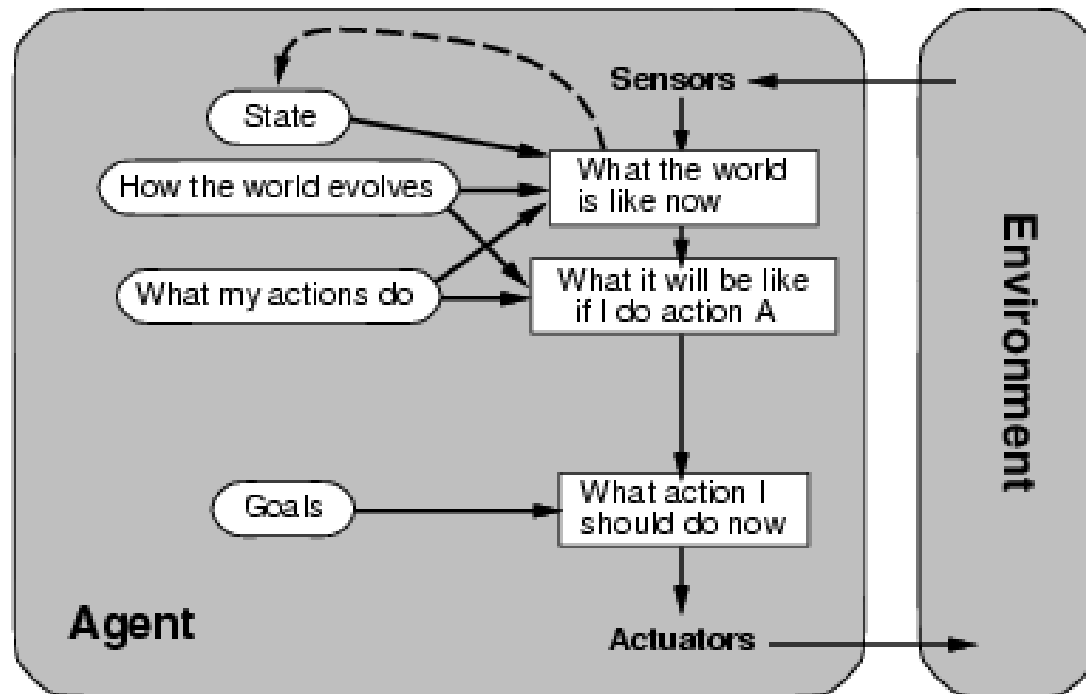


40

Goals for Agents

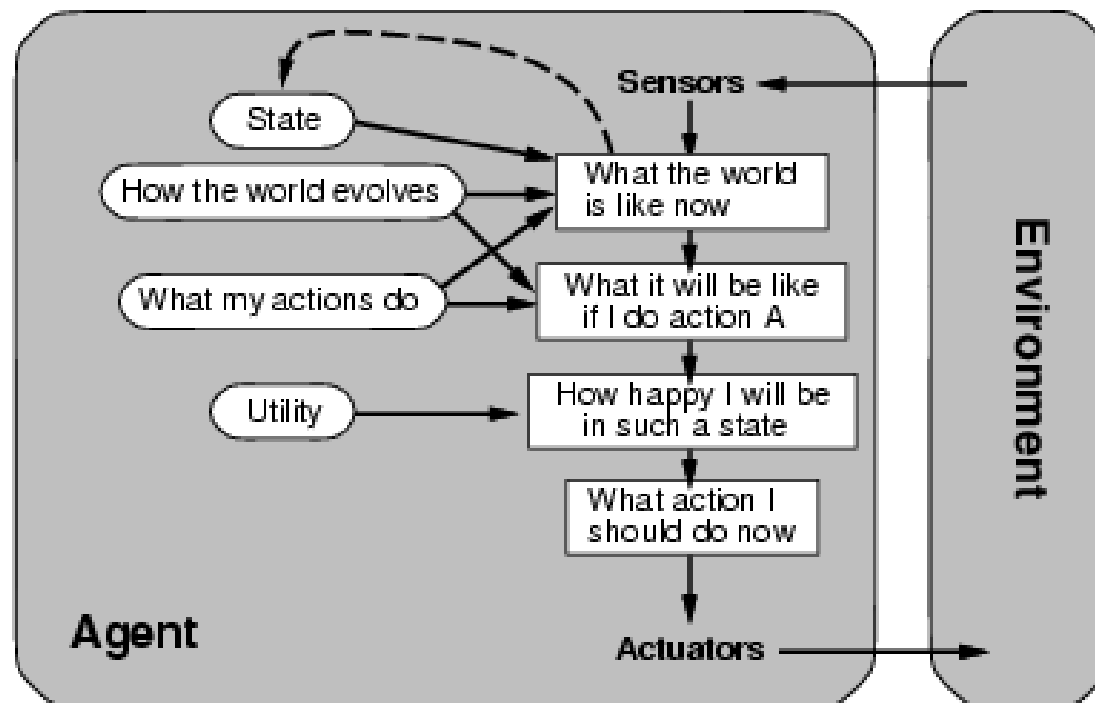
- We build agents in order to reach a goal for us
- The goals must be *specified* by us...
- But we want to tell agents what to do *without* telling them how to do it
→ Planning

Goal-Based Agents

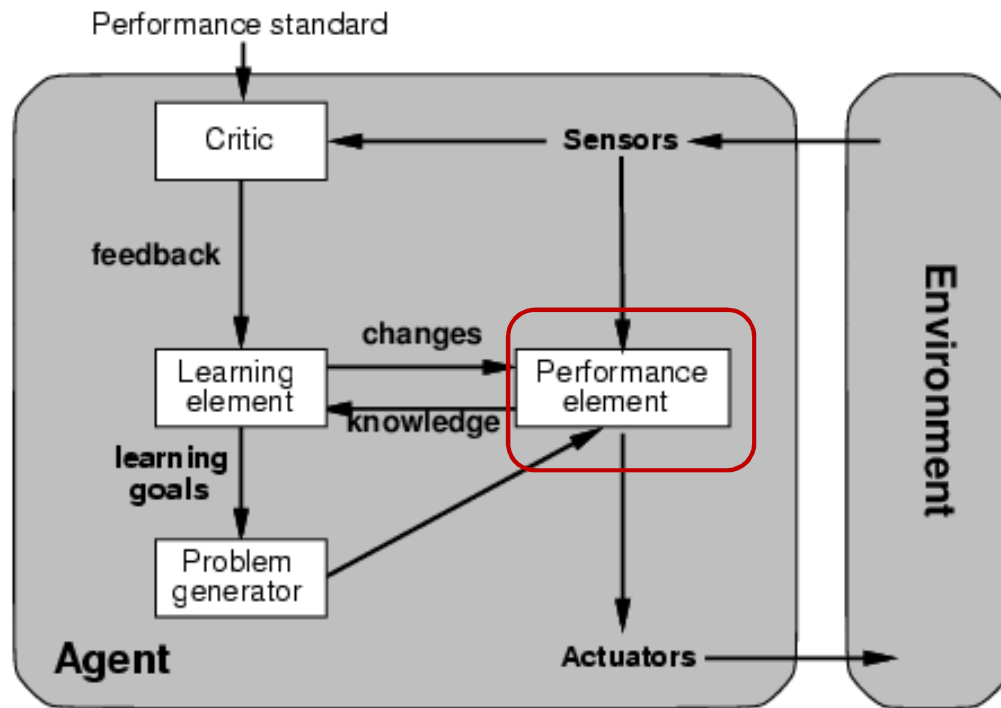


42

Utility-Based Agents



Learning Agents



45

Representation of agent states

- How to represent the states of agents?

- Atomic



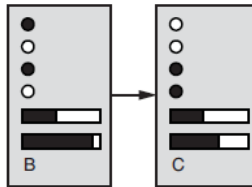
- State is a black box.

- Example:

Want to travel from city B to C. Cities are represented as names.

Representation of agent states

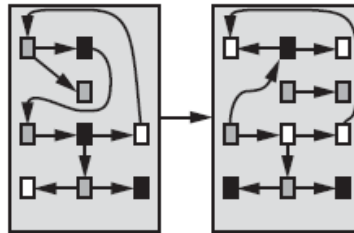
- Factored



- A state consists of a vector of attribute values
- Example: A car with GPS location, Fuel, radio station, ...

Representation of agent states

- Structured



- A state includes objects, each of which may have attributes of its own as well as relationships to other objects
- Example: natural language understanding

AI first approach: Deductive Reasoning Agents

- How can an agent decide what to do using theorem proving?
- Basic idea is to use logic to encode a theory stating the *best* action to perform in any given situation
- Let:
 - ♦ ρ be this theory (e.g. a set of rules)
 - ♦ Δ be a logical database that describes the current state of the world
 - ♦ Ac be the set of actions the agent can perform
 - ♦ $\Delta \vdash_{\rho} \phi$ mean that ϕ can be proven from Δ using ρ

49

Deductive Reasoning Agents

```
/* try to find an action explicitly prescribed */  
for each  $a \in Ac$  do  
  if  $\Delta \vdash_{\rho} Do(a)$  then  
    return  $a$   
  end-if  
end-for  
/* try to find an action not excluded */  
for each  $a \in Ac$  do  
  if  $\Delta \not\vdash_{\rho} \neg Do(a)$  then  
    return  $a$   
  end-if  
end-for  
return null /* no action found */
```

Deductive Reasoning Agents

- Problems:
 - ♦ how to convert video camera input to logical description?
 - ♦ decision making assumes a *static* environment
 - ♦ decision making using first-order logic is *undecidable*!
- Even when we use *propositional* logic, decision making in the worst case means solving NP-complete problems (= bad news!)

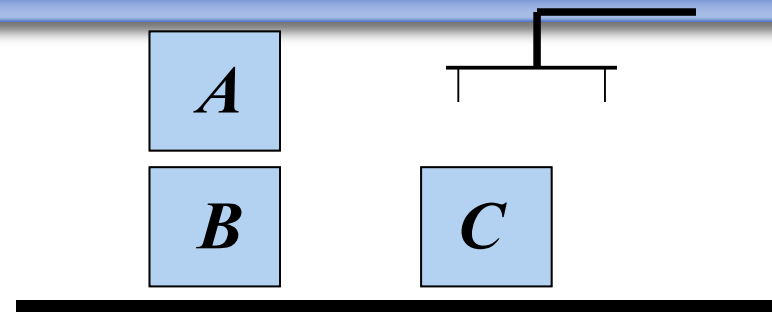
Practical Reasoning

- Practical reasoning consists of two activities:
 - ♦ *deliberation*
deciding *what* state of affairs we want to achieve
 - ♦ *means-ends reasoning*
deciding *how* to achieve these states of affairs

What is Means-End Reasoning?

- Basic idea is to give an agent:
 - ◆ representation of goal to achieve
 - ◆ representation of actions it can perform
 - ◆ representation of the environmentand have it generate a *plan* to achieve the goal
- Essentially, this is
automatic programming

The Blocks World



- We'll illustrate the techniques with the *blocks world*
- Contains a robot arm, 3 blocks (A, B, and C) of equal size, and a table

The Blocks World Ontology

- To represent this environment, need an *ontology*

On(x, y)

obj x on top of obj y

OnTable(x)

obj x is on the table

Clear(x)

nothing is on top of obj x

Holding(x)

arm is holding x

The Blocks World

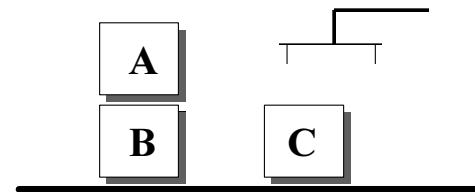
- Here is a representation of the blocks world described above:

Clear(A)

On(A, B)

OnTable(B)

OnTable(C)

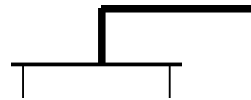


- Use the *closed world assumption*: anything not stated is assumed to be *false*

The Blocks World

- A *goal* is represented as a set of formula
- Here is a goal:

$OnTable(A) \wedge OnTable(B) \wedge OnTable(C)$



B

A

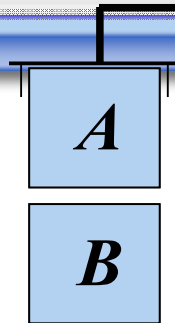
C

The Blocks World

- *Actions* are represented using a technique that was developed in the STRIPS planner
- Each action has:
 - ♦ a *name*
which may have arguments
 - ♦ a *pre-condition list*
list of facts which must be true for action to be executed
 - ♦ a *delete list*
list of facts that are no longer true after action is performed
 - ♦ an *add list*
list of facts made true by executing the action

Each of these may contain *variables*

The Blocks World Operators



- Example 1:
The *stack* action occurs when the robot arm places the object x it is holding is placed on top of object y .

	$Stack(x, y)$
pre	$Clear(y) \wedge Holding(x)$
del	$Clear(y) \wedge Holding(x)$
add	$ArmEmpty \wedge On(x, y)$

The Basic STRIPS Idea

- Place goal on goal stack:

Goal1

- Considering top Goal1, place onto it its subgoals:

GoalS1-2

GoalS1-1

Goal1

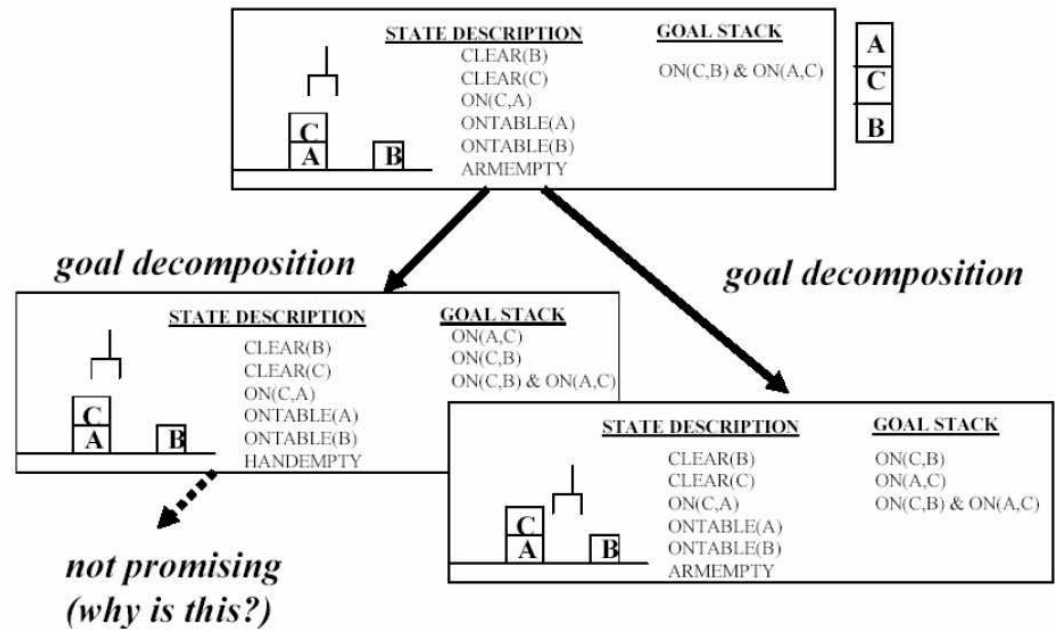
- Then try to solve subgoal GoalS1-2, and continue...

Stack Manipulation Rules, STRIPS

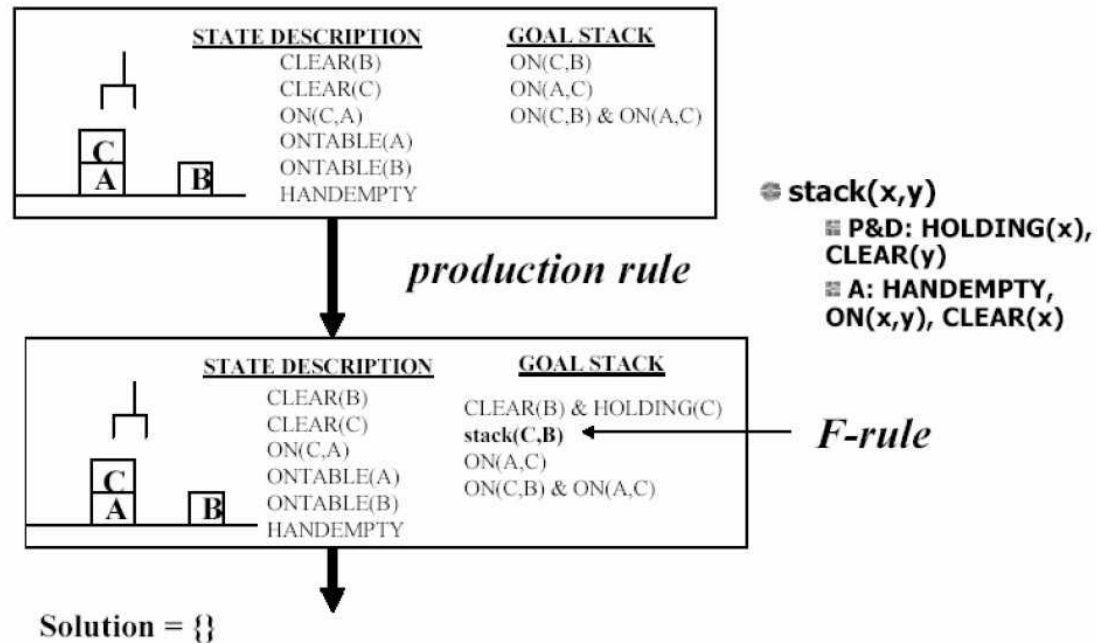
<u>If on top of goal stack:</u>	<u>Then do:</u>
Compound or single goal matching the current state description	Remove it
Compound goal <i>not</i> matching the current state description	<ol style="list-style-type: none">1. Keep original compound goal on stack2. List the unsatisfied component goals on the stack in some <i>new</i> order
Single-literal goal not matching the current state description	Find rule whose instantiated add-list includes the goal, and <ol style="list-style-type: none">1. Replace the goal with the instantiated rule;2. Place the rule's instantiated precondition formula on top of stack
Rule	<ol style="list-style-type: none">1. Remove rule from stack;2. Update database using rule;3. Keep track of rule (for solution)
Nothing	Stop

62

STRIPS in Action

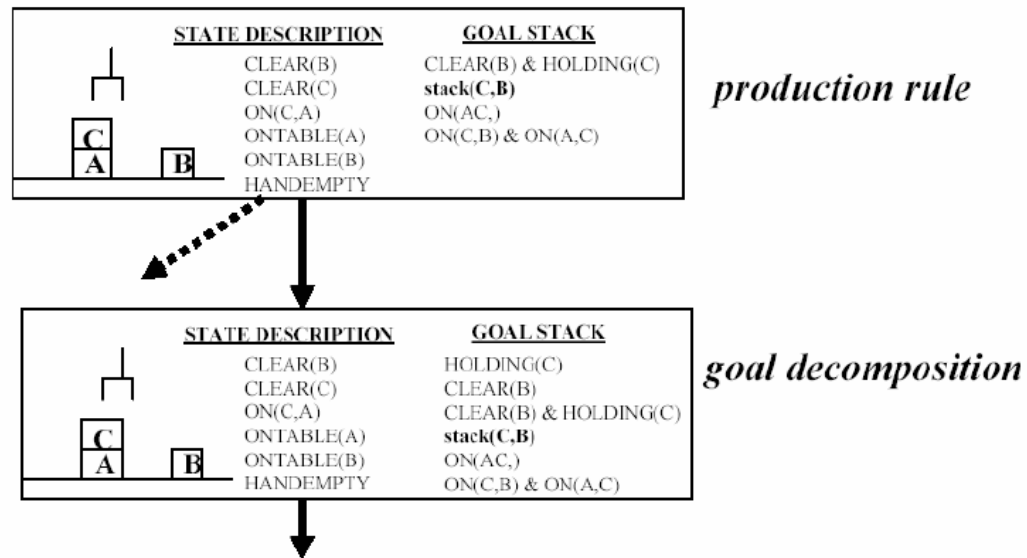


STRIPS in Action



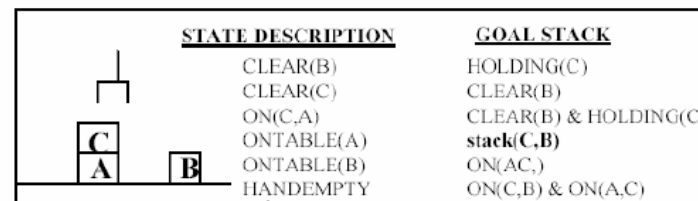
64

STRIPS in Action

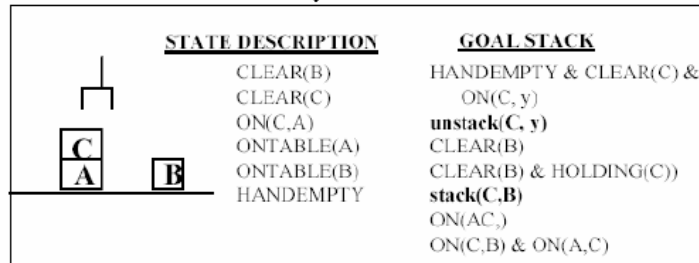


Solution = {}

STRIPS in Action



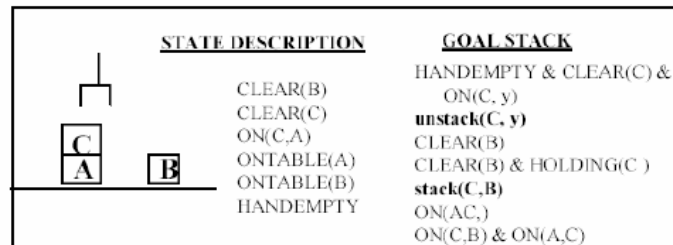
production rule



Solution = {}

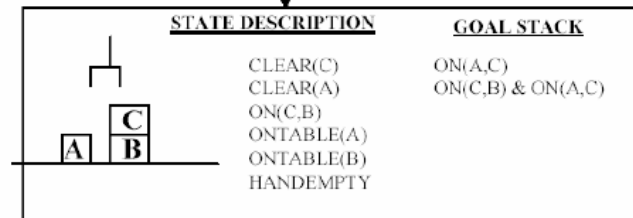
- **unstack(x,y)**
 - P&D:
HANDEEMPTY,
CLEAR(x), ON(x,y)
 - A: HOLDING(x),
CLEAR(y)

STRIPS in Action



☞ **unstack(x,y)**
 ☞ P&D: HANDEEMPTY,
 CLEAR(x), ON(x,y)
 ☞ A: HOLDING(x),
 CLEAR(y)

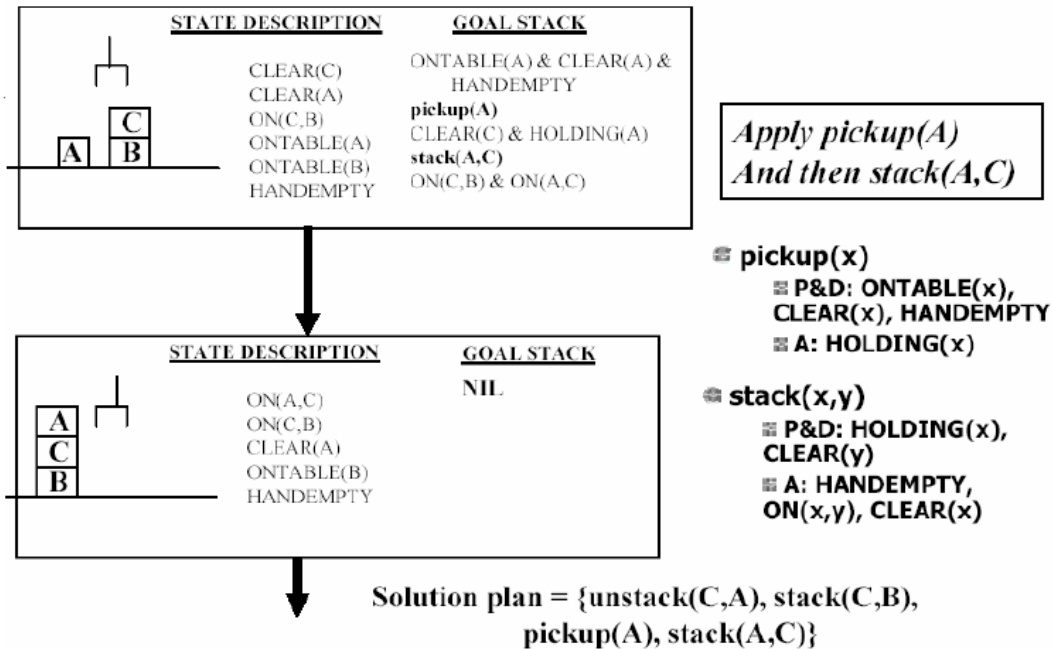
*Substitute {A/y}, then apply
 unstack(C,A) then stack(C,B)*



☞ **stack(x,y)**
 ☞ P&D: HOLDING(x),
 CLEAR(y)
 ☞ A: HANDEEMPTY,
 ON(x,y), CLEAR(x)

Solution = {unstack(C,A), stack(C,B)}

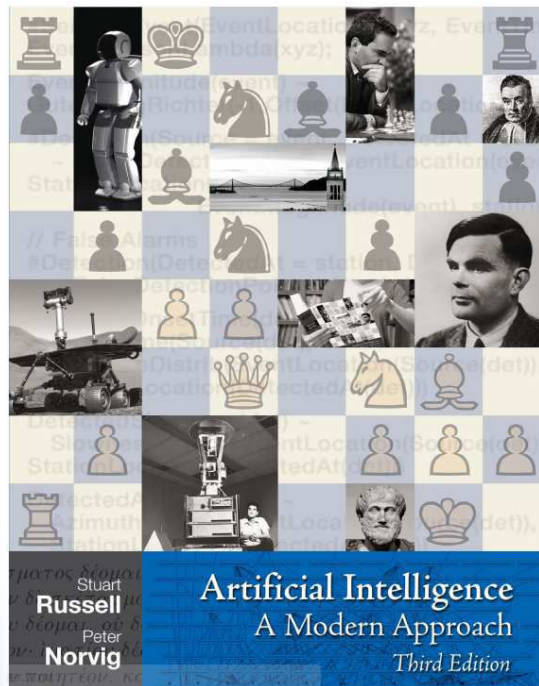
STRIPS in Action



**Intelligent Autonomous Agents
and Cognitive Robotics
Solving problems by searching**

Rainer Marrone
Hamburg University of Technology
Slides based on Hwee Tou Ng's

Literature



- Chapter 3

Problem types

- **Single-state problem: Deterministic, fully observable**
 - ◆ Agent knows exactly which state it will be in; can calculate optimal action sequence to reach the goal
- **Multiple state problem: Deterministic, partially/not observable**
 - ◆ Agent must reason about sequences of actions and states assumed while working towards goal state.
- **Contingency problem: Nondeterministic and partially observable**
 - ◆ Percepts provide **new** information about current state
 - ◆ Solution is a contingent plan or policy
 - ◆ Often interleave search and execution
- **Exploration problem: Unknown state space**
 - ◆ Discover and learn about environment while taking actions

Example: vacuum world

- **Single-state**, start in #5.

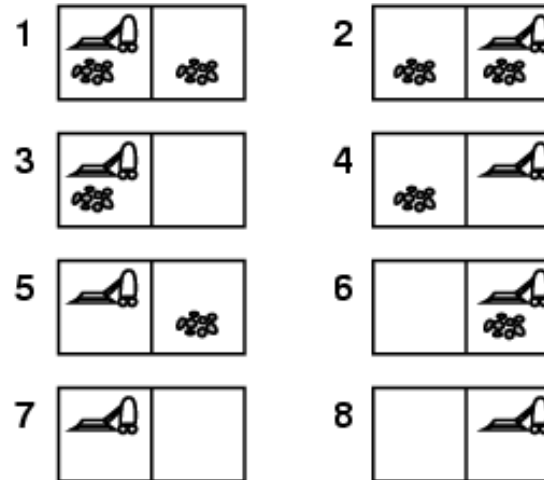
Solution?

[Right, Suck]

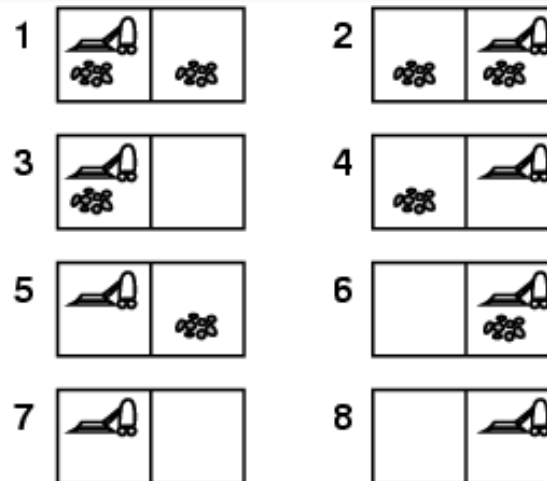
- **Multiple-state**, start in {1,2,3,4,5,6,7,8} e.g.,
Right goes to {2,4,6,8}

Solution?

[Right, Suck, Left, Suck]



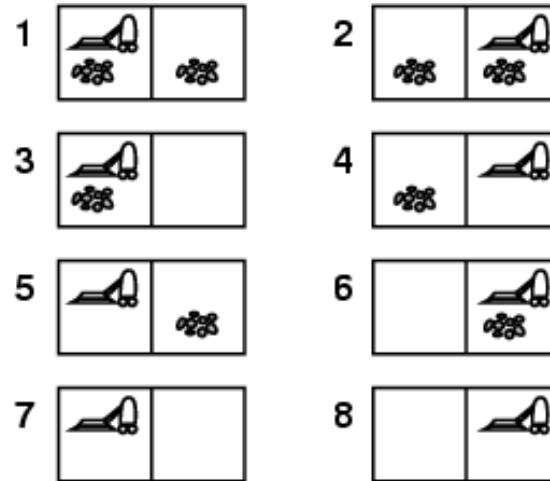
Example: vacuum world



- **Contingency**
 - ◆ Nondeterministic: *Suck* may dirty a clean carpet
 - ◆ Partially observable: location, dirt at current location.
 - ◆ Percept: $[L, Clean]$, i.e., start in #5 or #7

Solution

Example: vacuum world



- Contingency

- ◆ Nondeterministic: *Suck* may dirty a clean carpet
- ◆ Partially observable: location, dirt at current location.
- ◆ Percept: $[L, Clean]$, i.e., start in #5 or #7

Solution? $[Right, \textit{if dirt then Suck}]$

Solving problems by searching

- We will discuss solutions for all the different settings.
- We start with simple searches and modify them for more complex settings

Tree search algorithms

- Basic idea:
 - ♦ offline, simulated exploration of state space by generating successors of already-explored states (a.k.a. **expanding** states)

```
Function Tree-Search(problem, strategy) returns a solution, or failure
  initialize the search tree using the initial state problem
  loop do
    if there are no candidates for expansion then return failure
    choose a leaf node for expansion according to strategy
    if the node contains a goal state then return the corresponding solution
    else expand the node and add resulting nodes to the search tree
  end
```

8

Measuring search performance

- A search strategy is defined by picking the **order of node expansion**
- Strategies are evaluated along the following dimensions:
 - ♦ **completeness**: does it always find a solution if one exists?
 - ♦ **time complexity**: number of nodes generated
 - ♦ **space complexity**: maximum number of nodes in memory
 - ♦ **optimality**: does it always find a least-cost solution?
- Time and space complexity are measured in terms of
 - ♦ *b*: maximum branching factor of the search tree
 - ♦ *d*: depth of the least-cost solution
 - ♦ *m*: maximum depth of the state space (may be ∞)

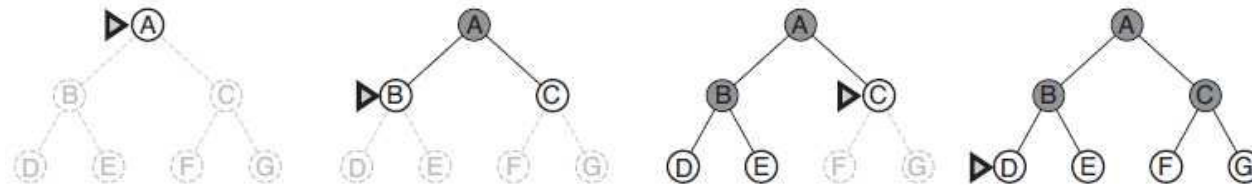
Uninformed search strategies

Uninformed search strategies use only the information available in the problem definition

- Breadth-first search
- Depth-first search
- Depth-limited search
- Iterative deepening search

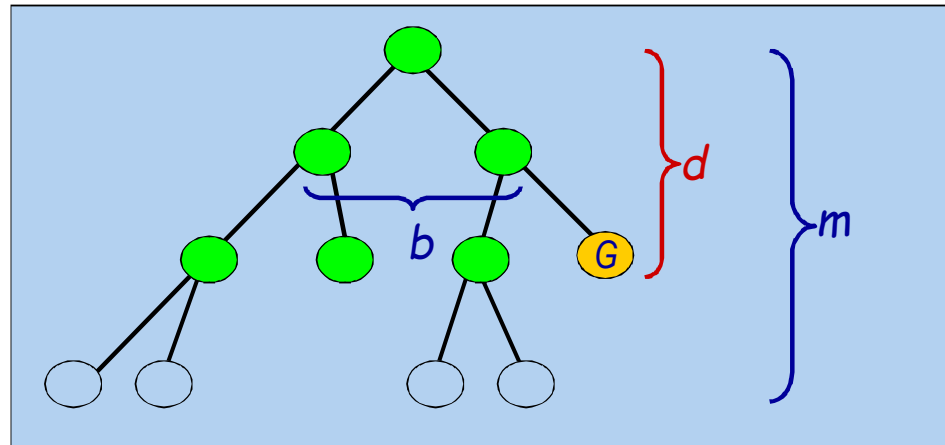
Breadth-first search

- Expand shallowest unexpanded node
- **Implementation:**
 - ♦ *fringe* is a FIFO queue, i.e., new successors go at end



Time complexity of breadth-first search

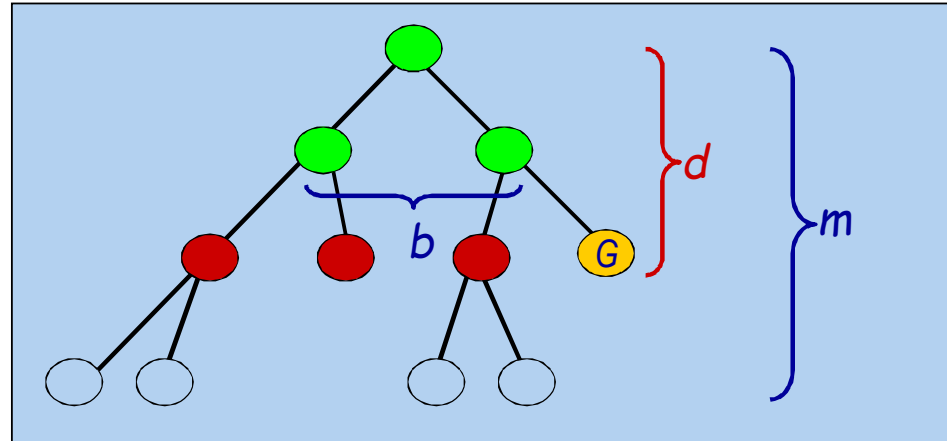
- If a goal node is found on depth d of the tree, all nodes up till that depth are created.



- Thus: $O(b^d)$

Space complexity of breadth-first

- *Largest number of nodes in QUEUE is reached on the level d of the goal node.*



- *QUEUE contains all ● and ● nodes. (Thus: 4).*
- *In General: b^d*

Properties of breadth-first search

- Complete? Yes (if b is finite)
- Time? $1+b+b^2+b^3+\dots +b^d = O(b^d)$
- Space? $O(b^{d+1})$ (keeps every node in memory)
 $O(b^d)$ (*only fringe*)
- Optimal? Yes (*if cost = 1 per step*)

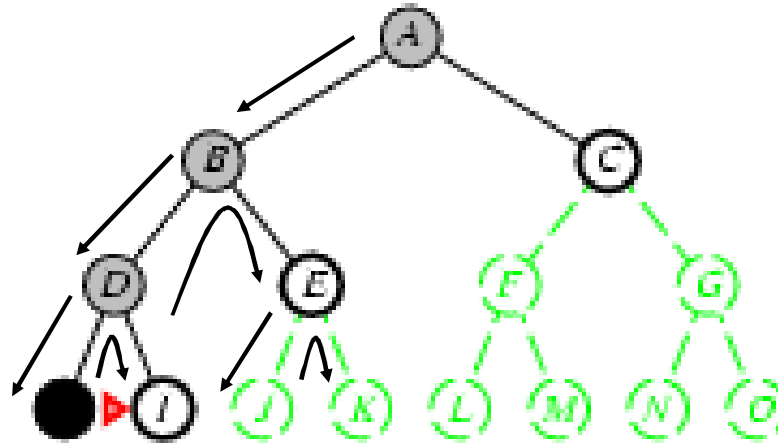
Complexity example

Depth	Nodes	Time	Memory
2	110	.11 milliseconds	107 kilobytes
4	11,110	11 milliseconds	10.6 megabytes
6	10^6	1.1 seconds	1 gigabyte
8	10^8	2 minutes	103 gigabytes
10	10^{10}	3 hours	10 terabytes
12	10^{12}	13 days	1 petabyte
14	10^{14}	3.5 years	99 petabytes
16	10^{16}	350 years	10 exabytes

Figure 3.13 Time and memory requirements for breadth-first search. The numbers shown assume branching factor $b = 10$; 1 million nodes/second; 1000 bytes/node.

Depth-first search

- Expand deepest unexpanded node



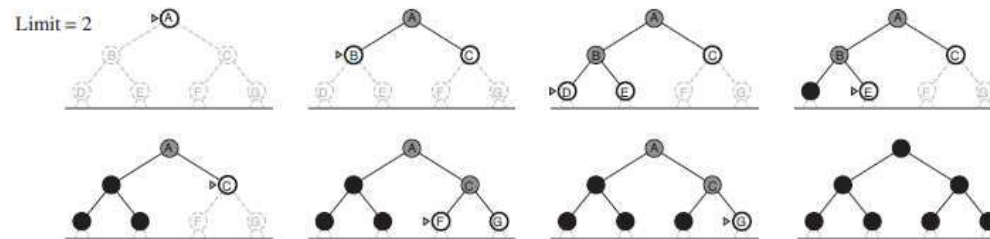
16

Properties of depth-first search

- Complete? No: fails in infinite-depth spaces
→ complete in finite spaces
- Time? $O(b^m)$: terrible if m is much larger than d
 - ◆ but if solutions are dense, may be much faster than breadth-first
- Space? $O(bm)$, i.e., linear space!
- Optimal? No

Depth-limited search

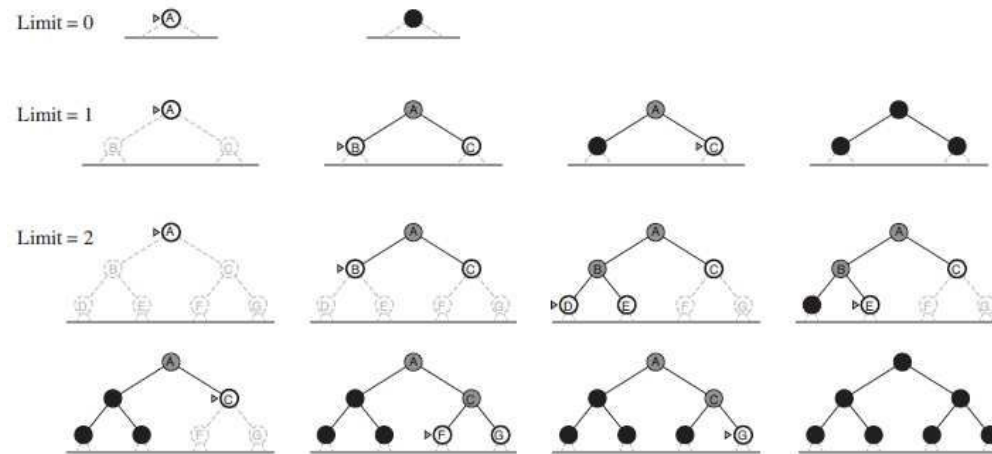
- depth-first search with depth limit l , i.e., nodes at depth l have no successors
- Solves infinite path problem
- Incomplete if $l < d$ (shallowest goal node)
- Nonoptimal if $l > d$



Iterative deepening search

```
function ITERATIVE-DEEPENING-SEARCH(problem) returns a solution, or fail-  
ure  
  inputs: problem, a problem  
  for depth ← 0 to ∞ do  
    result ← DEPTH-LIMITED-SEARCH(problem, depth)  
    if result ≠ cutoff then return result
```

Iterative deepening search



20

Iterative deepening search

- Number of nodes generated in a depth-limited search to depth d with branching factor b :

$$N_{DLS/BFS} = b^0 + b^1 + b^2 + \dots + b^{d-2} + b^{d-1} + b^d$$

- Number of nodes generated in an iterative deepening search to depth d with branching factor b :

$$N_{IDS} = (d+1)b^0 + d b^1 + (d-1)b^2 + \dots + 3b^{d-2} + 2b^{d-1} + 1b^d$$

- For $b = 10$, $d = 5$,
 - ♦ $N_{DLS} = 1 + 10 + 100 + 1,000 + 10,000 + 100,000 = 111,111$
 - ♦ $N_{IDS} = 6 + 50 + 400 + 3,000 + 20,000 + 100,000 = 123,456$
- Overhead = $(123,456 - 111,111)/111,111 = 11\%$

Properties of iterative deepening search

- Complete? Yes
- Time? $(d+1)b^0 + d b^1 + (d-1)b^2 + \dots + b^d = O(b^d)$
- Space? $O(bd)$
- Optimal? Yes, if step cost = 1

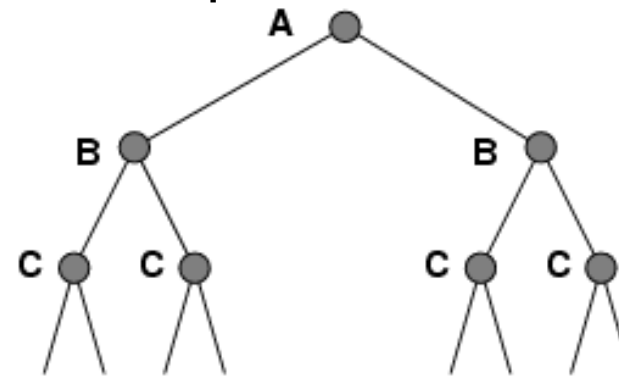
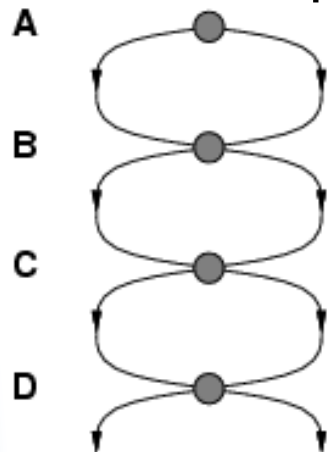
Summary of algorithms

Criterion	Breadth-First	Uniform-Cost	Depth-First	Depth-Limited	Iterative Deepening	Bidirectional (if applicable)
Complete?	Yes ^a	Yes ^{a,b}	No	No	Yes ^a	Yes ^{a,d}
Time	$O(b^d)$	$O(b^{1+\lceil C^*/\epsilon \rceil})$	$O(b^m)$	$O(b^l)$	$O(b^d)$	$O(b^{d/2})$
Space	$O(b^d)$	$O(b^{1+\lceil C^*/\epsilon \rceil})$	$O(bm)$	$O(bl)$	$O(bd)$	$O(b^{d/2})$
Optimal?	Yes ^c	Yes	No	No	Yes ^c	Yes ^{c,d}

Figure 3.21 Evaluation of tree-search strategies. b is the branching factor; d is the depth of the shallowest solution; m is the maximum depth of the search tree; l is the depth limit. Superscript caveats are as follows: ^a complete if b is finite; ^b complete if step costs $\geq \epsilon$ for positive ϵ ; ^c optimal if step costs are all identical; ^d if both directions use breadth-first search.

Repeated states

- Failure to detect repeated states can turn a linear problem into an exponential one!



Graph search

```
function GRAPH-SEARCH(problem, fringe) returns a solution, or failure
  closed ← an empty set
  fringe ← INSERT(MAKE-NODE(INITIAL-STATE[problem]), fringe)
  loop do
    if fringe is empty then return failure
    node ← REMOVE-FRONT(fringe)
    if GOAL-TEST[problem](STATE[node]) then return SOLUTION(node)
    if STATE[node] is not in closed then
      add STATE[node] to closed
      fringe ← INSERTALL(EXPAND(node, problem), fringe)
```

Remember nodes visited

Beyond classical search

- Informed search
 - ◆ Greedy best-first search
 - ◆ A* search
- Admissible heuristics, creating heuristics
- Local search algorithms
 - ◆ Hill-climbing search
 - ◆ Simulated annealing search
 - ◆ Local beam search
 - ◆ Genetic algorithms
- Searching with nondeterministic actions

Best-first search

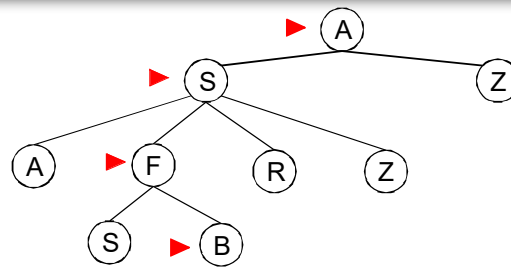
- Idea: use a **heuristic evaluation function** $f(n)$ for each node
 - ♦ estimate of "desirability"
 - Expand most desirable unexpanded node
- Implementation:
Order the nodes in fringe in decreasing order of desirability
- Special cases:
 - ♦ greedy best-first search
 - ♦ A* search

Greedy best-first search

- Evaluation function
 $f(n) = h(n)$ (**h**euristic)= estimate of cost from n to *goal*
e.g., $h_{SLD}(n)$ = straight-line distance from n to goal node
- Greedy best-first search expands the node that **appears** to be closest to the goal
- Stop if the goal node appears on the fringe

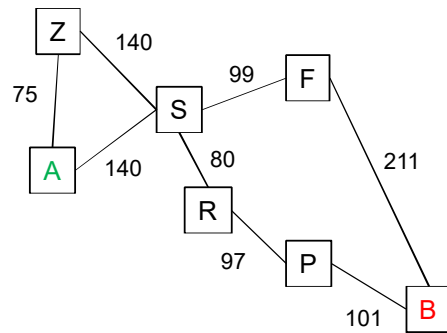
28

Greedy best-first search example: Go from **A** to **B**



140
 99
 211

 450



	B
A	366
Z	374
S	253
R	193
P	100
F	176
B	0

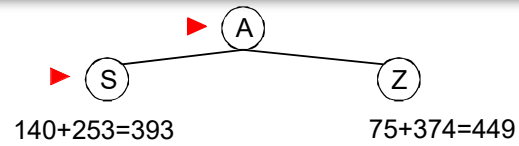
Properties of greedy best-first search

- Complete? No – can get stuck in loops, but can use graph search
- Time? $O(b^m)$, but a good heuristic can give dramatic improvement
- Space? $O(b^m)$ -- keeps all nodes in memory
- Optimal? No

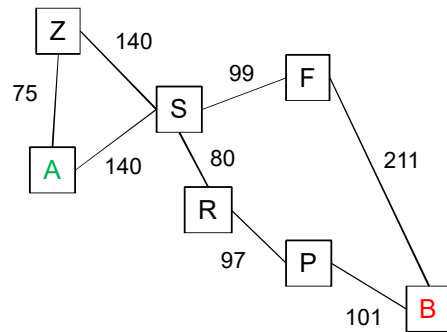
A* search

- Idea: avoid expanding nodes that are already expensive
- Evaluation function $f(n) = g(n) + h(n)$
 $g(n)$ = cost so far to reach n
 $h(n)$ = estimated cost from n to goal
- Goal node must also be expanded

A* search example: Go from A to B



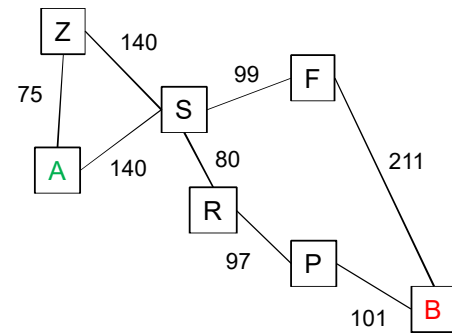
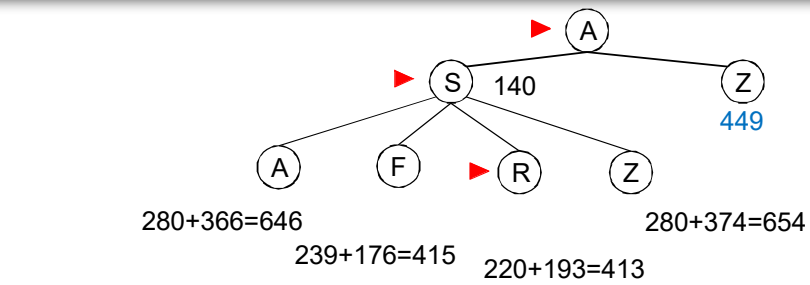
$$f(n) = g(n) + h(n)$$



	B
A	366
Z	374
S	253
R	193
P	100
F	176
B	0

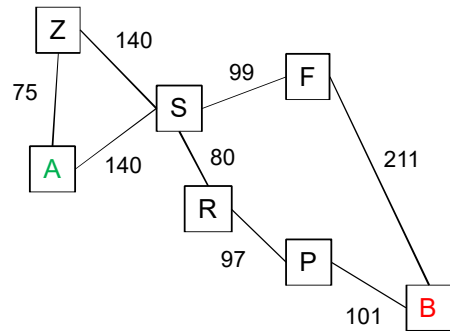
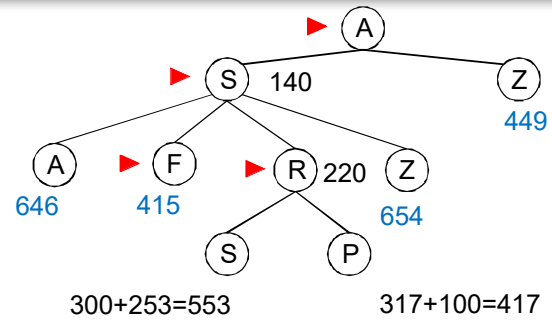
33

A* search example: Go from A to B



	B
A	366
Z	374
S	253
R	193
P	100
F	176
B	0

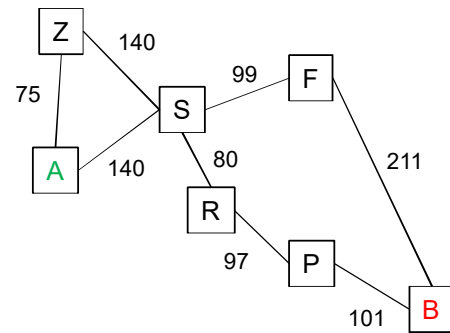
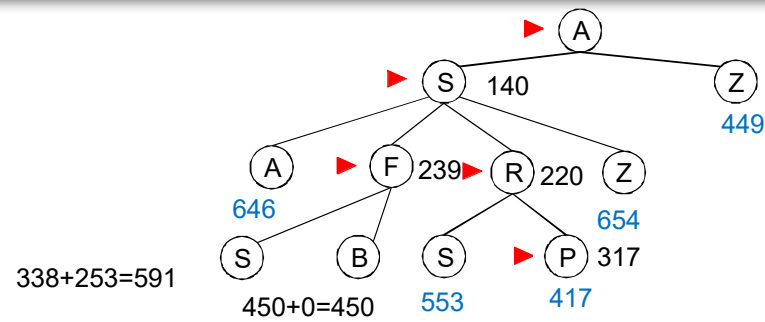
A* search example: Go from A to B



	B
A	366
Z	374
S	253
R	193
P	100
F	176
B	0

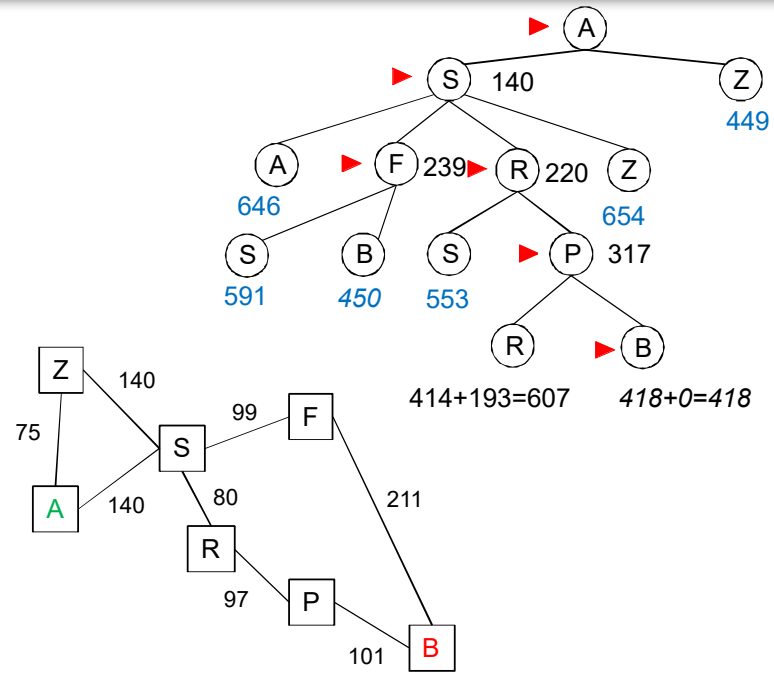
35

A* search example: Go from A to B



	B
A	366
Z	374
S	253
R	193
P	100
F	176
B	0

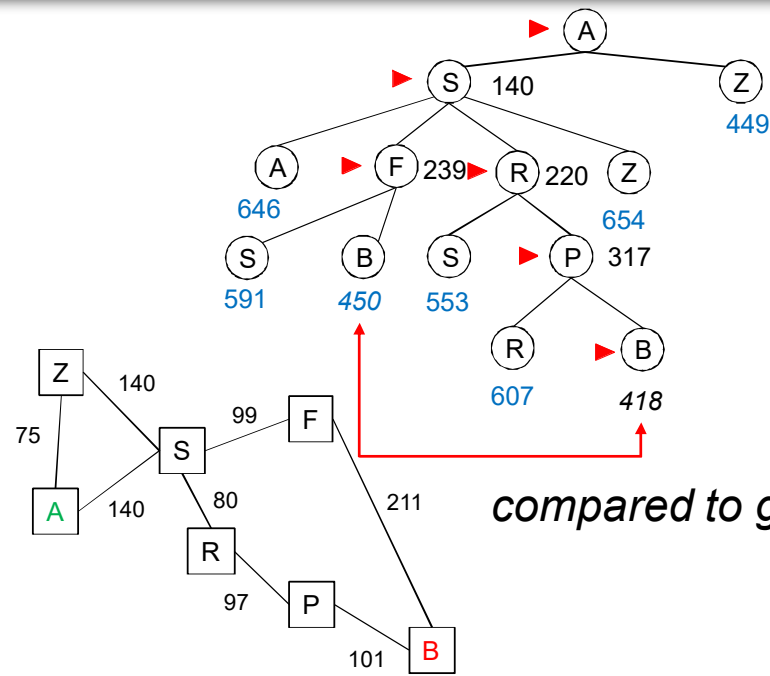
A* search example: Go from A to B



	B
A	366
Z	374
S	253
R	193
P	100
F	176
B	0

37

A* search example: Go from A to B



	B
A	366
Z	374
S	253
R	193
P	100
F	176
B	0

Admissible heuristics

- A heuristic $h(n)$ is **admissible** if for every node n , $h(n) \leq h^*(n)$, where $h^*(n)$ is the **true** cost to reach the goal state from n . An admissible heuristic **never overestimates** the cost to reach the goal, i.e., it is **optimistic**
- Example: $h_{SLD}(n)$ (never overestimates the actual road distance)
- **Theorem**: If $h(n)$ is admissible, A^* using TREE-SEARCH is optimal
- For graph searches we need a stronger criteria

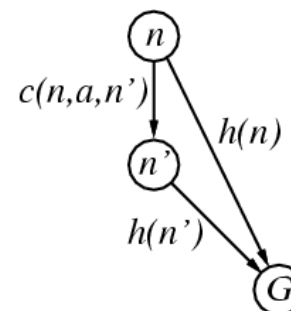
Consistent heuristics

- “each side of a triangle cannot be longer than the sum of the other two sides”
- A heuristic is **consistent** if for every node n , every successor n' of n generated by any action a ,

$$h(n) \leq c(n,a,n') + h(n')$$

- If h is consistent, we have

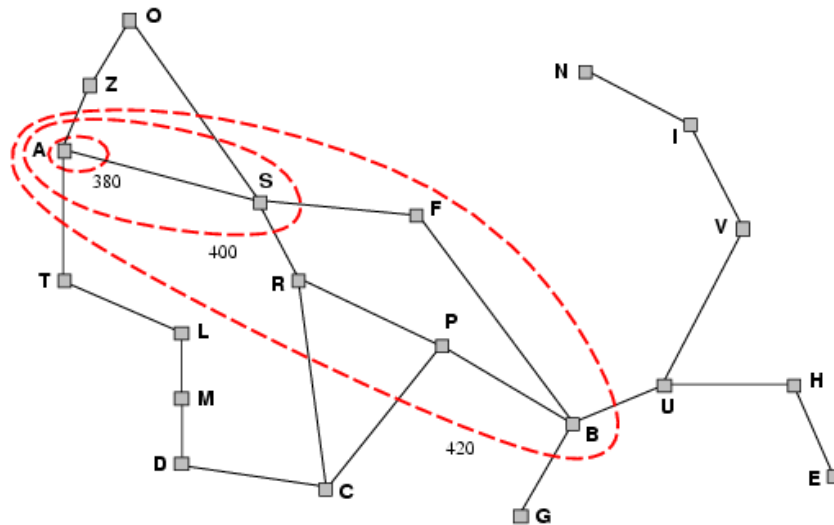
$$\begin{aligned} f(n') &= g(n') + h(n') \\ &= g(n) + c(n,a,n') + h(n') \\ &\geq g(n) + h(n) \\ &= f(n) \end{aligned}$$



- i.e., $f(n)$ is non-decreasing along any path.
- **Theorem:** If $h(n)$ is consistent, A* using GRAPH-SEARCH is optimal

Optimality of A*

- A* expands nodes in order of increasing f value
- A* will search all path with $f(n) < C^*$ (completeness)
- A* never expands nodes with $f > C^*$ (the true cost)



Map, showing contours at $f=380$, $f=400$, and $f=420$.

Properties of A*

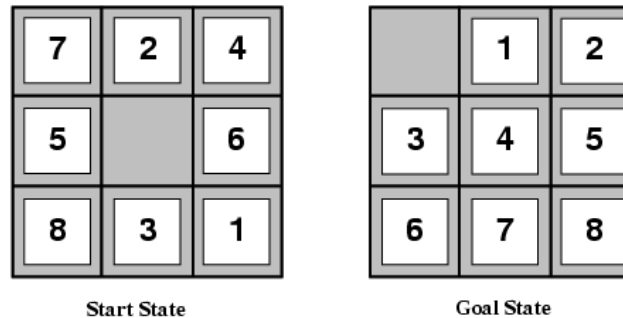
- Complete? Yes
- Time? The number of states in the goal contour can still be exponential.
- Space?
Keeps all generated nodes in memory,
as do all graph search algorithms.
- Optimal? Yes

Not practical for very large scale problems

Admissible heuristics

E.g., for the 8-puzzle:

- $h_1(n)$ = number of misplaced tiles
- $h_2(n)$ = total Manhattan distance
(i.e., no. of squares from desired location of each tile)



- $h_1(S) = ?$ 8
- $h_2(S) = ?$ $3+1+2+2+2+3+3+2 = 18$

Empirical Evaluation

- d = distance from goal
- Average over 100 instances
- IDS: Iterative Deepening Search (the best you can do without any heuristic)

nodes expanded

d	Search Cost		
	IDS	$A^*(h_1)$	$A^*(h_2)$
2	10	6	6
4	112	13	12
6	680	20	18
8	6384	39	25
10	47127	93	39
12	364404	227	73
14	3473941	539	113
16	–	1301	211
18	–	3056	363
20	–	7276	676
22	–	18094	1219
24	–	39135	1641

Dominance

- If $h_2(n) \geq h_1(n)$ for all n (both admissible) then h_2 **dominates** h_1
- *Is h_2 always better than h_1 ?*
- $f(n) < C^*$ (true cost)
- Every node $h(n) < C^* - g(n)$ will surely get expanded
- Because $h_2(n) \geq h_1(n)$ every node of h_2 will also be expanded from h_1 , and h_1 will cause other nodes to be expanded

45

Relaxed problems

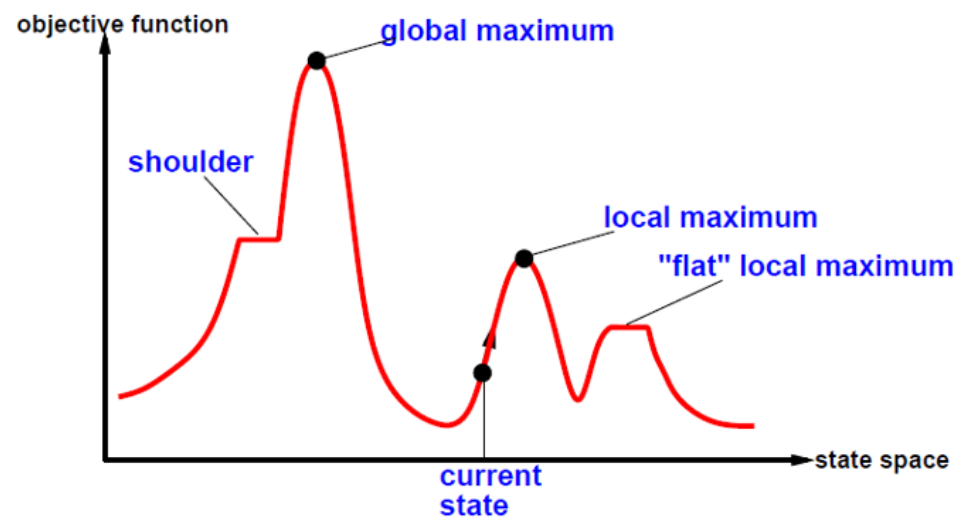
- A problem with fewer restrictions on the actions is called a **relaxed problem**
- The cost of an optimal solution to a relaxed problem is an admissible heuristic for the original problem
- If the rules of the 8-puzzle are relaxed so that a tile can move **anywhere**, then $h_1(n)$ gives the shortest solution
- If the rules are relaxed so that a tile can move to **any adjacent square**, then $h_2(n)$ gives the shortest solution

Local search algorithms

- In many optimization problems, the **path** to the goal is irrelevant; the goal state itself is the solution
- State space = set of "complete" configurations
- Find configuration satisfying constraints, e.g., n-queens; integrated-circuit design; factory-floor layout,
- In such cases, we can use **local search algorithms**. Keep a single "current" state, try to improve it

State space and objective function

Useful to consider state space landscape



Hill-climbing search

- "Like climbing Everest in thick fog with amnesia" (Russell, Norvig)

```
function HILL-CLIMBING(problem) returns a state that is a local maximum
  inputs: problem, a problem
  local variables: current, a node
                  neighbor, a node

  current ← MAKE-NODE(INITIAL-STATE[problem])
  loop do
    neighbor ← a highest-valued successor of current
    if VALUE[neighbor] ≤ VALUE[current] then return STATE[current]
    current ← neighbor
  end
```

50

Hill-climbing search: 8-queens problem

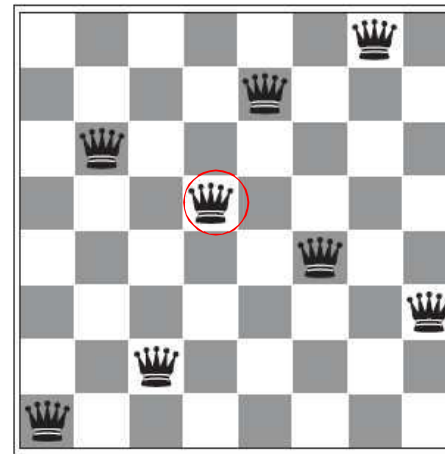
18	12	14	13	13	12	14	14
14	16	13	15	12	14	12	16
14	12	18	13	15	12	14	14
15	14	14	♙	13	16	13	16
♙	14	17	15	♙	14	16	16
17	♙	16	18	15	♙	15	♙
18	14	♙	15	15	14	♙	16
14	14	13	17	12	14	12	18

- The successors of a state are all possible states generated by moving a single queen to another square in the same column (so each state has $8 \times 7 = 56$ successors)
- Cost function: h = number of pairs of queens that are attacking each other, either directly or indirectly
- $h = 17$ for the above state, best moves are marked.

Hill-climbing search: 8-queens problem

18	12	14	13	13	12	14	14
14	16	13	15	12	14	12	16
14	12	18	13	15	12	14	14
15	14	14	♙	13	16	13	16
♙	14	17	15	♙	14	16	16
17	♙	16	18	15	♙	15	♙
18	14	♙	15	15	14	♙	16
14	14	13	17	12	14	12	18

(a)



(b)

Figure 4.3 (a) An 8-queens state with heuristic cost estimate $h = 17$, showing the value of h for each possible successor obtained by moving a queen within its column. The best moves are marked. (b) A local minimum in the 8-queens state space; the state has $h = 1$ but every successor has a higher cost.

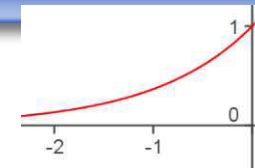
Observations

- Get stuck 86% vs 14% success
- Taking 4 steps only if successful
3 steps if getting stuck (17 Million states)

- If sideways are allowed (100), success in 94%. Increase of cost 21 steps.
- Variants
 - Stochastic hill climbing
 - First-choice hill climbing
 - Random restart

Simulated annealing search

Idea: escape local maxima by allowing some “bad” moves
but gradually decrease their size and frequency

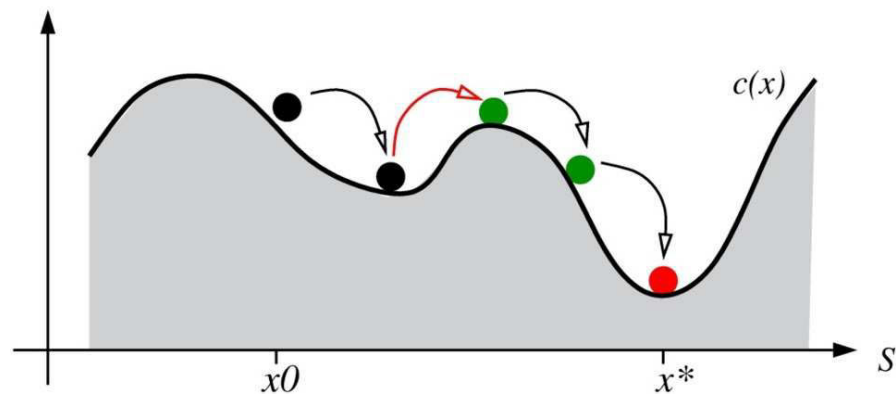


```
function SIMULATED-ANNEALING(problem, schedule) returns a solution state
  inputs: problem, a problem
           schedule, a mapping from time to “temperature”
  local variables: current, a node
                   next, a node
                   T, a “temperature” controlling prob. of downward steps

  current ← MAKE-NODE(INITIAL-STATE[problem])
  for t ← 1 to ∞ do
    T ← schedule[t]
    if T = 0 then return current  Stopping criteria
    next ← a randomly selected successor of current
     $\Delta E$  ← VALUE[next] - VALUE[current]
    if  $\Delta E > 0$  then current ← next
    else current ← next only with probability  $e^{\Delta E/T}$ 
```

54

Simulated Annealing



Properties of simulated annealing search

- One can prove:
If T decreases slowly enough, then simulated annealing search will find a global optimum/minimum with probability approaching 1
- Widely used in VLSI layout, airline scheduling, etc.

Local beam search

- Keep track of k states rather than just one
- Start with k randomly generated states
- At each iteration, all the successors of all k states are generated
- If any one is a goal state, stop; else select the k best successors from the complete list and repeat.

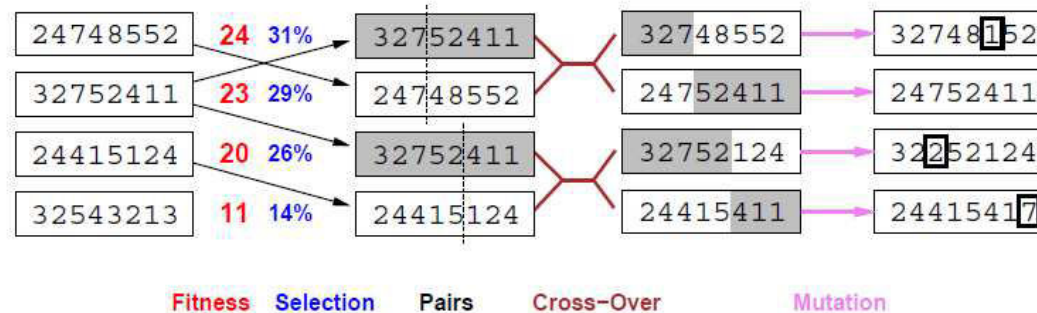
Genetic algorithms

- A variant of stochastic beam search. But a successor state is generated by different operations.
- Start with k randomly generated states (**population**)
- A state is represented as a string over a finite alphabet (often a string of 0s and 1s)
- Evaluation function (**fitness function**). Higher values for better states.
- Produce the next generation of states by selection, crossover, and mutation

58

Genetic algorithms

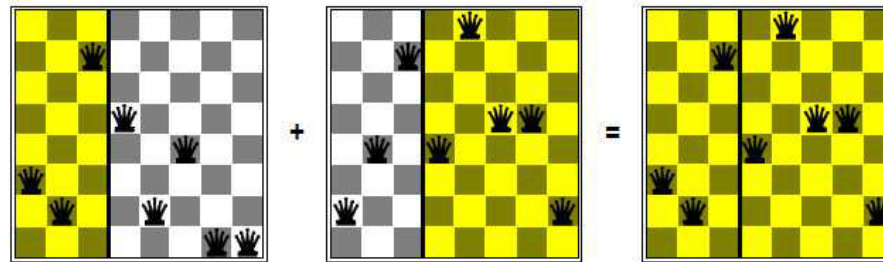
= stochastic local beam search + generate successors from **pairs** of states



- Fitness function: number of non-attacking pairs of queens (min = 0, max = 7 + 6 + 5 + 4 + 3 + 2 + 1 = 28)
- $24/(24+23+20+11) = 31\%$
- $23/(24+23+20+11) = 29\%$ etc

Genetic algorithms

Crossover helps iff substrings are meaningful components



GAs \neq evolution: e.g., real genes encode replication machinery!

- *How many crossover, mutations*
- *How to encode the problem, fitness function*
- *One (more popular) vs. two child's*

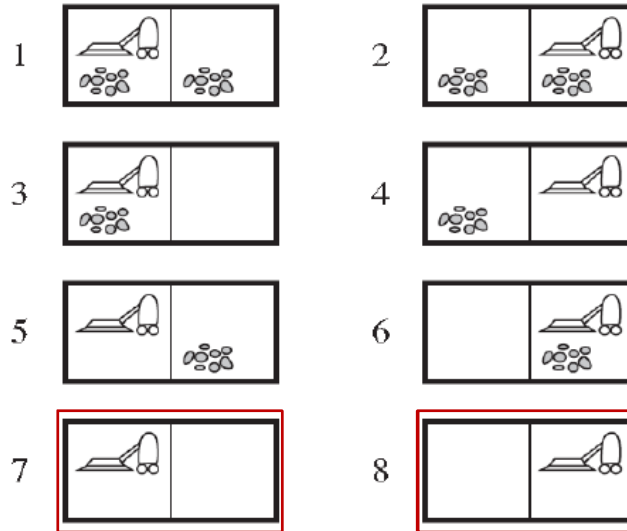
60

Nondeterministic/Uncertain actions

- What if the outcome of actions is non deterministic
- Erratic vacuum cleaner
 - ♦ When applied to a *dirty* square the square is *cleaned and adjacent square sometimes also*.
 - ♦ When applied to a *clean square*, *sometimes dirt is deposited* on that square
 - ➔ need to have **contingency plan/strategy**

Possible states

- The eight possible states of the erratic vacuum world – states 7 and 8 are goal states

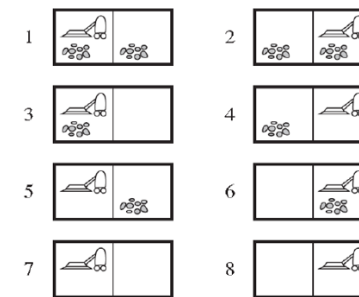


62

Multiple States

- The result of an action is a set of states
- *Suck* in state 1 returns the set {5,7}
- We also need to generalize the concept of solution, since for example, if we **start** in state **1** there is no single sequence of actions to solve the problem instead we need a contingency plan like:

[*Suck*, if *State*=5 then [*Right*, *Suck*] else []]



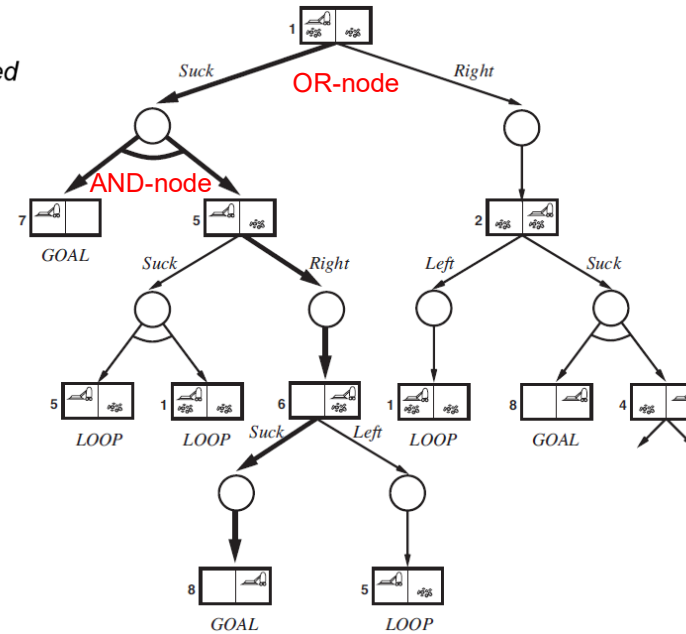
AND-OR Search trees

- Branching is also introduced by the environment choice of the outcome of actions.

When applied to a *dirty* square the square is *cleaned* and *adjacent square sometimes also*.

When applied to a *clean* square, *sometimes dirt is deposited* on that square

- This leads to **AND-OR trees**
- The bold path is the current plan



64

AND-OR Search trees

- A solution is a subtree
 - ◆ has a goal node at every leaf
 - ◆ specifies one action at OR-nodes
 - ◆ Includes every outcome branch at AND-nodes
- Leads to *if then else* or *case* if more than two outcomes

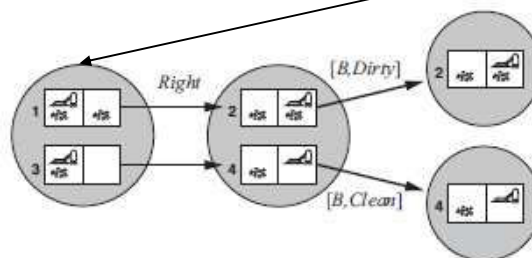
AND-OR Search trees

- Can also be explored by BFS and best-first methods
- Heuristic functions must be modified to estimate cost of a contingent solution rather than a sequence
- The notion of admissibility carries over.

Partial Observable Env.

- The vacuum cleaner has only partial information, e.g., if he is in the left square he does not see the state of the right square.

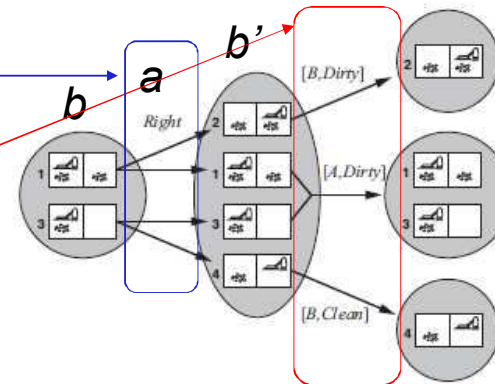
If the initial state is left and dirt, we have a **belief state** rather than a physical state



- But we also have uncertain actions: Move action may fail

Uncertain actions & partial observable

- Prediction:
 $b' = \text{Predict}(b, a)$
- Possible observations in b'
 $\text{Percepts}(b') = \{o : o = \text{PERCEPT}(b')\}$

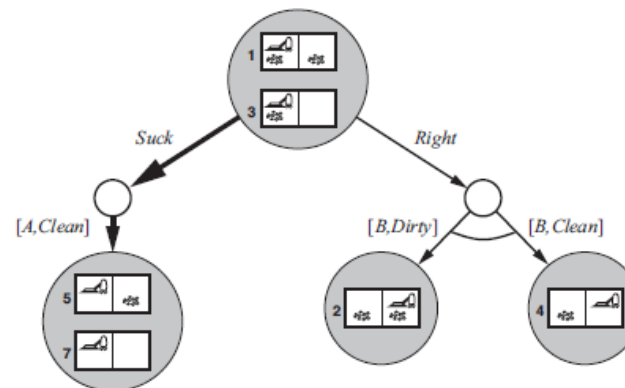


- Update of belief state:
 $b_o = \text{UPDATE}(b', o) = \{s : o = \text{PERCEPT}(s) \text{ and } s \in b'\}$
- Putting all together:

$$\text{RESULTS}(b, a) = \{b_o : b_o = \text{UPDATE}(\text{PREDICT}(b, a), o) \text{ and } o \in \text{POSSIBLE-PERCEPTS}(\text{PREDICT}(b, a))\}$$

Structure

- Can use different search structures
- E.g. And-Or-Graphs



Intelligent Autonomous Agents and Cognitive Robotics: Adversarial Agents

Ralf Möller, Rainer Marrone
Hamburg University of Technology

Adversarial Agents

- In this chapter we cover **competitive environments**, in which the agents goals are in conflict, giving rise to adversarial search problems often known as games.
- Mathematical game theory, a branch of economics, views any multi-agent environment as a game, regardless of whether the agents are cooperative or competitive.

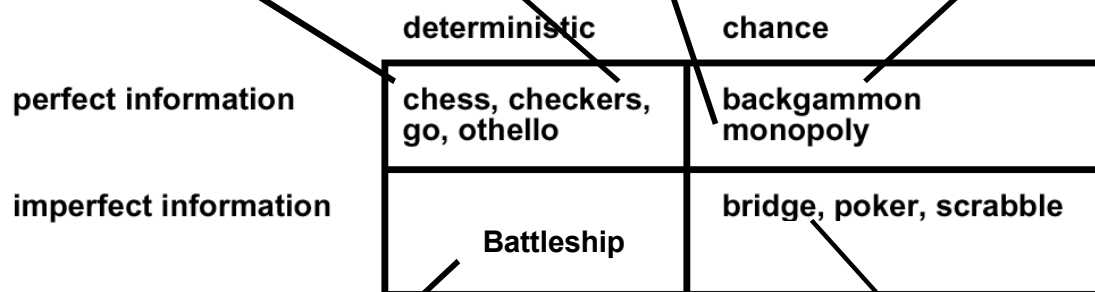
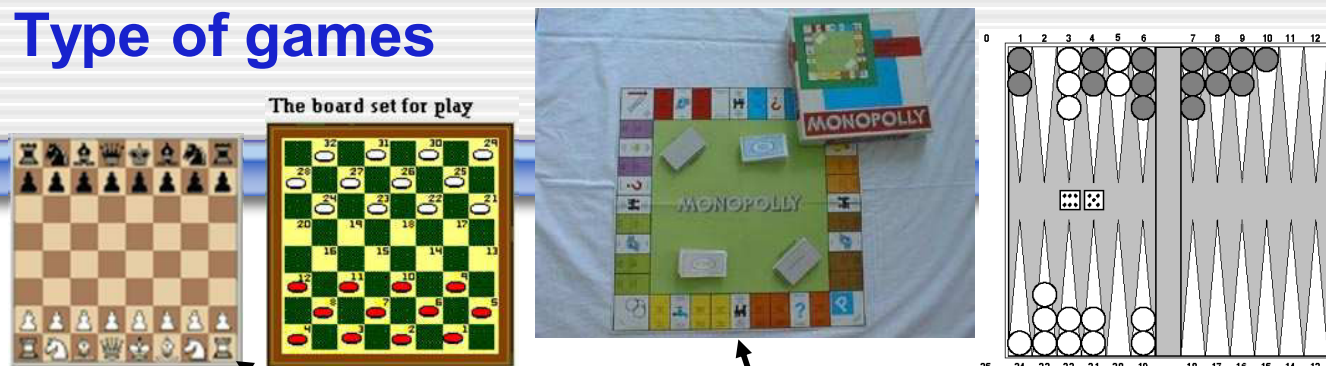
Multi-Agent Games

- Agents must ***anticipate*** what other agents do
- Criteria:
 - ♦ **Abstraction:** To describe a game we must capture every *relevant* aspect of the game.
 - ♦ **Accessible environments:** Such games are characterized by perfect information
 - ♦ **Search:** game-playing then consists of a search through possible game positions *with actions of other agents*
 - ♦ **Unpredictable opponent:** introduces **uncertainty** thus game-playing must deal with **contingency problems**

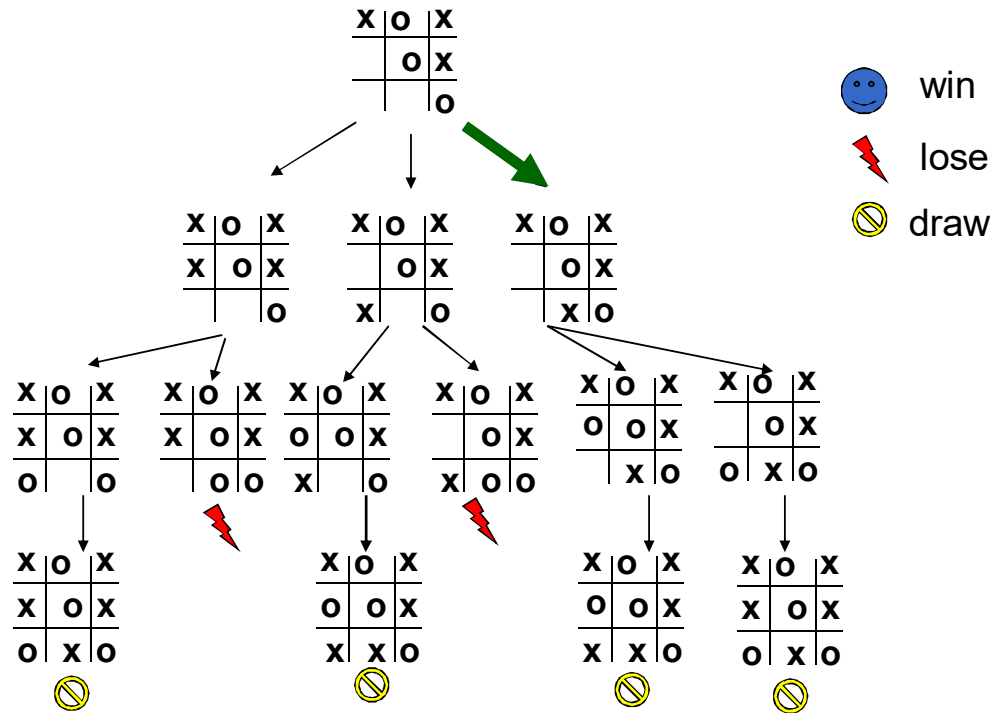
Two-player games

- A game formulated as a *search problem*:
 - ◆ *Initial state*: board position and turn
 - ◆ *Actions/Transition model*: definition of legal moves
 - ◆ *Terminal state*: conditions for when game is over
 - ◆ ***Utility function***:
a numeric value that describes the outcome of the game. E.g., -1, 0, 1 for loss, draw, win (AKA **payoff function**)

Type of games



What is a good move?



The minimax algorithm

- Perfect play for deterministic environments with perfect information
- **Basic idea:** choose move with highest minimax value
= best achievable payoff against **best play**
- **Algorithm:**
 1. Generate game tree completely
 2. Determine utility of each terminal state
 3. Propagate the utility values upward in the tree by applying MIN and MAX operators on the nodes in the current level
 4. At the root node use minimax decision to select the move with the max (of the min) utility value

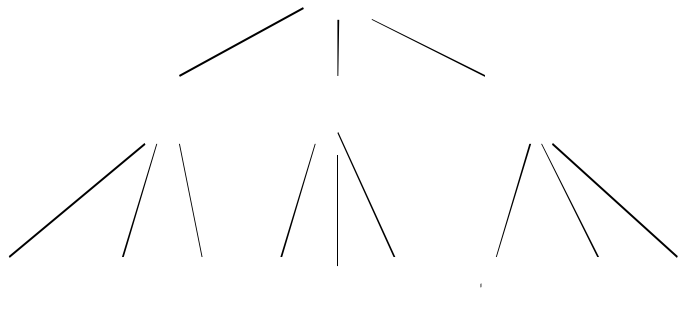
Minimax algorithm



MAX



MIN



- Minimize opponent's chance
- Maximize your chance

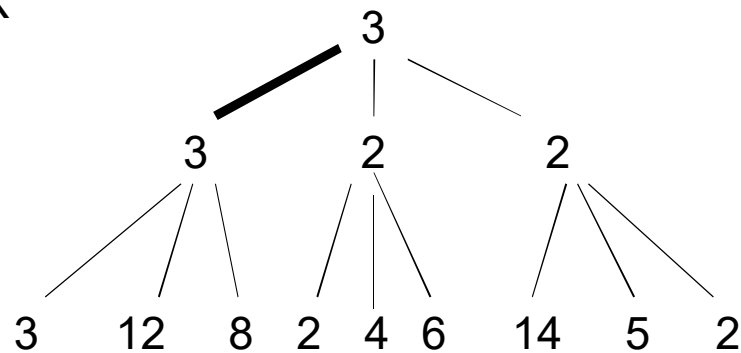
Minimax



MAX



MIN



- Minimize opponent's chance
- Maximize your chance

MINIMAX-VALUE(n) =

$$\begin{cases} \text{UTILITY}(n) & \text{if } n \text{ is a terminal state} \\ \max_{s \in \text{Successors}(n)} \text{MINIMAX-VALUE}(s) & \text{if } n \text{ is a MAX node} \\ \min_{s \in \text{Successors}(n)} \text{MINIMAX-VALUE}(s) & \text{if } n \text{ is a MIN node.} \end{cases}$$

Minimax: Recursive implementation

```
function MINIMAX-DECISION(state) returns an action  
  return  $\arg \max_{a \in \text{ACTIONS}(s)} \text{MIN-VALUE}(\text{RESULT}(\text{state}, a))$ 
```

```
function MAX-VALUE(state) returns a utility value  
  if TERMINAL-TEST(state) then return UTILITY(state)  
   $v \leftarrow -\infty$   
  for each a in ACTIONS(state) do  
     $v \leftarrow \text{MAX}(v, \text{MIN-VALUE}(\text{RESULT}(s, a)))$   
  return v
```

```
function MIN-VALUE(state) returns a utility value  
  if TERMINAL-TEST(state) then return UTILITY(state)  
   $v \leftarrow \infty$   
  for each a in ACTIONS(state) do  
     $v \leftarrow \text{MIN}(v, \text{MAX-VALUE}(\text{RESULT}(s, a)))$   
  return v
```

Complete: Yes, for finite state-space **Time complexity:** $O(b^m)$
Optimal: Yes, if winning is the goal **Space complexity:** $O(bm)$ or $O(m)$

10

Game vs. search problem

- Unpredictable opponent → contingency plan (MINIMAX assumes best playing opponent)
- Time limits → cannot explore complete state space, approximate
- Pruning (McCarthy, 1956)
- Finite horizon, approximate (Zuse, 1945; Shannon 1950,...)

Searching for the next move

- **Complexity:** many games have a huge search space
 - ♦ **Chess:** $b = 35, m=100 \Rightarrow \text{nodes} = 35^{100}$
means more than 10^{154} in a search tree and more than 10^{40} nodes in a search graph. Take several millennia to compute moves. $35^{100} = 10^{\log(35^{100})} = 10^{100 \cdot \log(35)} = 10^{100 \cdot 1,54} = 10^{154}$
- **Resource (e.g., time, memory) limit:** optimal solution not feasible/possible, thus must approximate
 1. **Pruning:** makes the search more efficient by discarding portions of the search tree that cannot improve quality.
 2. **Evaluation functions:** heuristics to evaluate utility of a state without exhaustive search.

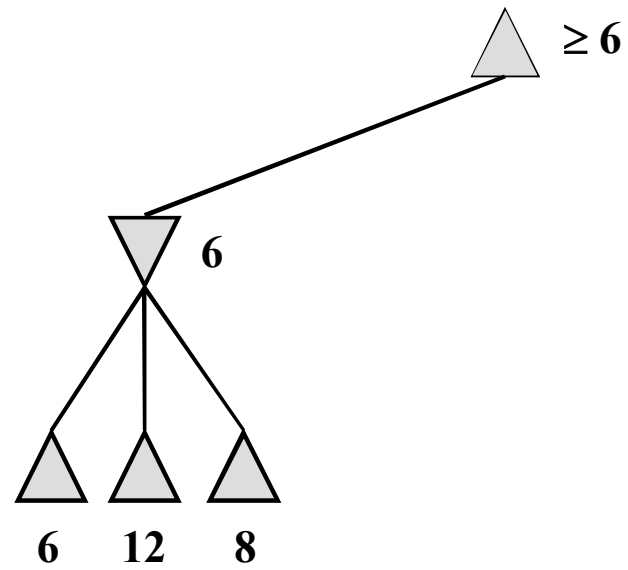
1. α - β pruning

- **α - β pruning**: the basic idea is to prune portions of the search tree that cannot improve the utility value of the *max* or *min* node, by just considering the values of nodes seen so far.
- Does it work? Yes, it roughly cuts the branching factor from b to \sqrt{b} resulting in double as far look-ahead than pure minimax.

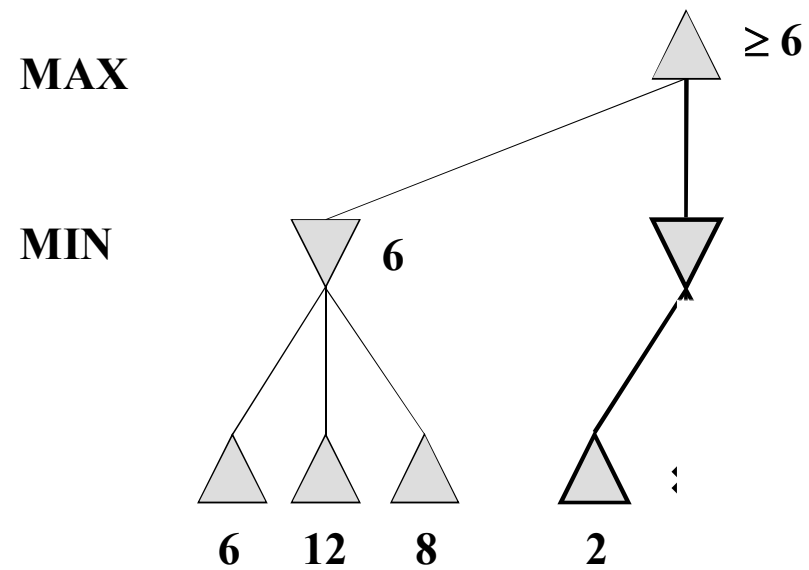
α - β pruning: example

MAX

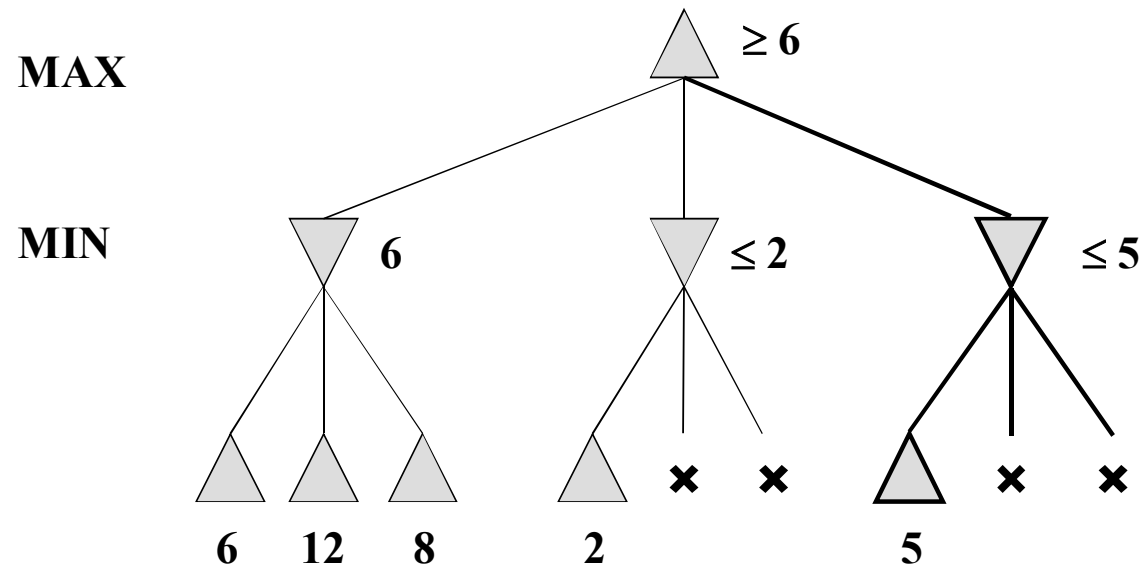
MIN



α - β pruning: example



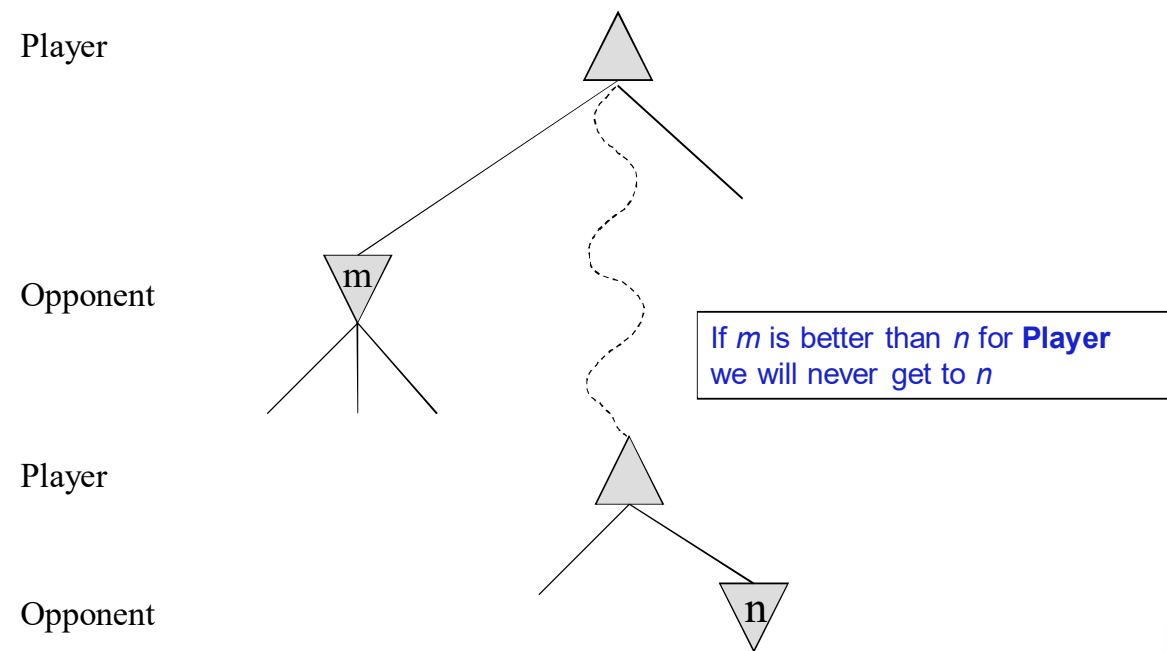
α - β pruning: example



$$\begin{aligned} \text{MINIMAX}(\text{root}) &= \max(\min(6, 12, 8), \min(2, a, b), \min(5, b, d)) \\ &= \max(6, z, y) \quad \text{where } z = \min(2, a, b) \leq 2 \text{ and } y = \min(5, b, d) \leq 5 \\ &= 6 \end{aligned}$$

16

α - β pruning: general principle



More on the α - β algorithm

- Because minimax is depth-first, let's consider nodes along a given path in the tree. Then, as we go along this path, we keep track of:
 - ♦ α : the value of the best (i.e., highest-value) choice we have found so far at any choice point **along the path** for **MAX**
 - ♦ β : the value of the best (i.e., lowest-value) choice we have found so far at any choice point **along the path** for **MIN**

The α - β algorithm:

function ALPHA-BETA-SEARCH(*state*) **returns** an action
 $v \leftarrow \text{MAX-VALUE}(\text{state}, -\infty, +\infty)$
return the *action* in ACTIONS(*state*) with value *v*

function MAX-VALUE(*state*, α , β) **returns** a utility value
if TERMINAL-TEST(*state*) **then return** UTILITY(*state*)
 $v \leftarrow -\infty$
for each *a* **in** ACTIONS(*state*) **do**
 $v \leftarrow \text{MAX}(v, \text{MIN-VALUE}(\text{RESULT}(s,a), \alpha, \beta))$
if $v \geq \beta$ **then return** *v*
 $\alpha \leftarrow \text{MAX}(\alpha, v)$
return *v*

function MIN-VALUE(*state*, α , β) **returns** a utility value
if TERMINAL-TEST(*state*) **then return** UTILITY(*state*)
 $v \leftarrow +\infty$
for each *a* **in** ACTIONS(*state*) **do**
 $v \leftarrow \text{MIN}(v, \text{MAX-VALUE}(\text{RESULT}(s,a), \alpha, \beta))$
if $v \leq \alpha$ **then return** *v*
 $\beta \leftarrow \text{MIN}(\beta, v)$
return *v*

More on the α - β algorithm

In Min-Value:

```
v ← +∞  
for each a in ACTIONS(state) do  
  v ← MIN(v, MAX-VALUE(RESULT(s,a), α, β))  
  if v ≤ α then return v  
  β ← MIN(β, v)  
return v
```

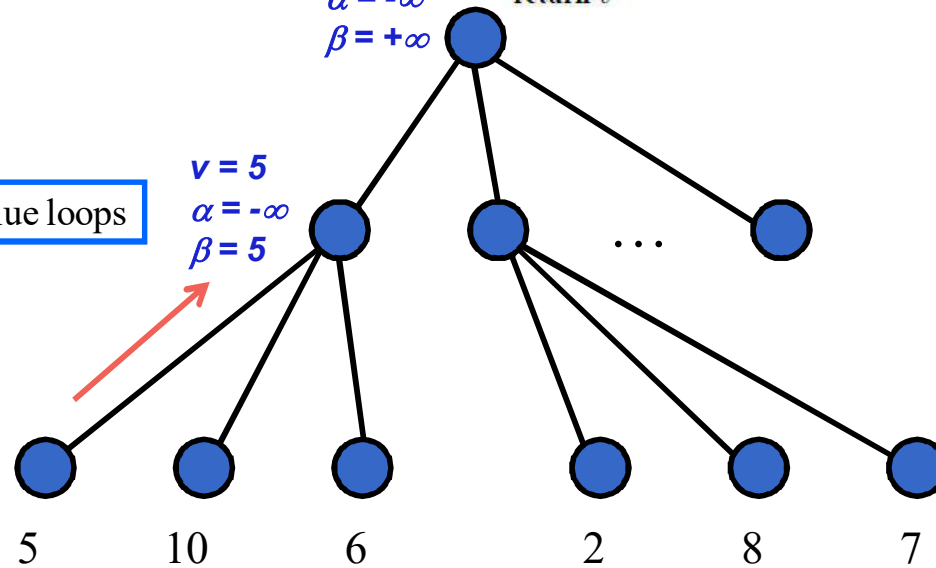
MAX

$v = -\infty$
 $\alpha = -\infty$
 $\beta = +\infty$

MIN Min-Value loops

$v = 5$
 $\alpha = -\infty$
 $\beta = 5$

MAX



More on the α - β algorithm

In Min-Value:

```

v ← +∞
for each a in ACTIONS(state) do
  v ← MIN(v, MAX-VALUE(RESULT(s.a), α, β))
  if v ≤ α then return v
  β ← MIN(β, v)
return v

```

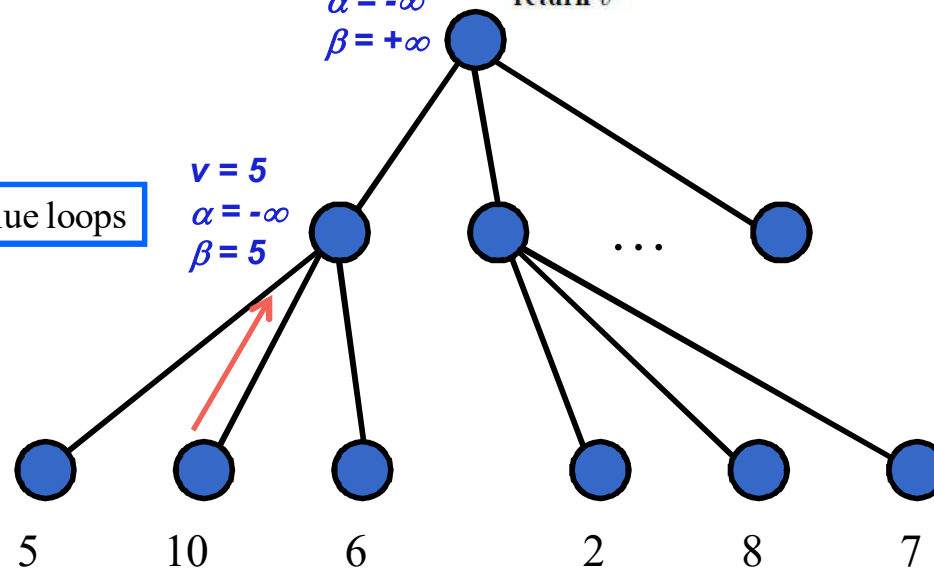
MAX

$v = -\infty$
 $\alpha = -\infty$
 $\beta = +\infty$

MIN Min-Value loops

$v = 5$
 $\alpha = -\infty$
 $\beta = 5$

MAX



More on the α - β algorithm

In Min-Value:

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v ← +∞  
for each a in ACTIONS(state) do  
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  if v ≤ α then return v  
  β ← MIN(β, v)  
return v
```

MAX

$v = -\infty$
 $\alpha = -\infty$
 $\beta = +\infty$

MIN Min-Value loops

$v = 5$
 $\alpha = -\infty$
 $\beta = 5$

MAX



More on the α - β algorithm

In Max-Value:

```

v ← -∞
for each a in ACTIONS(state) do
  v ← MAX(v, MIN-VALUE(RESULT(s,a), α, β))
  if v ≥ β then return v
  α ← MAX(α, v)
return v

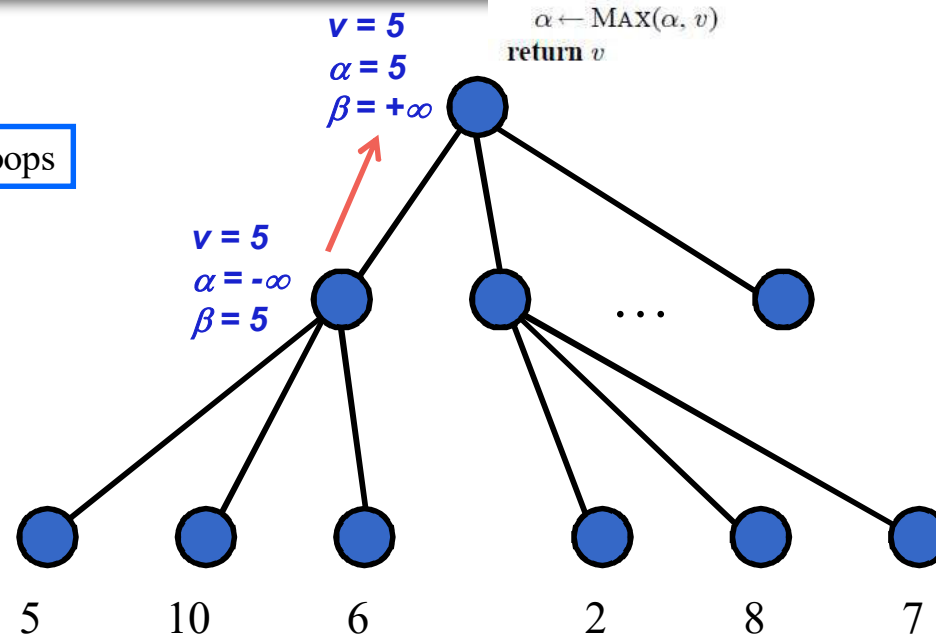
```

MAX

Max-Value loops

MIN

MAX



More on the α - β algorithm

In Max-Value:

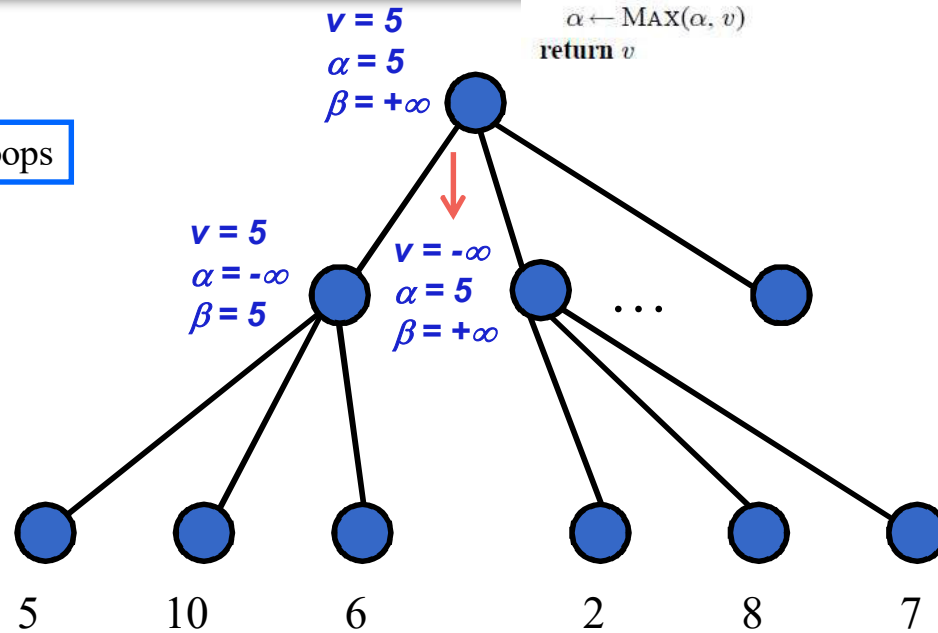
```
v ← -∞  
for each a in ACTIONS(state) do  
  v ← MAX(v, MIN-VALUE(RESULT(s,a), α, β))  
  if v ≥ β then return v  
  α ← MAX(α, v)  
return v
```

MAX

Max-Value loops

MIN

MAX



More on the α - β algorithm

In Min-Value:

```

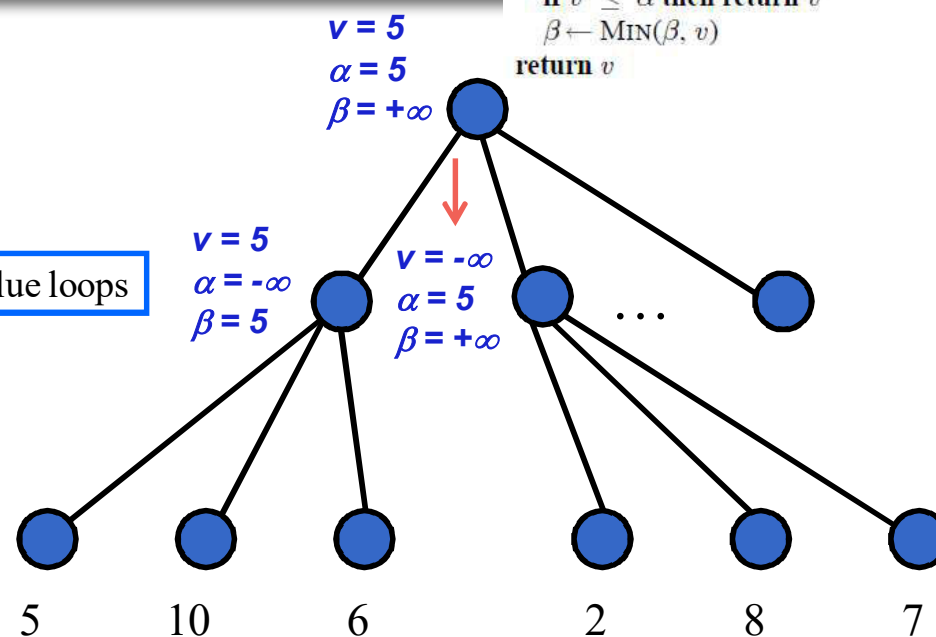
v ← +∞
for each a in ACTIONS(state) do
  v ← MIN(v, MAX-VALUE(RESULT(s.a), α, β))
  if v ≤ α then return v
  β ← MIN(β, v)
return v

```

MAX

MIN Min-Value loops

MAX



More on the α - β algorithm

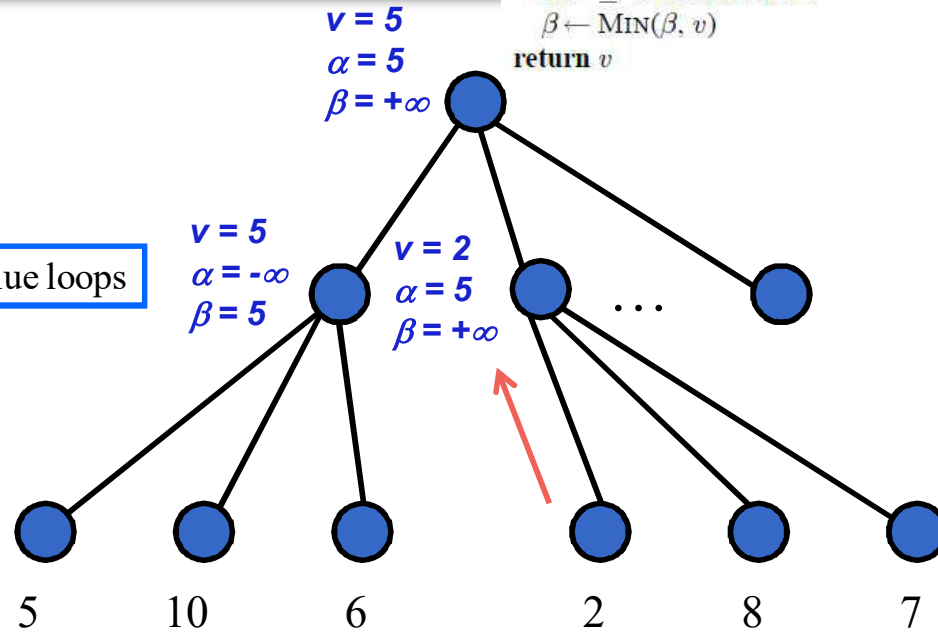
In Min-Value:

```
v ← +∞  
for each a in ACTIONS(state) do  
  v ← MIN(v, MAX-VALUE(RESULT(s,a), α, β))  
  if v ≤ α then return v  
  β ← MIN(β, v)  
return v
```

MAX

MIN Min-Value loops

MAX



More on the α - β algorithm

In Min-Value:

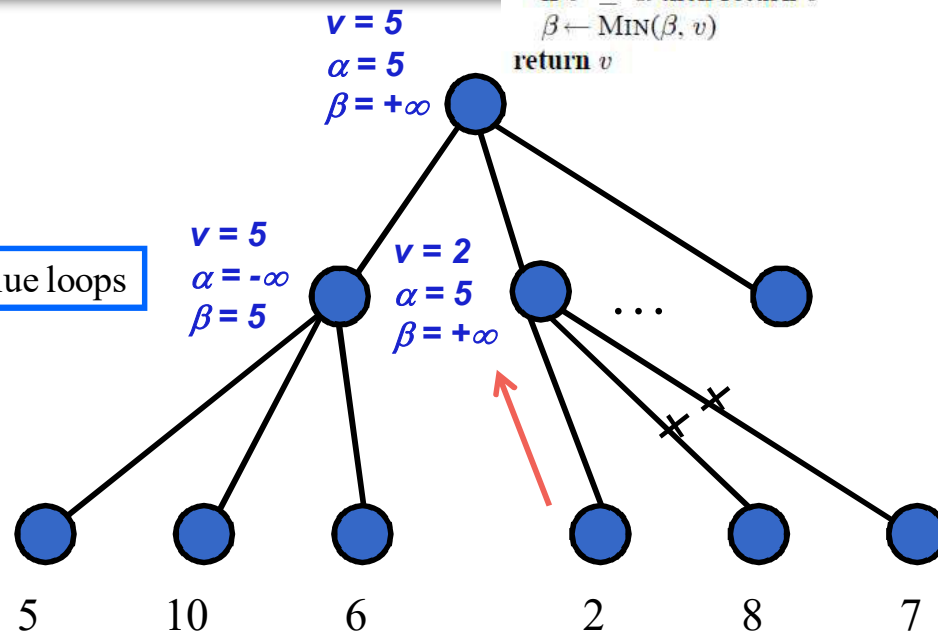
```

v ← +∞
for each a in ACTIONS(state) do
  v ← MIN(v, MAX-VALUE(RESULT(s.a), α, β))
  if v ≤ α then return v
  β ← MIN(β, v)
return v
    
```

MAX

MIN Min-Value loops

MAX



More on the α - β algorithm

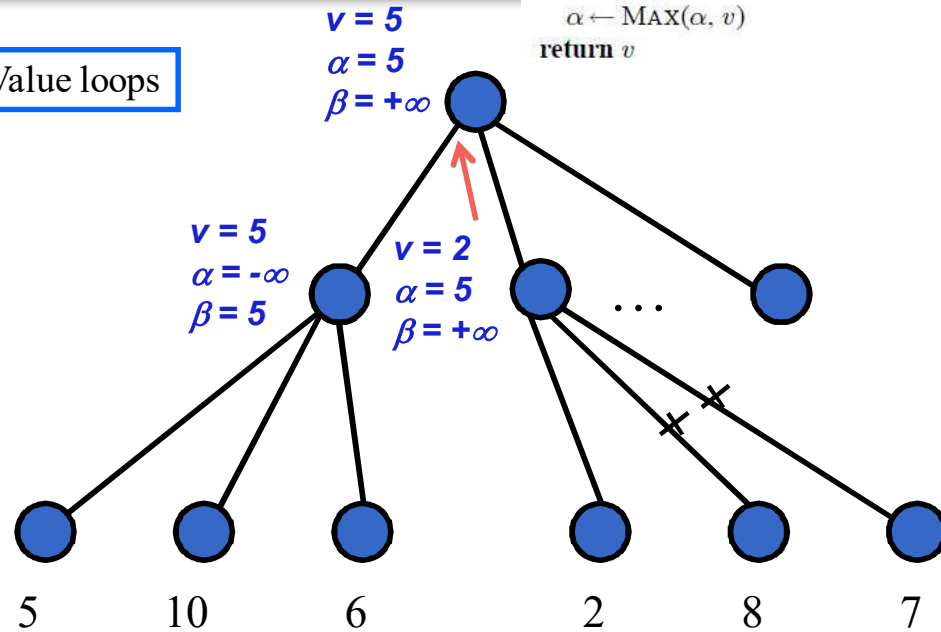
In Max-Value:

```
v ← -∞  
for each a in ACTIONS(state) do  
  v ← MAX(v, MIN-VALUE(RESULT(s,a), α, β))  
  if v ≥ β then return v  
  α ← MAX(α, v)  
return v
```

MAX Max-Value loops

MIN

MAX



Properties of α - β

- Pruning does not affect the final result!!!
- Good move ordering improves effectiveness of pruning
- With *perfect ordering*, time complexity = $O(b^{m/2})$
 - ♦ doubles depth of search
 - ♦ need a heuristic how to order
 - ♦ can easily reach depth 8 => good chess
- A simple example of the value of reasoning about which computations are relevant (a form of *metareasoning*)

2. Move evaluation without complete search

- The minimax algorithm generates the entire game search space, whereas the alpha-beta algorithm allows us to prune large parts of it.
- Complete search is often too complex and impractical. alpha-beta is still DFS.
- **Evaluation function:** evaluates value of state using **heuristics** and cuts off search
- **New MINIMAX:**
 - ♦ **CUTOFF-TEST:** cutoff test to replace the termination condition (e.g., deadline, depth-limit, etc.)
 - ♦ **EVAL:** evaluation function to replace utility function (e.g., number of chess pieces taken)

30

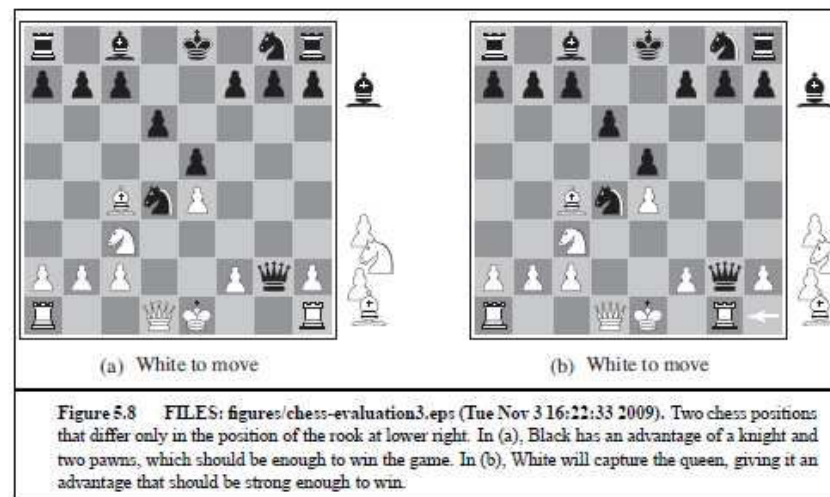
Evaluation function

- The evaluation function should order the *terminal* states in the same way as the true utility function ($a < b < c \dots$).
- The computation must not take too long!
Significant compared to minimax?
- For nonterminal states, the evaluation function should be strongly correlated with the actual chances of winning.

Evaluation functions

- Most calculate features – e.g., number of pawns
- From that we can form categories, equivalence classes.
- Any category represent states that win, lose or result in draws.
- If we know 72% lead to win (+1), 20% to loss (0), 8% drawn (1/2).
Expected value:
 - $(0,72 * +1) + (0,20 * 0) + (0,08 * 1/2) = 0,76$

Evaluation functions

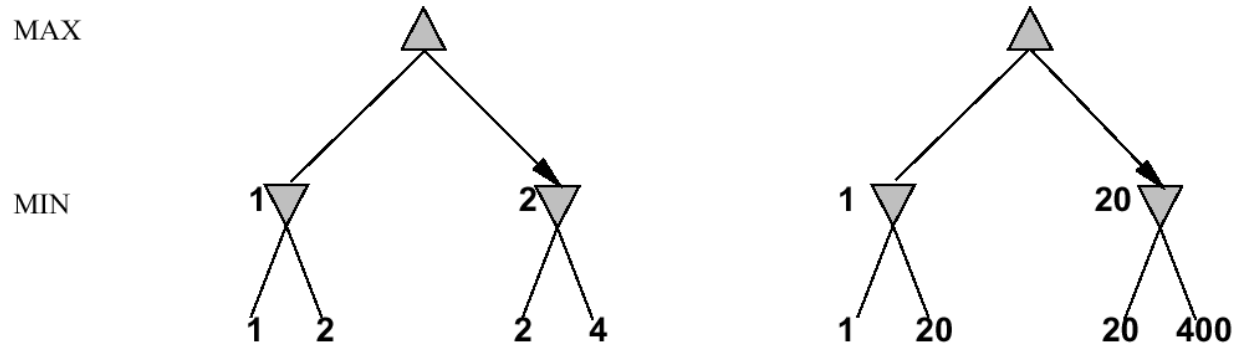


- **Weighted linear evaluation function:** to combine n heuristics

$$f = w_1 f_1 + w_2 f_2 + \dots + w_n f_n$$

E.g, w 's could be the values of pieces (1 for pawn, 3 for bishop etc.)
 f 's could be the number of type of pieces on the board

Note: exact values do not matter



Behaviour is preserved under any *monotonic* transformation of EVAL

Only the order matters:

payoff in deterministic games acts as an *ordinal utility* function

With cutoff and eval

function MAX-VALUE(*state*, α , β) *returns a utility value*

inputs: *state*, current state in game

α , the value of the best alternative for MAX along the path to *state*

β , the value of the best alternative for MIN along the path to *state*

if CUTOFF-TEST(*state*, *depth*) **then return** EVAL(*state*)

$v \leftarrow -\infty$

for *a*, *s* in SUCCESSORS(*state*) **do**

$v \leftarrow \text{MAX}(v, \text{MIN-VALUE}(s, \alpha, \beta))$

if $v \geq \beta$ **then return** *v*

$\alpha \leftarrow \text{MAX}(\alpha, v)$

return *v*

Minimax with cutoff: viable algorithm?

MINIMAXCUTOFF is identical to MINIMAXVALUE except

1. TERMINAL? is replaced by CUTOFF?
2. UTILITY is replaced by EVAL

Does it work in practice?

$$b^m = 10^6, \quad b = 35 \quad \Rightarrow \quad m = 4$$

4-ply lookahead is a hopeless chess player!

4-ply \approx human novice

8-ply \approx typical PC, human master

12-ply \approx Deep Blue, Kasparov

Assume we have
100 seconds,
evaluate 10^4
nodes/s; can
evaluate 10^6
nodes/move

Other Cutoff methods

- Quiescent search
apply eval only to positions that are quiescent, have no big change of value in the near future.
- Forward pruning
considers not all moves in a concrete position.
Beam search is one approach to forward pruning.
- ProbCut
probabilistic alpha-beta with statistical prior knowledge

Games of chance

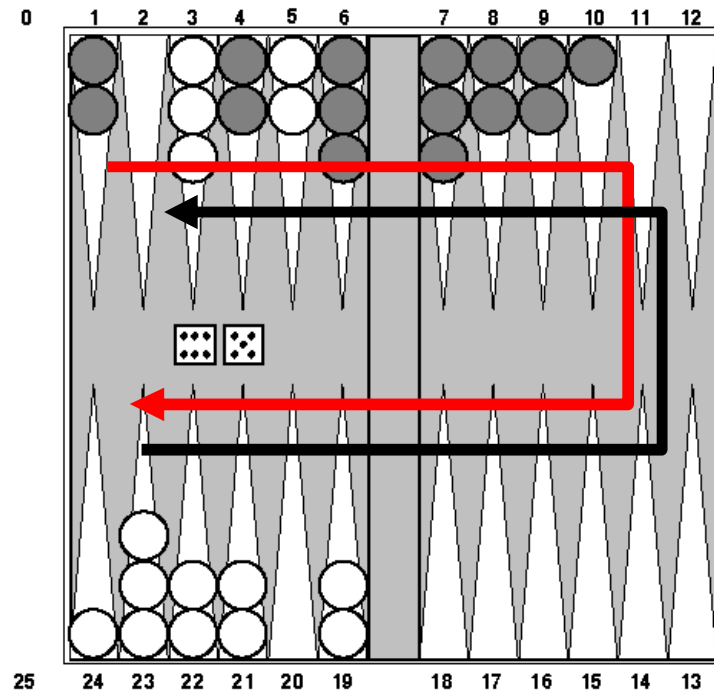
- *Backgammon is a two-player game with **uncertainty**.*

- *Players roll dices to determine what moves to make.*

- *White/red arrow has just rolled 5 and 6 and has four legal moves:*

- 5-10, 5-11
- 5-11, 19-24
- 5-10, 10-16
- 5-11, 11-16

- *Such games are good for exploring decision making in adversarial problems involving skill and luck.*



Game Trees with Chance Nodes

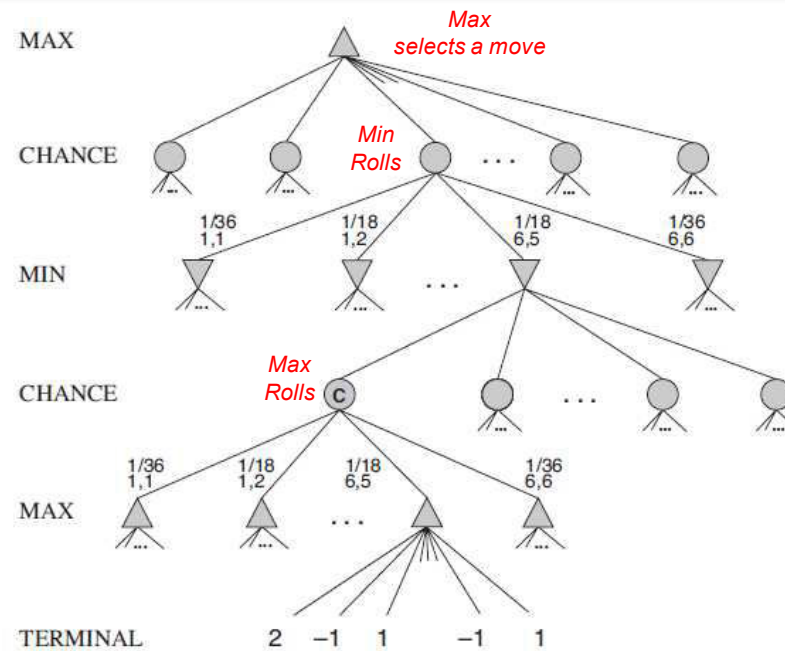
- Use minimax to compute values for MAX and MIN nodes
- Use **expected values** for chance nodes

- For chance nodes over a max node, as in C:

$$\text{expectimax}(C) = \sum_i (P(d_i) * \text{maxvalue}(i))$$

- For chance nodes over a min node:

$$\text{expectimin}(N) = \sum_i (P(d_i) * \text{minvalue}(i))$$



Algorithm for nondeterministic games

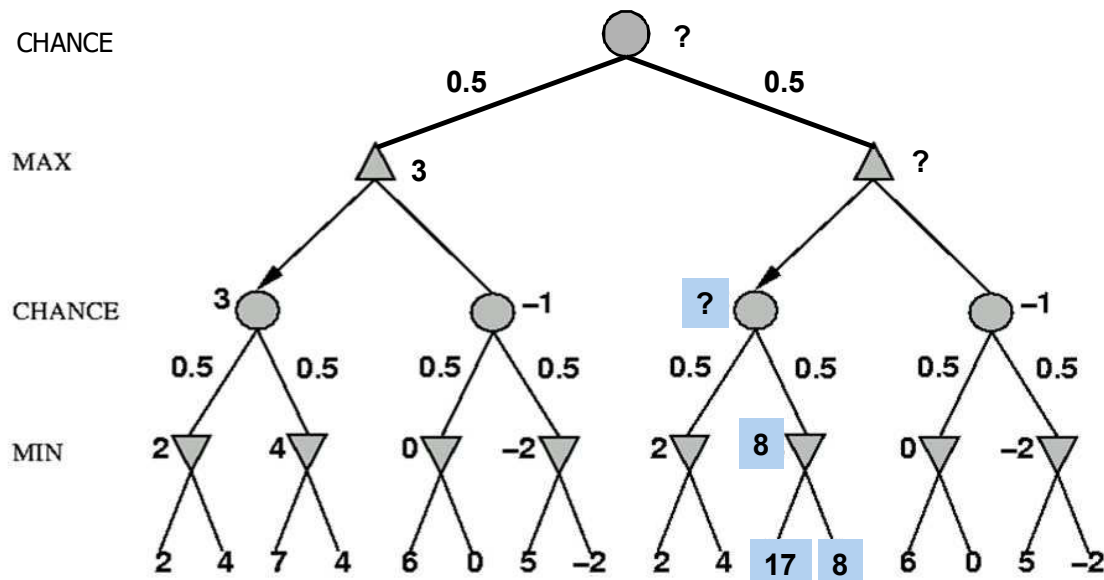
EXPECTIMINIMAX gives perfect play.

$$\text{EXPECTIMINIMAX}(s) = \begin{cases} \text{UTILITY}(s) & \text{if } \text{TERMINAL-TEST}(s) \\ \max_a \text{EXPECTIMINIMAX}(\text{RESULT}(s, a)) & \text{if } \text{PLAYER}(s) = \text{MAX} \\ \min_a \text{EXPECTIMINIMAX}(\text{RESULT}(s, a)) & \text{if } \text{PLAYER}(s) = \text{MIN} \\ \sum_r P(r) \text{EXPECTIMINIMAX}(\text{RESULT}(s, r)) & \text{if } \text{PLAYER}(s) = \text{CHANCE} \end{cases}$$

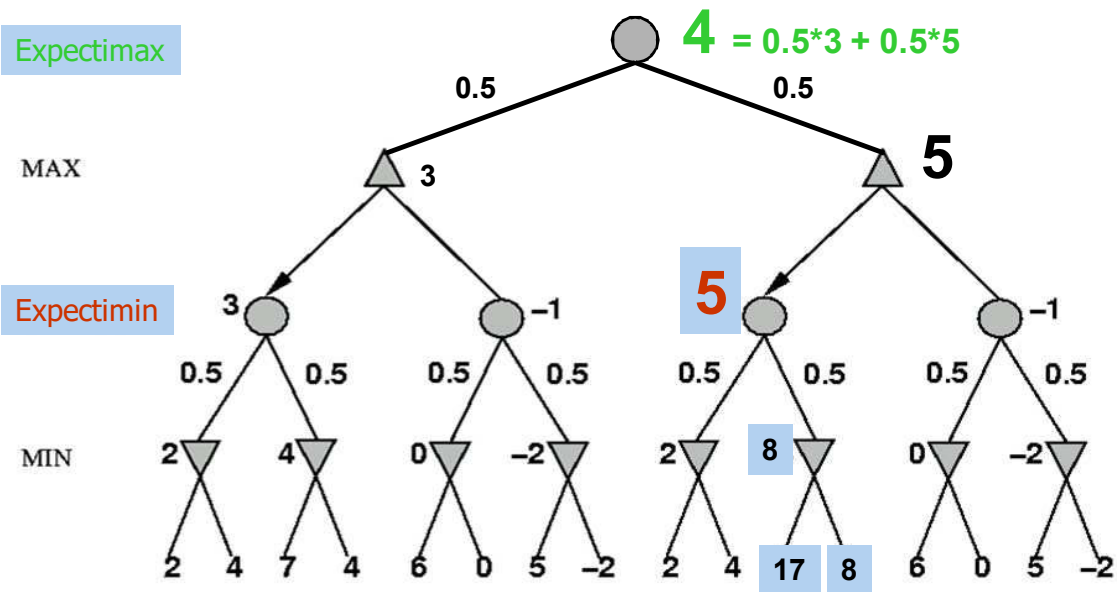
A version of α - β is possible but only if leaf values are bounded. WHY??

Nondeterministic games: the element of chance

expectimax and **expectimin**, expected values over all possible outcomes



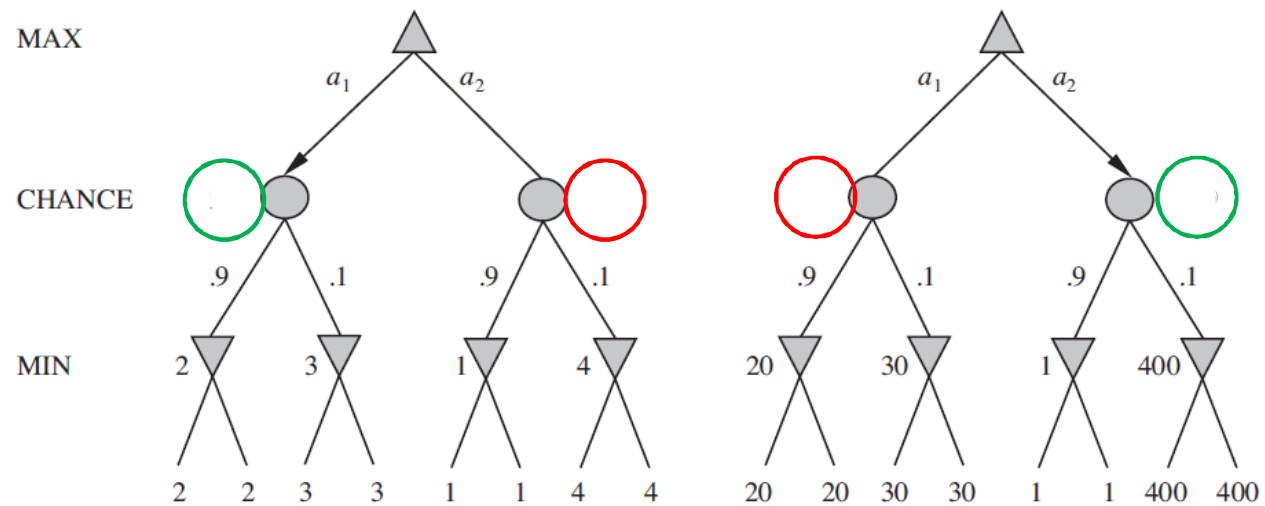
Nondeterministic games: the element of chance



42

Evaluation functions

Order-preserving transformation do not necessarily behave the same!



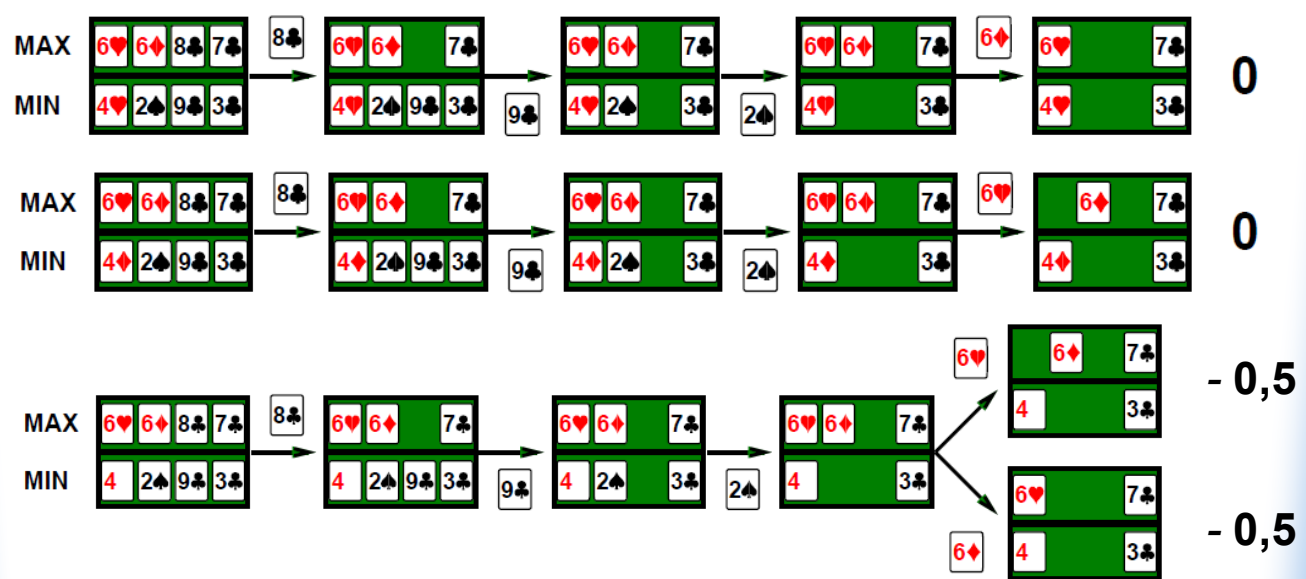
Games of imperfect information

- E.g., card games, where opponent's initial cards are unknown
- Typically we can calculate a probability for each possible deal
- Seems just like having one big dice roll at the beginning of the game
- Idea: compute the minimax value of each action in each deal, then choose the action with highest expected value over all deals
- Special case: if an action is optimal for all deals, it's optimal.

- GIB, current best bridge program, approximates this idea by
 - ♦ generating 100 deals consistent with bidding information
 - ♦ picking the action that wins most tricks on average

Example

- Four card bridge, MAX to play first



Proper analysis

- Intuition that the value of an action is the average of its values in all actual states is *WRONG*
- With partial observability, value of an action depends on the *information state* or *belief state* the agent is in
- Can generate and search a tree of information states
- Leads to rational behaviors such as
 - ◆ Acting to obtain information
 - ◆ Signalling to one's partner
 - ◆ Acting randomly to minimize information disclosure

Summary

- Games are fun to work on!
- They illustrate several important points about AI
 - ♦ perfection is unattainable → must approximate
 - ♦ good idea to think about what to think about
 - ♦ uncertainty constrains the assignment of values to states
 - ♦ optimal decisions depend on information state, not real state

**Intelligent Autonomous Agents
and Cognitive Robotics
Topic 3: Constraint Satisfaction
Problems**

Slides partly from Hwee Tou Ng's
Chapter 5 of AIMA

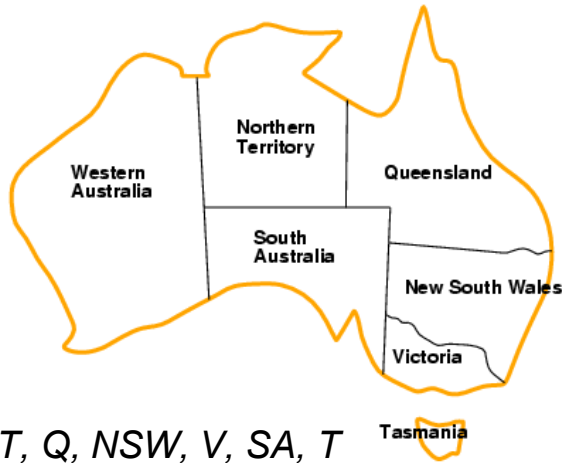
Outline

- Constraint Satisfaction Problems (CSP)
- Backtracking search for CSPs
- Multi-Agents → distributed backtracking

Constraint satisfaction problems (CSPs)

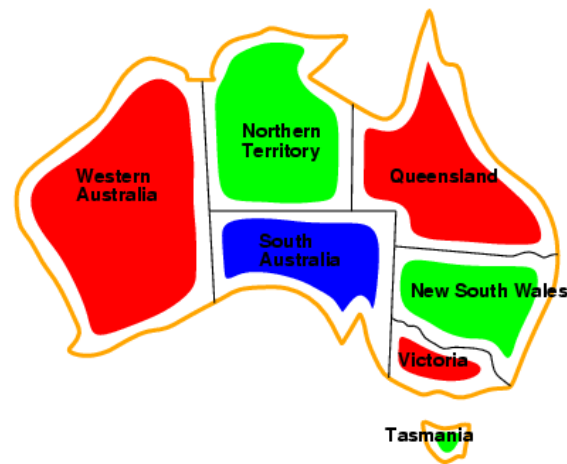
- Standard search problem:
 - ♦ **state** is a "black box" – any data structure that supports successor function, heuristic function, and goal test
- CSP:
 - ♦ **state** is defined by **variables** X_i ($i=1..n$) with
 - ♦ **values** from **domain** D_i
 - ♦ **goal test** is a set of **constraints** C_m ($m=1..z$) specifying allowable combinations of values for subsets of variables
- Simple example of a **formal representation language**
- Allows useful **general-purpose** algorithms with more power than standard search algorithms

Visual example: Map-Coloring



- **Variables:** WA, NT, Q, NSW, V, SA, T
- **Domains:** $\forall i, D_i = \{\text{red, green, blue}\}$
- **Constraints:** adjacent regions must have different colors
 - ♦ e.g., $WA \neq NT$
 - ♦ or $(WA, NT) \in \{(\text{red, green}), (\text{red, blue}), (\text{green, red}), (\text{green, blue}), (\text{blue, red}), (\text{blue, green})\}$

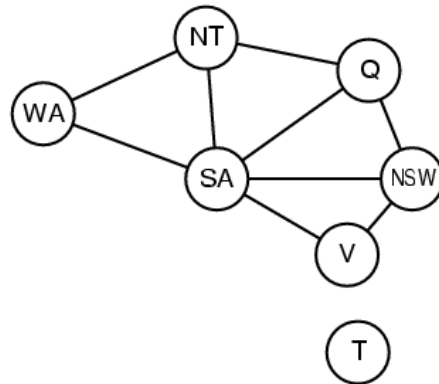
Example: Map-Coloring



- Solutions are complete and consistent assignments, e.g.,
 - ♦ WA = red, NT = green, Q = red, NSW = green, V = red, SA = blue, T = green

Constraint graph

- **Binary CSP:** each constraint relates two variables
- **Constraint graph:** nodes are variables, arcs are constraints



General-purpose CSP algorithms use the graph structure to speed up search. E.g., Tasmania is an independent sub problem!

Varieties of CSPs

- Discrete variables
 - ♦ finite domains:
 - n variables, domain size $d \rightarrow O(d^n)$ complete assignments
 - e.g., n-queens problem
 - ♦ infinite domains:
 - integers, strings, etc.
 - e.g., job scheduling, variables are start/end days for each job
 - need a constraint language, e.g., $StartJob_1 + 5 \leq StartJob_3$
- Continuous variables
 - ♦ e.g., start/end times for Hubble Space Telescope observations
 - ♦ linear constraints solvable in polynomial time by linear programming

Varieties of constraints

- **Unary** constraints involve a single variable,
 - ◆ e.g., $SA \neq \text{green}$
- **Binary** constraints involve pairs of variables,
 - ◆ e.g., $SA \neq WA$
- **Higher-order** constraints involve 3 or more variables

Real-world CSPs

- Assignment problems
 - ◆ e.g., who teaches what class
- Timetabling problems
 - ◆ e.g., which class is offered when and where; preferences
- Hardware configuration
- Transportation scheduling
- Factory scheduling

Constraint propagation

- In CSP an algorithm can do
 - ◆ Constraint propagation = inference
 - ◆ Search
 - ◆ Intertwined or as preprocessing
- The key idea is to create *local* consistency

Node consistency

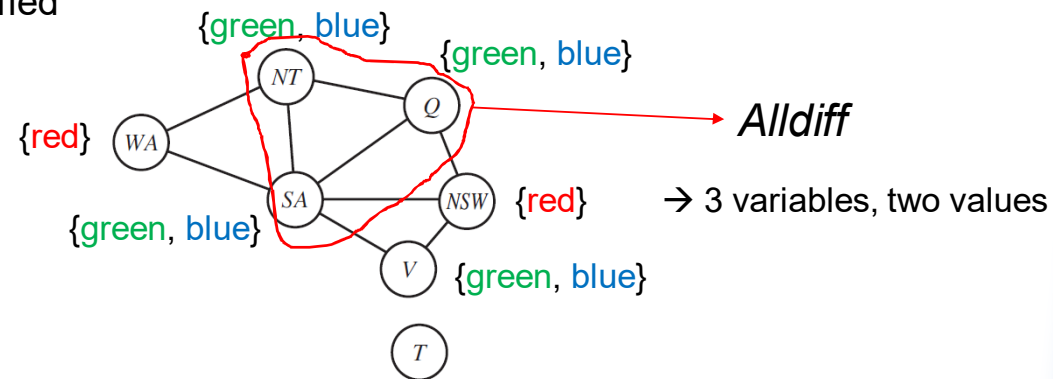
- A variable is node-consistent if all the values satisfy the unary constraints
- Infer the values that are legal for a variable,
 - ♦ e.g. if South Australia does not like green, eliminate it {red, blue}
 - ♦ e.g. don't want to teach at 8 pm

Global Constraints

- *Alldiff* (many algorithms)
 - ♦ Idea: If m variables have n values and $m > n \rightarrow$ can not be satisfied
 - Remove any variable with singleton domain and propagate this into other domains. Repeat as long as there are singleton domains.
 - If an empty domain is produced or $m > n$, then an inconsistency has been detected

Global Constraints

- *Alldiff* (many algorithms)
 - ♦ Idea: If m variables have n values and $m > n \rightarrow$ can not be satisfied

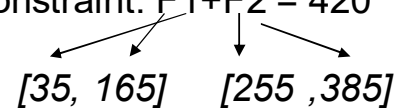


Resource Constraints

- Resource constraints: *Atmost*
We can detect an inconsistency simply by checking the sum of the minimum values of the current domains:
Atmost(10, P1, P2, P3, P4) persons for tasks.
 - ♦ Each variable has domain {3, 4, 5, 6}
→ can not be satisfied
 - ♦ Each variable has domain {2, 3, 4, 5, 6}
→ delete 5 and 6

Resource Constraints

- Bounds propagation/bounds consistent
 - ♦ In complex problems often not possible to enumerate domain values
 - ♦ Constraints:
 - Plane capacities for $F1=[0, 165]$, $F2[0, 385]$
 - Constraint: $F1+F2 = 420$



Resource Constraints

- Bounds propagation/bounds consistent
 - ♦ In complex problems often not possible to enumerate domain values
 - ♦ Constraints:
 - Plane capacities for $F1=[0, 165]$, $F2[0, 385]$
 - Constraint: $F1+F2 = 420$
 $\rightarrow F1[35, 165]$ and $F2[255, 385]$
 - ♦ We say that a CSP is **bounds consistent** if for every variable X , and for both the lower-bound and upper-bound values of X , there exists some value of Y that satisfies the constraint between X and Y for every variable Y . (Often used in praxis)

Standard search formulation

Let's start with the straightforward approach, then fix it

States are defined by the values assigned so far

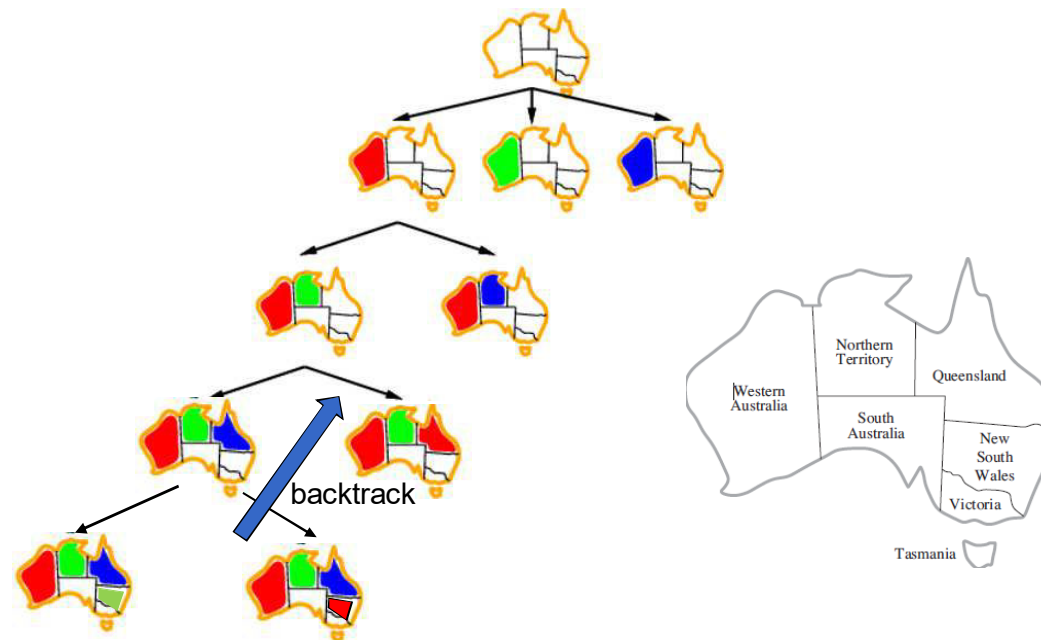
- **Initial state**: the empty assignment { }
- **Successor function**: assign a value to an unassigned variable that does not conflict with current assignment
→ fail if no legal assignments
- **Goal test**: the current assignment is complete

1. Every solution appears at depth n with n variables
→ use depth-first search
2. Path is irrelevant
3. At the root we have n variables and d values $b = nd$
4. At depth l we have $b = (n - l)d$
5. All combinations $n! \cdot d^n$ leaves

Backtracking search

- Variable assignments are **commutative**
[WA = red then NT = green] same as [NT = green then WA = red]
- Only need to consider assignments to a single variable at each node
→ $b = d$ branching factor, n variables → d^n leaves
- Depth-first search for CSPs with single-variable assignments is called **backtracking** search
- Backtracking search is the basic uninformed algorithm for CSPs

Backtracking example



Backtracking search

```
function BACKTRACKING-SEARCH(csp) returns a solution, or failure  
  return BACKTRACK({ }, csp)
```

```
function BACKTRACK(assignment, csp) returns a solution, or failure  
  if assignment is complete then return assignment  
  var ← SELECT-UNASSIGNED-VARIABLE(csp)  
  for each value in ORDER-DOMAIN-VALUES(var, assignment, csp) do  
    if value is consistent with assignment then  
      add {var = value} to assignment  
      inferences ← INFERENCE(csp, var, value)  
      if inferences ≠ failure then  
        add inferences to assignment  
        result ← BACKTRACK(assignment, csp)  
        if result ≠ failure then  
          return result  
      remove {var = value} and inferences from assignment  
  return failure
```

Improving backtracking efficiency

- **General-purpose** methods can give huge gains in speed:
 - ◆ Which variable should be assigned next
SELECT-UNASSIGNED-VARIABLE?
 - ◆ In what order should its values be tried
ORDER-DOMAIN-VALUES?
 - ◆ What inferences should be performed at each step in
the search INFERENCE?
 - ◆ Can we detect inevitable failure early?

Most constrained variable

- Most constrained variable:
choose the variable with the fewest legal values



- a.k.a. **minimum remaining values (MRV)**
heuristic



Most constrained variable

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choose the variable with the fewest legal values



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heuristic

What about the first state

MRV does not help in the first state

Degree heuristic

- Tie-breaker among most constrained variables:
Degree heuristic
- Most constraining variable:
 - ◆ choose the variable with the most constraints on remaining variables
 - ◆ used together with MRV



Least constraining value

- Given a variable, choose the least constraining value:
 - ♦ the one that rules out the fewest values in the remaining variables *Queensland is selected*



- Combining these heuristics makes 1000 queens feasible

Inference: Forward checking

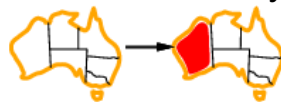
- **Idea:**
 - ♦ Keep track of remaining legal values for unassigned neighbors
 - ♦ Terminate search when any variable has no legal values



WA = red

Inference: Forward checking

- **Idea:**
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Inference: Forward checking

- **Idea:**
 - ◆ Keep track of remaining legal values for unassigned neighbors
 - ◆ Terminate search when any variable has no legal values



WA	NT	Q	NSW	V	SA	T
Red, Green, Blue	Red, Green, Blue	Red, Green, Blue	Red, Green, Blue	Red, Green, Blue	Red, Green, Blue	Red, Green, Blue
Red	Green, Blue	Red, Green, Blue	Red, Green, Blue	Red, Green, Blue	Green, Blue	Red, Green, Blue
Red	Blue	Green, Blue	Red, Blue	Red, Green, Blue	Blue	Red, Green, Blue



Q = green

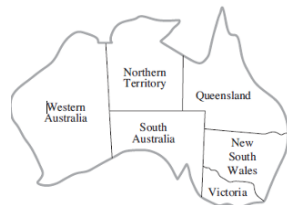
Victoria = blue

Inference: Forward checking

- **Idea:**
 - ◆ Keep track of remaining legal values for unassigned neighbors
 - ◆ Terminate search when any variable has no legal values



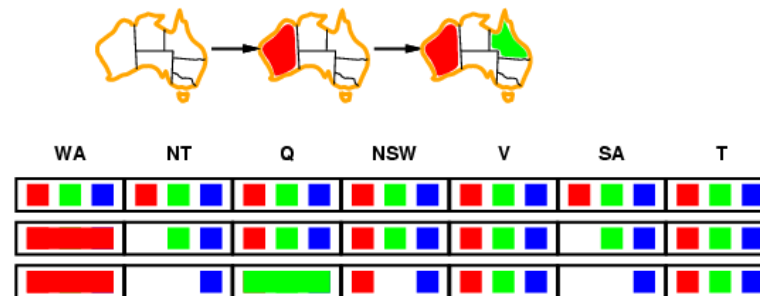
WA	NT	Q	NSW	V	SA	T
Red	Green	Blue	Red	Green	Blue	Red
Red	Green	Blue	Red	Green	Blue	Red
Red	Green	Blue	Red	Green	Blue	Red
Red	Green	Blue	Red	Green	Blue	Red



Victoria = blue

Forward checking

- Forward checking propagates information from assigned to unassigned variables, but doesn't provide early detection for all failures:



- NT and SA cannot both be blue!
- **Constraint propagation repeatedly** enforces constraints locally, to neighbors

Arc consistency

- Simplest form of propagation makes each arc **consistent**
 $X \rightarrow Y$ is consistent iff
 for **every** value x of X there is **some** allowed y of Y
- Constraint $Y=X^2$ and domain $\{0,1,..9\}$. Can write the constraint as
 $[(X, Y), \{(0, 0), (1, 1), (2, 4), (3, 9)\}]$
 Can reduce the domains
 $X = \{0, 1, 2, 3\}$
 $Y = \{0, 1, 4, 9\}$
- What about $(SA \neq WA)$ and domain $\{\text{red, green, blue}\}$
 $[(SA, WA), \{(\text{red}, \text{green}), (\text{red}, \text{blue}), (\text{green}, \text{red}), (\text{green}, \text{blue}), (\text{blue}, \text{red}), (\text{blue}, \text{green})\}]$



Arc consistency algorithm AC-3

```
function AC-3(csp) returns the CSP, possibly with reduced domains
  inputs: csp, a binary CSP with variables  $\{X_1, X_2, \dots, X_n\}$ 
  local variables: queue, a queue of arcs, initially all the arcs in csp

  while queue is not empty do
     $(X_i, X_j) \leftarrow \text{REMOVE-FIRST}(\textit{queue})$ 
    if REMOVE-INCONSISTENT-VALUES( $X_i, X_j$ ) then
      for each  $X_k$  in NEIGHBORS[ $X_i$ ] do
        add  $(X_k, X_i)$  to queue



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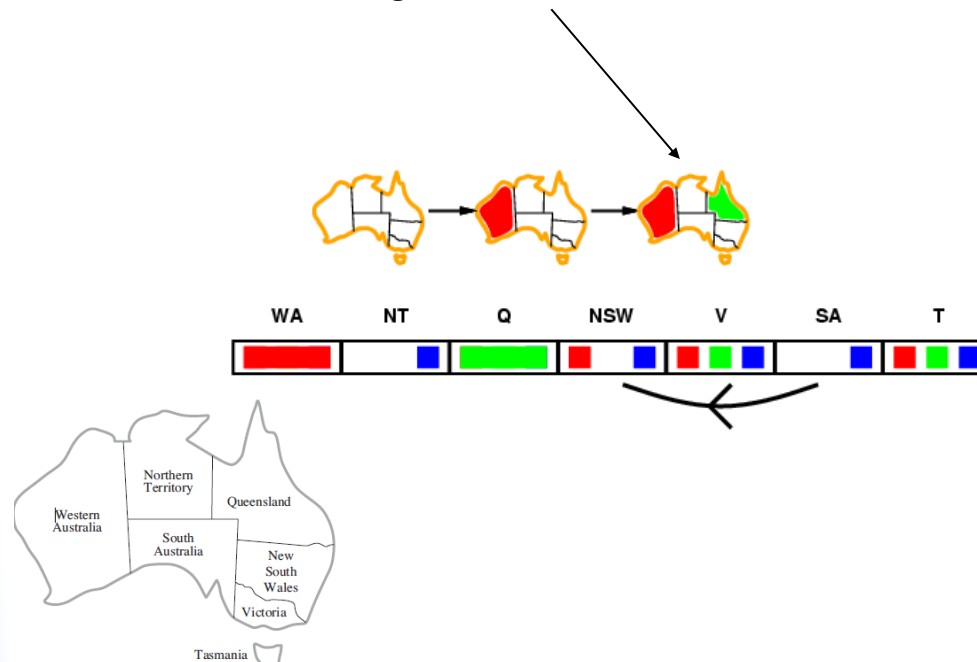

function REMOVE-INCONSISTENT-VALUES( $X_i, X_j$ ) returns true iff succeeds
  removed  $\leftarrow$  false
  for each  $x$  in DOMAIN[ $X_i$ ] do
    if no value  $y$  in DOMAIN[ $X_j$ ] allows  $(x, y)$  to satisfy the constraint  $X_i \leftrightarrow X_j$ 
      then delete  $x$  from DOMAIN[ $X_i$ ]; removed  $\leftarrow$  true
  return removed
```

- Time complexity: $O(cd^3)$

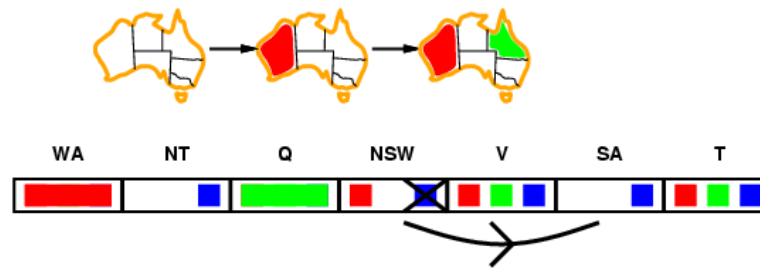
34

Arc consistency

- Assume we begin in state



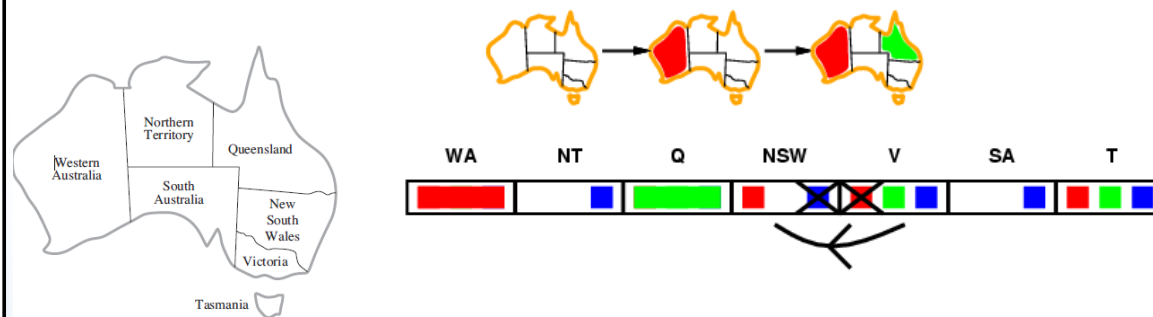
Arc consistency



If X loses a value, neighbors of X need to be rechecked

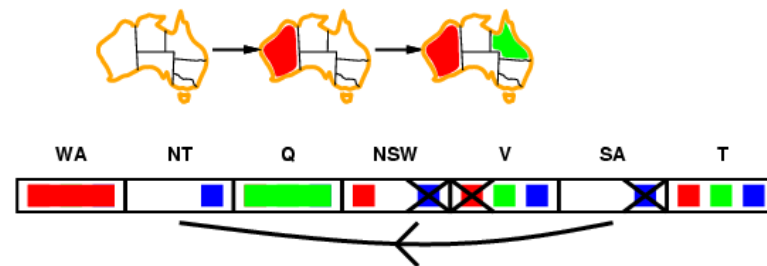


Arc consistency



- If X loses a value, neighbors of X need to be rechecked

Arc consistency



- If X loses a value, neighbors of X need to be rechecked
- Is run as a preprocessor
- Can also be modified to work with backtracking
 - ♦ On assignment put only (X_i, X_j) in the queue

Path consistency

- $\{X_i, X_j\}$ is path consistent with respect to X_m if for every consistent assignment there is an for X_m that is consistent.
 $\{X_i, X_m\}$ and $\{X_m, X_j\}$.
See the CSP graph for detecting paths
- Could also be extended to K-Consistency

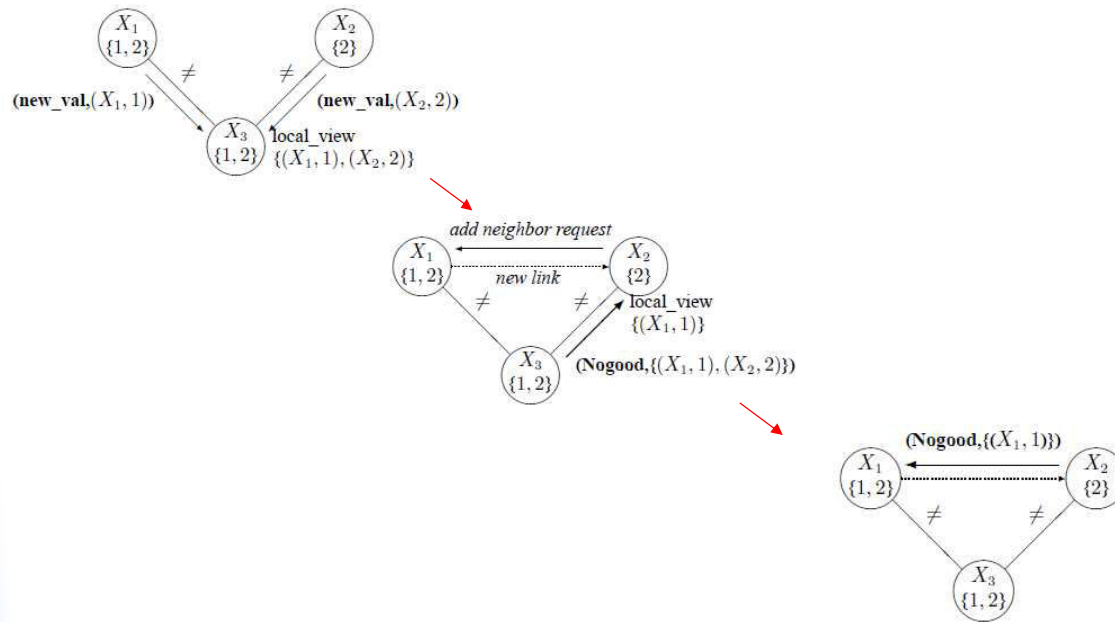
Multi-Agents CSP

- Also called distributed CSP
 - ♦ Variable and domain definition as before
 - ♦ Each agent owns a variable (many can be mapped to one)
 - ♦ Agents decides on value with relative autonomy
 - ♦ Has no global view on all dependencies
 - ♦ BUT! Can communicate with his neighbors in the constraint graph
- Many algorithms!! We only sketch one important algorithm

Multi-Agents CSP: Asynchronous Backtracking

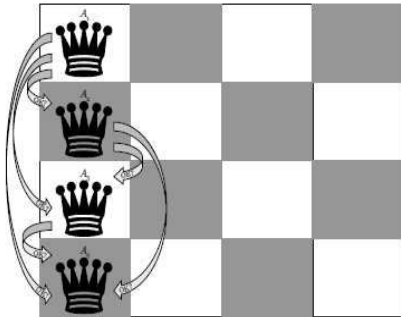
- The algorithm makes an ordering on agents and assigns them priority numbers. All agents set their initial value concurrently
- a higher-priority agent j informs all lower-priority agents k_i of its assignment if connected in constraint graph
- lower-priority agent k evaluates the shared C_{jk} constraint with its own assignment
 - ♦ if constraints are satisfied with the current assignment \rightarrow no action
 - ♦ otherwise, agent k looks for a different value consistent with choice of agent j
 - ♦ if such a consistent value exists \rightarrow agent j adopts this value and informs other low-priority agents
 - ♦ if such a consistent value does not exist, agent j updates *NoGood list* and sends the message to agent j and seek for a value that is consistent with all connected higher priority agents
 - ♦ j receives a NoGood mentioning i it is not connected with j . j asks i to set up a link

Adding edges



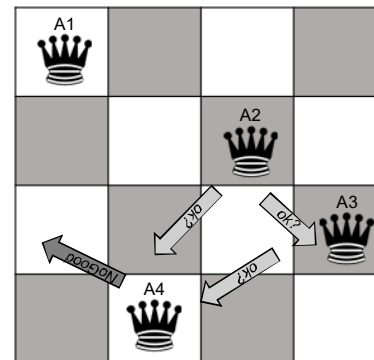
43

Example: 4-Queens



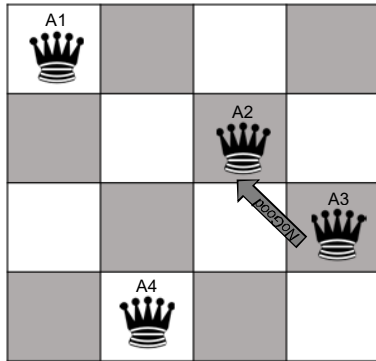
A1 knows no position
 A2 knows A1
 A3 knows A2 and A1
 A4 knows all positions

Based on local information each queen checks where to move or to resolve conflicts with upper queen. Afterwards do nothing, send "OK?" or "NoGood" messages.



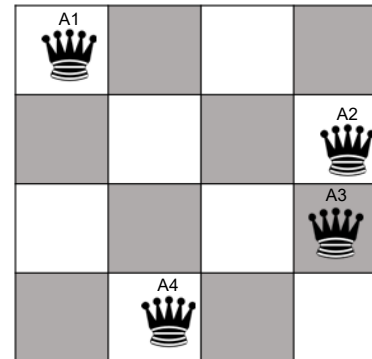
NoGood: $A1=1$ and $A2=1 \rightarrow A3 \neq 1$

Example: 4-Queens

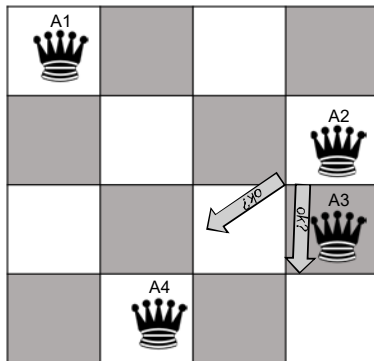


Only A3 is active

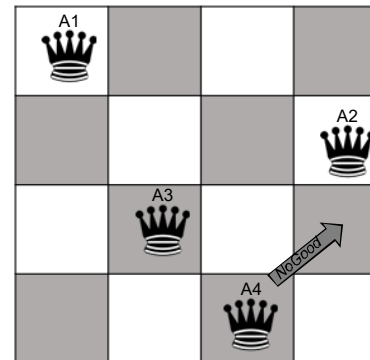
NoGood: $A1=1 \rightarrow A2 \neq 3$



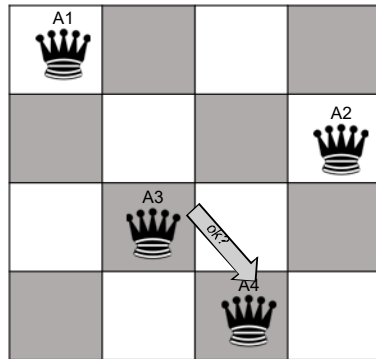
Example: 4-Queens



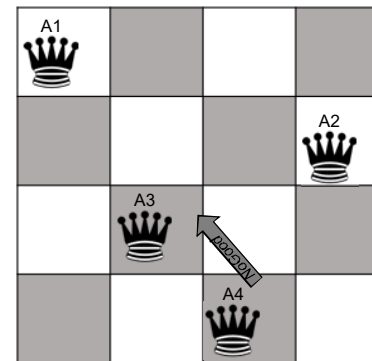
*A4 sends a NoGood message:
 $A1=1$ and $A2=4 \rightarrow A3 \neq 4$ (no
 longer valid)
 and moves.*



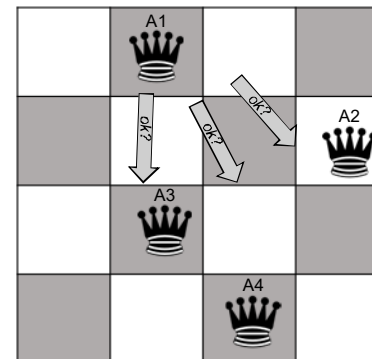
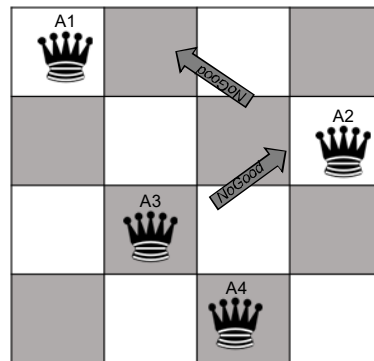
Example: 4-Queens



*A4 sends a NoGood message:
A1=1 and A2=4 \rightarrow A3 \neq 2
and does not move, no conflict.*



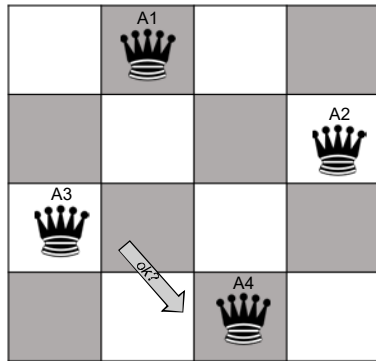
Example: 4-Queens



A3 has no option \rightarrow NoGood: $A1=1 \rightarrow A2 \neq 4$,

A2 had a former NoGood message from A3 not to stay in 3
 \rightarrow send NoGood: $A1 \neq 1$

Example: 4-Queens



No conflict for any queen → solved

**Intelligent Autonomous Agents
and Cognitive Robotics
Topic 5: Bayesian Networks**

Ralf Möller, Rainer Marrone
Hamburg University of Technology

Uncertainty in prior knowledge

- Diagnosis:
 - ◆ $\text{Toothache} \Rightarrow \text{Cavity} \vee \text{GumProblem} \vee \text{Abscess} \vee \dots$

 $\text{RootInfection} \vee \dots \vee \text{Cavity} \Rightarrow \text{Toothache}$
- The connection between toothaches and cavity is just not a logical consequence. For medical diagnosis logic does not seem to be appropriate.

Probability

Probabilistic assertions **summarize** effects of

- **laziness:**
It is too much work to list the complete set of antecedents or consequents needed to ensure an exceptionless rule and too hard to use such rules
- **theoretical ignorance:**
no complete theory, e.g., medical science has no complete theory for the domain.
- **practical ignorance:**
lack of relevant facts, initial conditions, tests, etc.

Making decisions under uncertainty

Suppose I believe the following:

$$P(A_{25} \text{ gets me there on time} \mid \dots) = 0.04$$

$$P(A_{90} \text{ gets me there on time} \mid \dots) = 0.70$$

$$P(A_{120} \text{ gets me there on time} \mid \dots) = 0.95$$

$$P(A_{1440=24h} \text{ gets me there on time} \mid \dots) = 0.9999$$

- Which action to choose?

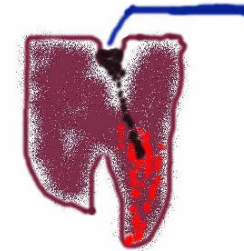
Depends on my **preferences** for missing flight vs. time spent waiting, etc.

- ♦ **Utility theory** is used to represent and use preferences
- ♦ **Decision theory** = **probability theory** + **utility theory**

↓
Later in this lecture

Example world

Example: *Dentist problem* with four variables:
Toothache (I have a toothache)
Cavity (I have a cavity)
Catch (steel probe catches in my tooth)
Weather (*sunny,rainy,cloudy,snow*)



Prior probability

- **Prior or unconditional probabilities** of propositions
e.g., $P(\text{Cavity} = \text{true}) = 0.1$ and $P(\text{Weather} = \text{sunny}) = 0.72$
correspond to belief prior to arrival of any (new) evidence
- **Probability distribution**
gives values for all possible assignments
(*sunny,rainy,cloudy,snow*):
 $P(\text{Weather}) = \langle 0.72, 0.1, 0.08, 0.1 \rangle$
(**normalized**, i.e., sums to 1 because one must be the case)

Full joint probability distribution

- **Joint probability distribution** for a set of random variables gives the probability of every atomic event on those random variables

$P(\text{Weather}, \text{Cavity})$ = a 4×2 matrix of values:

<i>Weather</i> =	sunny	rainy	cloudy	snow	
<i>Cavity</i> = true	0.144	0.02	0.016	0.02	= 0.2
<i>Cavity</i> = false	0.576	0.08	0.064	0.08	= 0.8
					= 1.0

- Full joint probability distribution: all random variables involved
 - ♦ $P(\text{Toothache}, \text{Catch}, \text{Cavity}, \text{Weather})$
- **Every question about a domain can be answered by the full joint distribution**

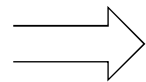
Conditional probability

- **Conditional** or **posterior probabilities** (after having received some information)
e.g., $P(\text{cavity} \mid \text{toothache}) = 0.8$
- Definition of **conditional probability** (in terms of uncond. prob.):
 $P(a \mid b) = P(a \wedge b) / P(b)$ if $P(b) > 0$
- **Product rule** gives an alternative formulation (\wedge is commutative):
 $P(a \wedge b) = P(a \mid b) P(b) = P(b \mid a) P(a)$
- **Chain rule** is derived by successive application of product rule:

$$\begin{aligned} P(X_1, \dots, X_n) &= P(X_1, \dots, X_{n-1}) P(X_n \mid X_1, \dots, X_{n-1}) \\ &= P(X_1, \dots, X_{n-2}) P(X_{n-1} \mid X_1, \dots, X_{n-2}) P(X_n \mid X_1, \dots, X_{n-1}) \\ &= \prod_{i=1}^n P(X_i \mid X_1, \dots, X_{i-1}) \end{aligned}$$

Bayes rule

Product rule: $P(x, y) = P(x|y)P(y) = P(y|x)P(x)$



$$P(\text{cause}|\text{effect}) = \frac{P(\text{effect}|\text{cause}) * P(\text{cause})}{P(\text{effect})} = \frac{\text{likelihood} * \text{prior}}{\text{evidence}}$$

$$P(X|Y) = \alpha P(Y|X)P(X)$$

Inference by enumeration

- Start with the joint probability distribution:

	<i>toothache</i>		\neg <i>toothache</i>	
	<i>catch</i>	\neg <i>catch</i>	<i>catch</i>	\neg <i>catch</i>
<i>cavity</i>	.108	.012	.072	.008
\neg <i>cavity</i>	.016	.064	.144	.576

- For any proposition ϕ , sum the atomic events where it is true: $P(\phi) = \sum_{\omega \in \phi} P(\omega)$

10

Inference by enumeration

- Start with the joint probability distribution:

	<i>toothache</i>		\neg <i>toothache</i>	
	<i>catch</i>	\neg <i>catch</i>	<i>catch</i>	\neg <i>catch</i>
<i>cavity</i>	.108	.012	.072	.008
\neg <i>cavity</i>	.016	.064	.144	.576

$$P(\text{cavity} \vee \text{toothache}) = 0.108 + 0.012 + 0.072 + 0.008 + 0.016 + 0.064 = 0.28$$

Inference by enumeration

- Start with the joint probability distribution:

	<i>toothache</i>		\neg <i>toothache</i>	
	<i>catch</i>	\neg <i>catch</i>	<i>catch</i>	\neg <i>catch</i>
<i>cavity</i>	.108	.012	.072	.008
\neg <i>cavity</i>	.016	.064	.144	.576

- Can also compute conditional probabilities:

$$\begin{aligned} P(\neg \text{cavity} \mid \text{toothache}) &= \frac{P(\neg \text{cavity} \wedge \text{toothache})}{P(\text{toothache})} \\ &\stackrel{\text{Product rule}}{=} \frac{0.016 + 0.064}{0.108 + 0.012 + 0.016 + 0.064} \\ &= 0.4 \end{aligned}$$

Normalization

	<i>toothache</i>		\neg <i>toothache</i>	
	<i>catch</i>	\neg <i>catch</i>	<i>catch</i>	\neg <i>catch</i>
<i>cavity</i>	.108	.012	.072	.008
\neg <i>cavity</i>	.016	.064	.144	.576

- Denominator $P(\mathbf{z})$ (or $P(\text{toothache})$ in the example before) can be viewed as a **normalization constant** α

$$\begin{aligned}
 P(\text{Cavity} \mid \text{toothache}) &= P(\text{Cavity}, \text{toothache}) / P(\text{toothache}) \\
 &= \alpha P(\text{Cavity}, \text{toothache}) \\
 &= \alpha [P(\text{Cavity}, \text{toothache}, \text{catch}) + P(\text{Cavity}, \text{toothache}, \neg \text{catch})] \\
 &= \alpha [0.108 + 0.016 + 0.012 + 0.064] \\
 &= \alpha [0.2 = 0.12 + 0.08] = \alpha [0.6 + 0.4]
 \end{aligned}$$

$$\begin{aligned}
 \alpha * (0.12 + 0.08) &= 1 \\
 \alpha * 0.2 &= 1 \\
 \alpha &= 1 / 0.2 = 5 \\
 5 * 0.12 &= 0.6 \\
 5 * 0.08 &= 0.4
 \end{aligned}$$

General idea: compute distribution on query variable by fixing **evidence variables** (toothache) and summing over **hidden variables** (Catch)

General inference procedure

Typically, we are interested in
the posterior joint distribution of the **query variables** Y
given specific values e for the **evidence variables** E
 X are all variables of the modeled world

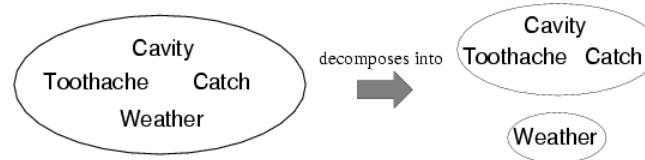
Let the **hidden variables** be $H = X - Y - E$ then the required summation of joint entries is done by summing out the hidden variables:

$$P(Y | E = e) = \alpha P(Y, E = e) = \alpha \sum_h P(Y, E = e, H = h)$$

- The terms in the summation are joint entries because Y , E and H together exhaust the set of random variables (X)
- Obvious problems:
 1. Space complexity $O(d^n)$ to store the joint distribution where d is the largest arity and n denotes the number of random variables
 2. Worst-case time complexity $O(d^n)$
 3. How to find the numbers for $O(d^n)$ entries?

Independence

- A and B are independent iff
 $P(A|B) = P(A)$ or $P(B|A) = P(B)$ or $P(A, B) = P(A) P(B)$



$$P(\textit{Toothache}, \textit{Catch}, \textit{Cavity}, \textit{Weather}) \\ = P(\textit{Toothache}, \textit{Catch}, \textit{Cavity}) P(\textit{Weather})$$

- 32 entries table can be constructed from 8 and 4 entries;
- Absolute independence powerful but rare
- How can we check whether we have independent variables in the full joint?

Example #1

Bread	Bagels	Butter	$p(r,a,u)$
0	0	0	0.24
0	0	1	0.06
0	1	0	0.12
0	1	1	0.08
1	0	0	0.12
1	0	1	0.18
1	1	0	0.04
1	1	1	0.16

Bread	$p(r)$
0	
1	

$$P(a,u)=P(a)P(u)?$$

$$P(r,a)=P(r)P(a)?$$

16

Example #1

Butter	p(u)
0	0.52
1	0.48

Bread	Bagels	Butter	p(r,a,u)
0	0	0	0.24
0	0	1	0.06
0	1	0	0.12
0	1	1	0.08
1	0	0	0.12
1	0	1	0.18
1	1	0	0.04
1	1	1	0.16

Bagels	p(a)
0	0.6
1	0.4

Bread	p(r)
0	0.5
1	0.5

Bagels	Butter	p(a,u)
0	0	
0	1	
1	0	
1	1	

?

$$P(a,u) = P(a)P(u)?$$

$$P(r,a) = P(r)P(a)?$$

Example #1

Butter	$p(u)$
0	0.52
1	0.48

Bagels	$p(a)$
0	0.6
1	0.4

Bread	$p(r)$
0	0.5
1	0.5

Bread	Bagels	Butter	$p(r,a,u)$
0	0	0	0.24
0	0	1	0.06
0	1	0	0.12
0	1	1	0.08
1	0	0	0.12
1	0	1	0.18
1	1	0	0.04
1	1	1	0.16

Bagels	Butter	$p(a,u)$
0	0	0.36
0	1	0.24
1	0	0.16
1	1	0.24

$\neq 0.52 \cdot 0.6 = 0.312$

Bread	Bagels	$p(r,a)$
0	0	0.3
0	1	0.2
1	0	0.3
1	1	0.2

$P(a,u) = P(a)P(u)$? **NO**

$P(r,a) = P(r)P(a)$? **YES**

18

Conditional independence

- $P(\textit{Toothache}, \textit{Cavity}, \textit{Catch})$ has $2^3 - 1 = 7$ independent entries
- If I have a cavity, the probability that the probe catches in doesn't depend on whether I have a toothache:
 - (1) $P(\textit{catch} \mid \textit{toothache}, \textit{cavity}) = P(\textit{catch} \mid \textit{cavity})$
 - (2) $P(\textit{catch} \mid \textit{toothache}, \neg \textit{cavity}) = P(\textit{catch} \mid \neg \textit{cavity})$
- *Catch* is **conditionally independent** of *Toothache* given *Cavity*:
 $P(\textit{Catch} \mid \textit{Toothache}, \textit{Cavity}) = P(\textit{Catch} \mid \textit{Cavity})$
- Equivalent statements:
 $P(\textit{Toothache} \mid \textit{Catch}, \textit{Cavity}) = P(\textit{Toothache} \mid \textit{Cavity})$
 $P(\textit{Toothache}, \textit{Catch} \mid \textit{Cavity}) = P(\textit{Toothache} \mid \textit{Cavity}) P(\textit{Catch} \mid \textit{Cavity})$

Conditional independence contd.

- Write out full joint distribution using chain rule:
$$\begin{aligned} & \mathbf{P}(\textit{Toothache}, \textit{Catch}, \textit{Cavity}) \\ &= \mathbf{P}(\textit{Toothache} \mid \textit{Catch}, \textit{Cavity}) \mathbf{P}(\textit{Catch}, \textit{Cavity}) \\ &= \mathbf{P}(\textit{Toothache} \mid \textit{Catch}, \textit{Cavity}) \mathbf{P}(\textit{Catch} \mid \textit{Cavity}) \mathbf{P}(\textit{Cavity}) \\ & \text{conditional independence} \\ &= \mathbf{P}(\textit{Toothache} \mid \textit{Cavity}) \mathbf{P}(\textit{Catch} \mid \textit{Cavity}) \mathbf{P}(\textit{Cavity}) \end{aligned}$$

i.e., $2 + 2 + 1 = 5$ independent numbers
- In most cases, the use of conditional independence reduces the size of the representation of the joint distribution from exponential in n to linear in n .
- Conditional independence is our most basic and robust form of knowledge about uncertain environments.

Car Example

- Three variables:
 - ♦ Gas, Battery, Starts
- $P(\text{Battery}|\text{Gas}) = P(\text{Battery})$
Gas and Battery are independent
- $P(\text{Battery}|\text{Gas}, \text{Starts}) \neq P(\text{Battery}|\text{Starts})$
Gas and Battery are not independent given Starts
- Independence does not imply conditional independence.
- Conditional independence does not imply independence

Question

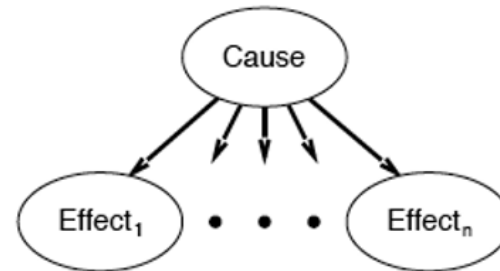
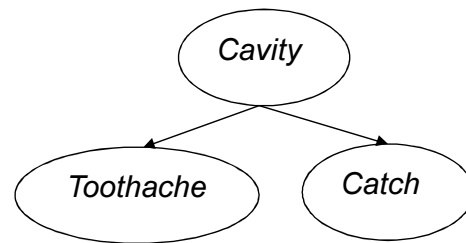
- How can we make use of
 - ◆ independence
 - ◆ and conditional independence

Need a model that can express this

Bayesian networks

- A simple, graphical notation for **conditional independence assertions** and hence for compact specification of the full joint distributions
- Syntax:
 - ♦ a set of nodes, one per variable
 - ♦ a directed, acyclic graph (link \approx "directly influences")
 - ♦ a conditional distribution for each node given its parents:
$$P(X_i | \text{Parents}(X_i))$$
- In the simplest case, conditional distribution represented as a **conditional probability table** (CPT) giving the distribution over X_i for each combination of parent values

Simplest Bayesian Network



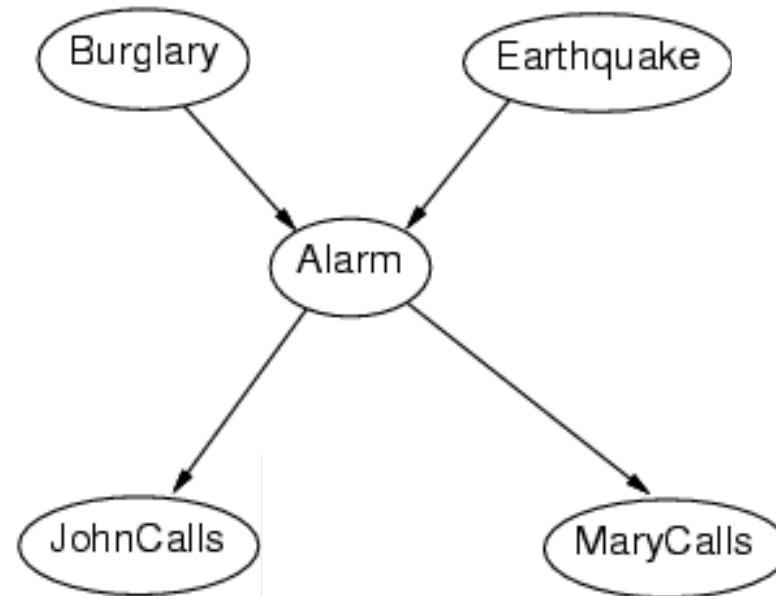
$$P(\text{Cause}|\text{Effect}_1, \text{Effect}_2, \dots) = \alpha P(\text{Cause}) \prod_i P(\text{Effect}_i|\text{Cause})$$

- also called Naïve Bayesian networks
- conditional independence of all effect variables

More complex example

- I'm at work, neighbor John calls to say my alarm is ringing, but neighbor Mary calls but not as often as John. Sometimes it's set off by minor earthquakes but also on burglary. Is there a burglar?
- Variables: *Burglary, Earthquake, Alarm, JohnCalls, MaryCalls*
- Network topology reflects "causal" knowledge:
 - ◆ A burglar can set the alarm off
 - ◆ An earthquake can set the alarm off
 - ◆ The alarm can cause Mary to call
 - ◆ The alarm can cause John to call

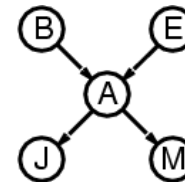
Example contd.



26

Compactness

- A CPT for Boolean X_i with k Boolean parents has 2^k rows for the combinations of parent values
- Each row requires one number p for $X_i = \text{true}$ (the number for $X_i = \text{false}$ is just $1-p$)
- If each of n Boolean variables has no more than k parents, the complete network requires $O(n \cdot 2^k)$ numbers i.e., grows linearly with n , vs. $O(2^n)$ for the full joint distribution
- For burglary net? $1 + 1 + 4 + 2 + 2 = 10$ numbers (vs. $2^5 - 1 = 31$)



k parents with **n** values each and **m** values for the child node of the parents?

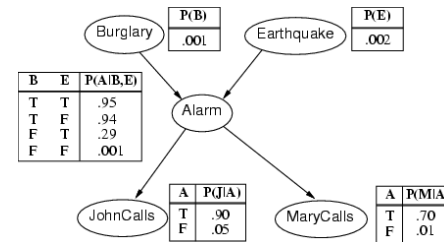
Number of independent values = $n^k \cdot (m-1)$

Semantics

The full joint distribution can be rewritten using the *chain rule*:

$$P(X_1, \dots, X_n) = \prod_{i=1}^n P(X_i | X_1, \dots, X_{i-1})$$

$$P(X_1, \dots, X_n) = \prod_{i=1}^n P(X_i | \text{parent}(X_i))$$



Assumption: Independence and Conditional independence assertions are correctly modeled

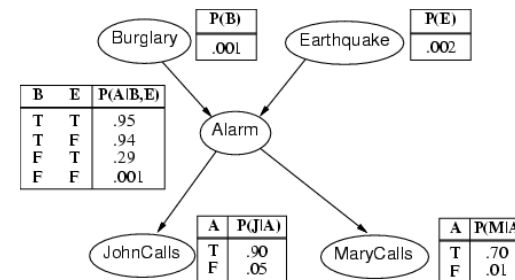
Semantics

The full joint distribution is defined as the product of the local conditional distributions:

$$P(X_1, \dots, X_n) = \prod_{i=1}^n P(X_i | \text{parent}(X_i))$$

e.g., $P(j \wedge m \wedge a \wedge \neg b \wedge \neg e)$

$$\begin{aligned} &= P(j | a) P(m | a) P(a | \neg b, \neg e) P(\neg b) P(\neg e) \\ &= 0.90 \times 0.7 \times 0.001 \times 0.999 \times 0.998 \\ &\approx 0.00063 \end{aligned}$$

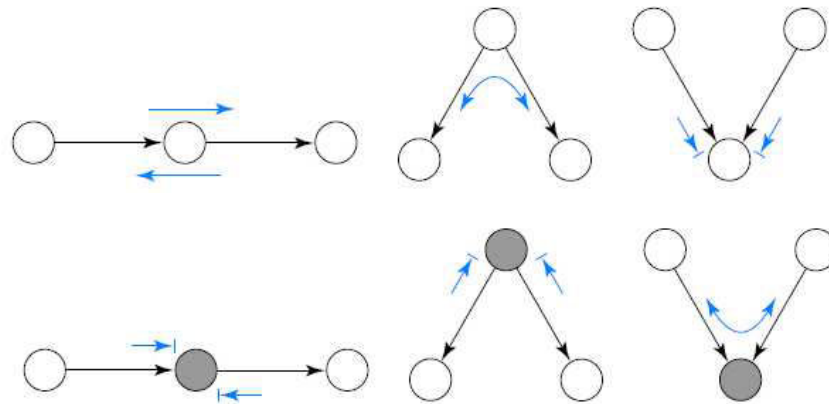


Encoding conditional independence via d-separation

- We can determine if conditional independence holds by a graph separation criterion called **d-separation** (*direction dependent separation*)
- X and Y are **d-separated** if there is no active path between them.
- The formal definition of **active** is somewhat involved. The Bayes Ball Algorithm gives a nice graphical definition.

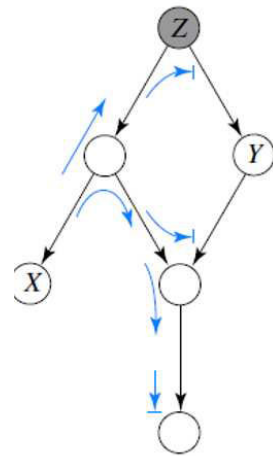
The six rules of Bayes Ball

An undirected path is active if a Bayes ball travelling along it never encounters the “stop” symbol: $\rightarrow \perp$



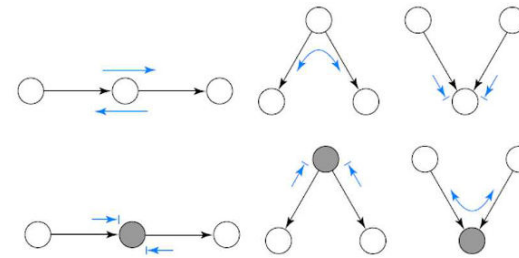
If there are no active paths from X to Y when $\{Z_1, \dots, Z_k\}$ are shaded, then $X \perp\!\!\!\perp Y \mid \{Z_1, \dots, Z_k\}$.

A double-header: two games of Bayes Ball



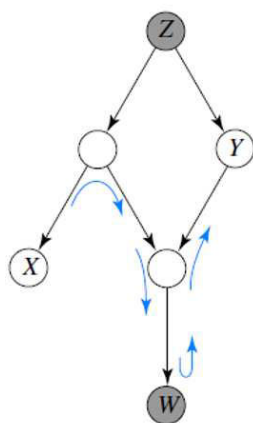
no active paths

$$X \perp\!\!\!\perp Y \mid Z$$

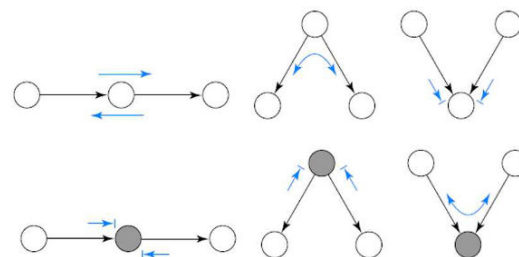


32

A double-header: two games of Bayes Ball

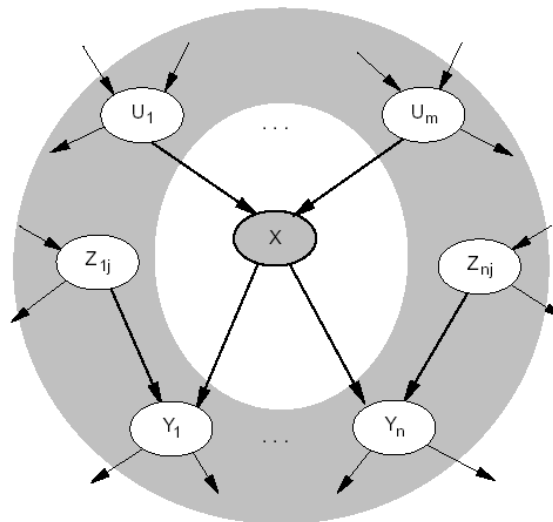


one active path
 $X \not\perp Y \mid \{W, Z\}$



Markov Blanket

- Markov blanket: Parents + children + children's parents
- Node is conditionally independent of all other nodes in network, given its Markov Blanket -> simplifies computation -> gather information on the nodes of the Markov Blanket?



34

Constructing Bayesian networks

- 1. Choose an ordering of variables X_1, \dots, X_n .
Cause should precede effects.
- 2. For $i = 1$ to n
 - ♦ add X_i to the network
 - ♦ select parents from X_1, \dots, X_{i-1} such that

$$P(X_i | \text{Parents}(X_i)) = P(X_i | X_1, \dots, X_{i-1})$$

This choice of parents guarantees:

$$\begin{aligned}
 P(X_1, \dots, X_n) &= \prod_{i=1}^n P(X_i | X_1, \dots, X_{i-1}) \\
 \text{(chain rule)} & \\
 &= \prod_{i=1}^n P(X_i | \text{Parents}(X_i)) \\
 \text{(by construction)} &
 \end{aligned}$$

Example

- Suppose we choose the ordering M, J, A, B, E

$P(J | M) = P(J)$? **No**

MaryCalls

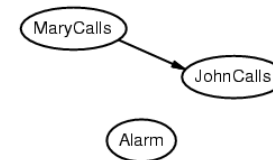
JohnCalls

Example

- Suppose we choose the ordering M, J, A, B, E

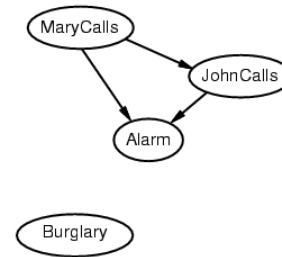
$P(A \mid J, M) = P(A)$? **No**

$P(A \mid J, M) = P(A \mid J)$? **No**



Example

- Suppose we choose the ordering M, J, A, B, E



$P(B | A, J, M) = P(B)$? **No**

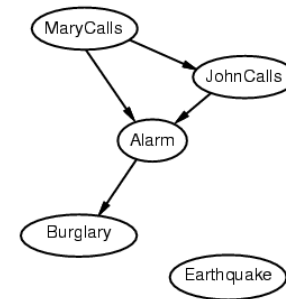
$P(B | A, J, M) = P(B | A)$? **Yes**

Example

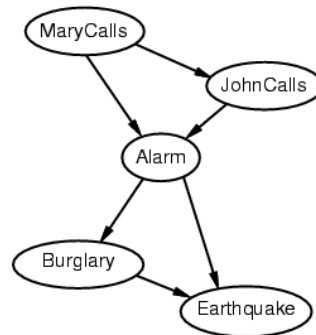
- Suppose we choose the ordering M, J, A, B, E

$P(E \mid B, A, J, M) = P(E \mid A)$? **No**

$P(E \mid B, A, J, M) = P(E \mid A, B)$? **Yes**



Example contd.

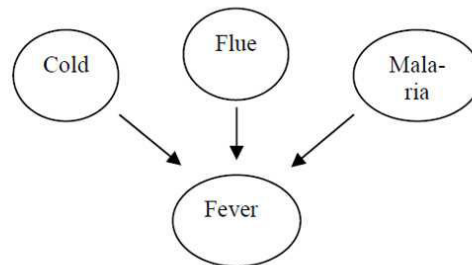


- Deciding conditional independence is hard in noncausal directions
- (Causal models and conditional independence seem hardwired for humans!)
- Network is less compact: $1 + 2 + 4 + 2 + 4 = 13$ numbers needed instead of 10.

Efficient implementation of CPTs

- The number of independent entries grow exponentially with the number of parents.
- Two ways to overcome this
 - ♦ Restrict the number of parents if possible
 - ♦ Instead of free distributions, often canonical (parameterized) distributions are suggested. One popular example of such a pattern is the noisy OR for discrete cases.

Example



The noisy OR is a generalization of the logical OR. Three assumptions:

1. All possible causes U_i for a event X are listed (you can add a *leak* node)
2. Negated causes $\neg U_i$ do not have any influence on X
3. Independent failure probability q_i for each cause alone.

$$q_{\text{cold}} = P(\neg \text{fever} \mid \text{cold}, \neg \text{flu}, \neg \text{malaria}) = 0.6 ,$$

$$q_{\text{flu}} = P(\neg \text{fever} \mid \neg \text{cold}, \text{flu}, \neg \text{malaria}) = 0.2 ,$$

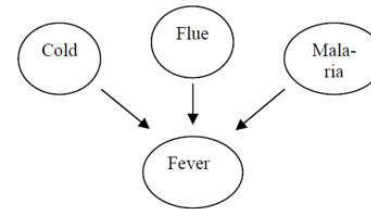
$$q_{\text{malaria}} = P(\neg \text{fever} \mid \neg \text{cold}, \neg \text{flu}, \text{malaria}) = 0.1 .$$

Example

$$q_{\text{cold}} = P(\neg \text{fever} \mid \text{cold}, \neg \text{flu}, \neg \text{malaria}) = 0.6,$$

$$q_{\text{flu}} = P(\neg \text{fever} \mid \neg \text{cold}, \text{flu}, \neg \text{malaria}) = 0.2,$$

$$q_{\text{malaria}} = P(\neg \text{fever} \mid \neg \text{cold}, \neg \text{flu}, \text{malaria}) = 0.1.$$



$$P(\neg x \mid o_1, o_2, \dots, o_r, \neg o_{r+1}, \dots, \neg o_n) = \prod_{r=1}^r q_i$$

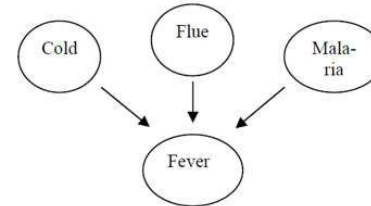
<i>Cold</i>	<i>Flu</i>	<i>Malaria</i>	$P(\text{Fever})$	$P(\neg \text{Fever})$
F	F	F		
F	F	T		0.1
F	T	F		0.2
F	T	T		
T	F	F		0.6
T	F	T		
T	T	F		
T	T	T		

Example

$$q_{\text{cold}} = P(\neg \text{fever} \mid \text{cold}, \neg \text{flu}, \neg \text{malaria}) = 0.6 ,$$

$$q_{\text{flu}} = P(\neg \text{fever} \mid \neg \text{cold}, \text{flu}, \neg \text{malaria}) = 0.2 ,$$

$$q_{\text{malaria}} = P(\neg \text{fever} \mid \neg \text{cold}, \neg \text{flu}, \text{malaria}) = 0.1 .$$

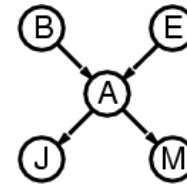


$$P(x \mid o_1, o_2, \dots, o_r, \neg o_{r+1}, \dots, \neg o_n) = 1 - \prod_{i=1}^r q_i$$

<i>Cold</i>	<i>Flu</i>	<i>Malaria</i>	$P(\text{Fever})$	$P(\neg \text{Fever})$
F	F	F	0.0	1.0
F	F	T	0.9	0.1
F	T	F	0.8	0.2
F	T	T	0.98	$0.02 = 0.2 \times 0.1$
T	F	F	0.4	0.6
T	F	T	0.94	$0.06 = 0.6 \times 0.1$
T	T	F	0.88	$0.12 = 0.6 \times 0.2$
T	T	T	0.988	$0.012 = 0.6 \times 0.2 \times 0.1$

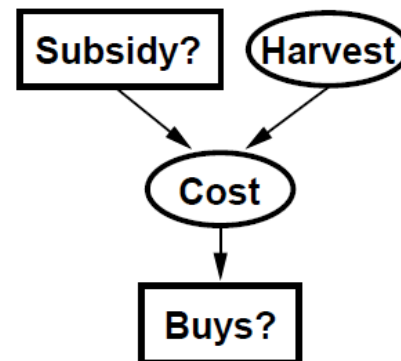
Last Time

- Structure and semantic of BN
 - ◆ Modelling of independence and conditional independence
 - ◆ Causal and non-causal networks
 - ◆ d-separation, Markov blanket
 - ◆ Efficient CPTs, e.g., noisy OR, trees, Min, Max, ...



Hybrid (discrete+contionous) networks

Discrete (*Subsidy?* and *Buys?*); continuous (*Harvest* and *Cost*)



Option 1: discretization—possibly large errors, large CPTs

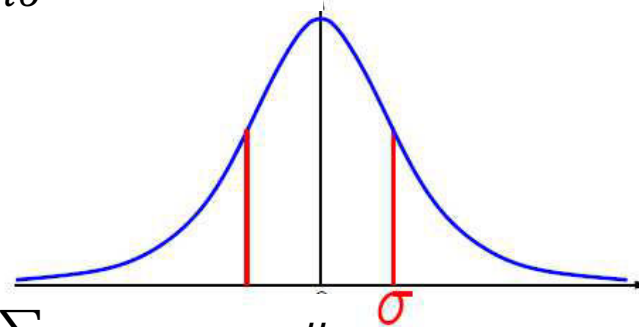
Option 2: finitely parameterized canonical families

1) Continuous variable, discrete+continuous parents (e.g., *Cost*)

2) Discrete variable, continuous parents (e.g., *Buys?*)

Continuous variables: Gaussian density

$$P(x) = \frac{1}{\sqrt{2\pi}\sigma} e^{-(x-\mu)^2 / 2\sigma^2}$$



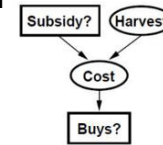
Mean $\mu = \frac{1}{n} \sum_n x_i$

Variance $\sigma^2 = \frac{1}{n-1} \sum_n (x_i - \mu)^2$

Standard deviation $\sigma = \sqrt{\sigma^2}$

Continuous child variables

- Need one *conditional density* function for child variable given
 - ♦ continuous parents
 - ♦ for each discrete value of parents



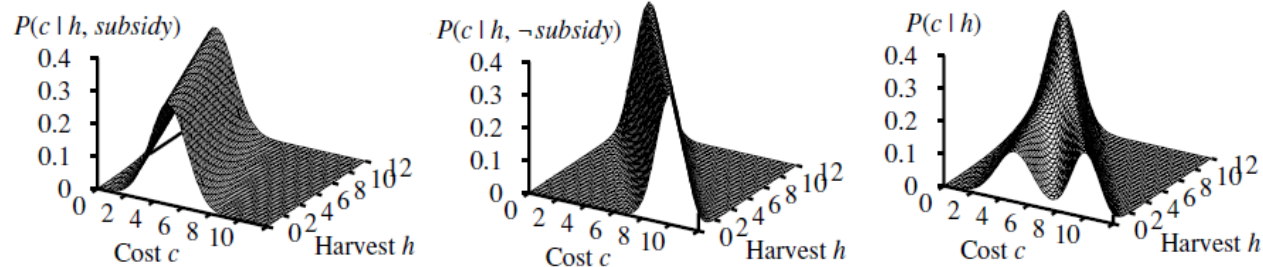
$$P(c | h, \text{subsidy}) = N(a_t h + b_t, \sigma_t^2)(c) = \frac{1}{\sigma_t \sqrt{2\pi}} e^{-\frac{1}{2} \left(\frac{c - (a_t h + b_t)}{\sigma_t} \right)^2}$$

$$P(c | h, \neg \text{subsidy}) = N(a_f h + b_f, \sigma_f^2)(c) = \frac{1}{\sigma_f \sqrt{2\pi}} e^{-\frac{1}{2} \left(\frac{c - (a_f h + b_f)}{\sigma_f} \right)^2}$$

- Mean **Cost** varies *linearly* with **Harvest**, variance fixed
- Linear variation is unreasonable over the full range but works if the likely range of **Harvest** is narrow

Continuous child variables

- Determine a Gaussian for *subsidy* and \neg *subsidy*
- What happens if subsidy is not given $P(c|h)$?



All-continuous network with LG distributions

⇒ full joint distribution is a multivariate Gaussian

Discrete+continuous LG network is a conditional Gaussian network i.e., a multivariate Gaussian over all continuous variables for each combination of discrete variable values

Discrete variabel cont. parents

- Probability of Buys given Cost should be a soft threshold

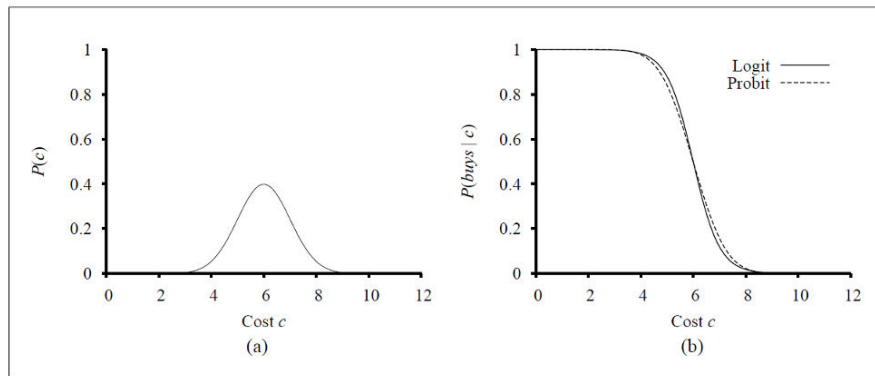


Figure 14.7 FILES: . (a) A normal (Gaussian) distribution for the cost threshold, centered on $\mu = 6.0$ with standard deviation $\sigma = 1.0$. (b) Logit and probit distributions for the probability of *buys* given *cost*, for the parameters $\mu = 6.0$ and $\sigma = 1.0$.

Use integral	$\Phi(x) = \int^x N(0, 1)(x)dx$	
Leads to	$P(\text{buys} \mid \text{Cost} = c) = \Phi((-c + \mu)/\sigma)$	Probit
Alternativ	$P(\text{buys} \mid \text{Cost} = c) = \frac{1}{1 + \exp(-2\frac{-c+\mu}{\sigma})}$	Logit

50

Inference tasks

- **Simple queries:** $P(X_1, \dots, X_n | e_2, e_4, e_5)$
- **Optimal decisions:** decision networks include utility information; inference must handle utility nodes.
- **Value of information:** which evidence to seek next?
- **Sensitivity analysis:** which probability values are more critical?
- **Explanation:** why do I need a new engine?

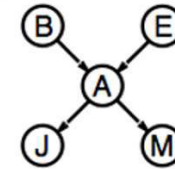
Inference by enumeration

$$P(b|j,m) = \alpha P(b,j,m)$$

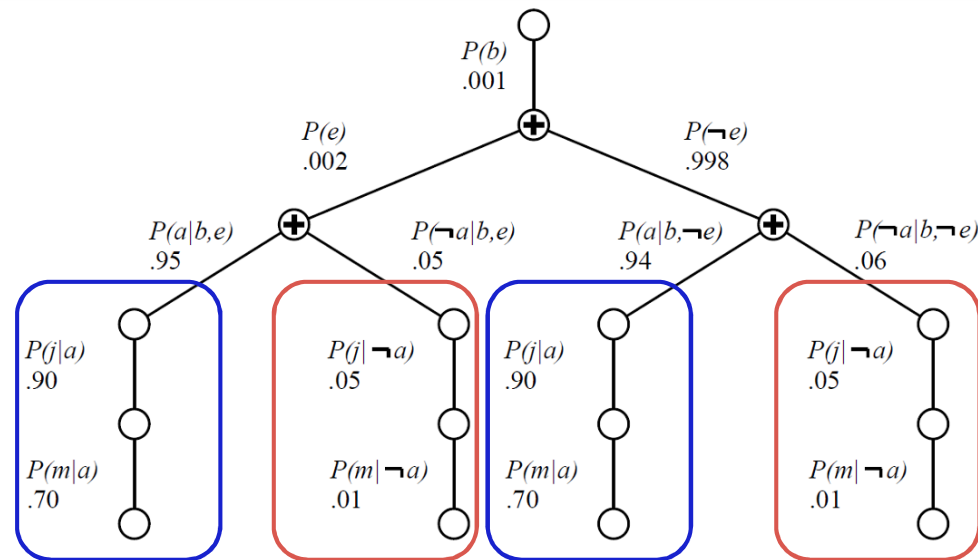
$$= \alpha \sum_a \sum_e P(b \wedge j \wedge m \wedge a \wedge e) \text{ [marginalization]}$$

$$= \alpha \sum_a \sum_e P(b)P(e)P(a|b,e)P(j|a)P(m|a) \text{ [BN]}$$

$$= \alpha P(b) \sum_e P(e) \sum_a P(a|b,e)P(j|a)P(m|a) \text{ [re-ordering]}$$



Evaluation Tree

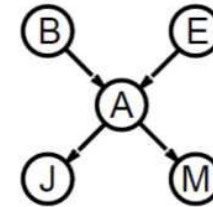


Enumeration is inefficient: repeated computation
 e.g., computes $P(j|a)P(m|a)$ for each value of e

Irrelevant variables

Consider the query $P(\text{JohnCalls} | \text{Burglary} = \text{true})$

$$P(J|b) = \alpha P(b) \sum_e P(e) \sum_a P(a|b, e) P(J|a) \sum_m P(m|a)$$



*What about **M**?*

*We sum over all possible values of **m***

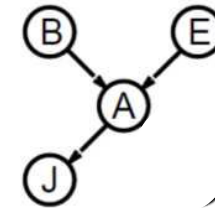
For each row it means that the value is 1

Irrelevant variables

Consider the query $P(\text{JohnCalls} | \text{Burglary} = \text{true})$

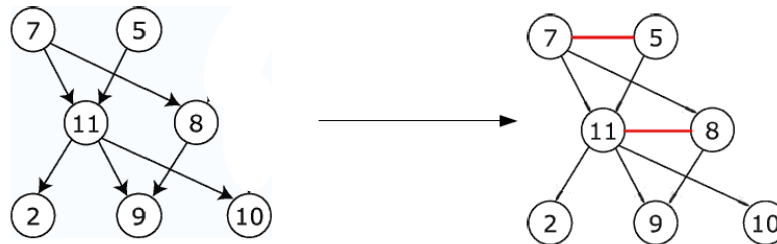
$$P(J|b) = \alpha P(b) \sum_e P(e) \sum_a P(a|b, e) P(J|a) \sum_m P(m|a)$$

For each row it means that the value is 1



Moral Graph: Markov Blanket

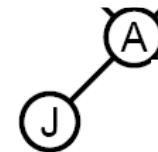
- The moral graph is an undirected graph that is obtained as follows:
 - ♦ connect all parents of all nodes
 - ♦ make all directed links undirected
- Note:
 - ♦ the moral graph connects each node to all nodes of its **Markov blanket**
 - it is already connected to parents and children
 - now it is also connected to the parents of its children



64

Irrelevant variables continued:

- m-separation:
 - ♦ A is m-separated from B by C iff it is separated by C in the moral graph
- Example:
 - ♦ J is m-separated from E by A



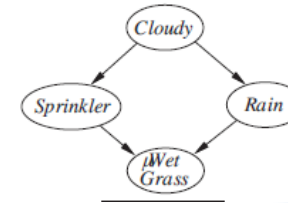
Theorem 2: Y is irrelevant if it is m-separated from X by E

- Example:

*For $P(\text{JohnCalls} | \text{Alarm}=\text{true})$,
Burglary, Earthquake and MarryCalls are irrelevant.*

Approximate Inference In Bayesian Network

- Singly connected networks (or polytrees):
 - ♦ any two nodes are connected by at most one (undirected) path
 - ♦ time and space cost of variable elimination linear in the size of the network (number of CPT entries; number of parents $O(d^k n)$).
- Multiply connected networks: NP-hard
- We need **approximate inference techniques!!!!!!**
- Monte Carlo algorithm
 - ♦ Widely used to estimate quantities that are difficult to calculate exactly
 - ♦ Randomized sampling algorithm
 - ♦ Accuracy depends on the number of samples
 - ♦ Two families
 - Direct sampling
 - Markov chain sampling



66

Inference by stochastic simulation

Basic idea:

- 1) Draw N samples from a sampling distribution S
- 2) Compute an approximate posterior probability \hat{P}
- 3) Show this converges to the true probability P



Outline:

- Sampling from an empty network
- Rejection sampling: reject samples disagreeing with evidence
- Likelihood weighting: use evidence to weight samples
- Markov chain Monte Carlo (MCMC): sample from a stochastic process whose stationary distribution is the true posterior

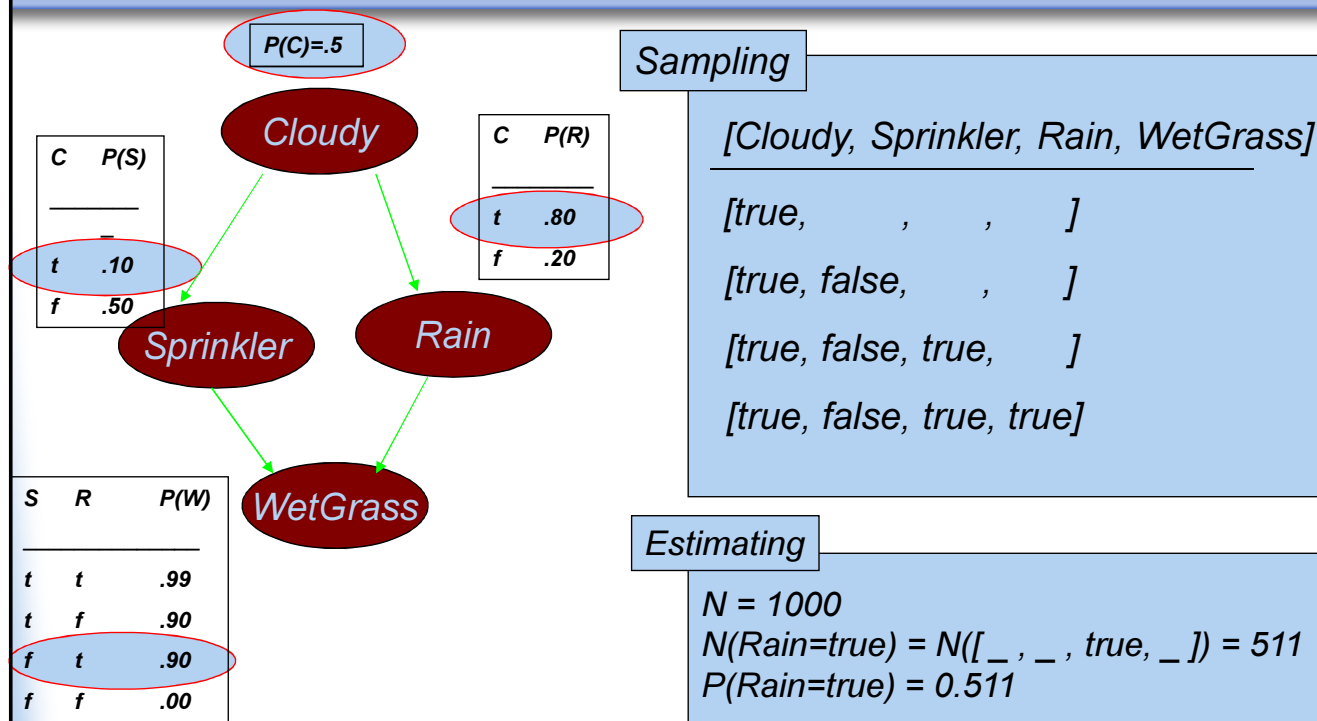
Sampling from empty network

- Generating samples from a network that has no evidence associated with it (*empty* network)
- Basic idea
 - ◆ sample a value for each variable in topological order
 - ◆ using the specified conditional probabilities

```
function PRIOR-SAMPLE(bn) returns an event sampled from bn
  inputs: bn, a belief network specifying joint distribution  $\mathbf{P}(X_1, \dots, X_n)$ 
  x ← an event with n elements
  for i = 1 to n do
    xi ← a random sample from  $\mathbf{P}(X_i \mid \text{parents}(X_i))$ 
      given the values of Parents(Xi) in x
  return x
```

68

Example in simple case



Sampling

[Cloudy, Sprinkler, Rain, WetGrass]

[true, , ,]

[true, false, ,]

[true, false, true,]

[true, false, true, true]

Estimating

$N = 1000$

$N(\text{Rain}=\text{true}) = N([_, _, \text{true}, _]) = 511$

$P(\text{Rain}=\text{true}) = 0.511$

Properties

Probability that PRIORSAMPLE generates a particular event

$$S_{PS}(x_1 \dots x_n) = \prod_{i=1}^n P(x_i | \text{parents}(X_i)) = P(x_1 \dots x_n)$$

i.e., the true prior probability

E.g., $S_{PS}(t, f, t, t) = 0.5 \times 0.9 \times 0.8 \times 0.9 = 0.324 = P(t, f, t, t)$

Let $N_{PS}(x_1 \dots x_n)$ be the number of samples generated for event x_1, \dots, x_n

Then we have

$$\begin{aligned} \lim_{N \rightarrow \infty} \hat{P}(x_1, \dots, x_n) &= \lim_{N \rightarrow \infty} N_{PS}(x_1, \dots, x_n) / N \\ &= S_{PS}(x_1, \dots, x_n) \\ &= P(x_1 \dots x_n) \end{aligned}$$

That is, estimates derived from PRIORSAMPLE are consistent

Shorthand: $\hat{P}(x_1, \dots, x_n) \approx P(x_1 \dots x_n)$

Rejection Sampling

- Used to compute conditional probabilities
- Procedure
 - ◆ Generating sample from prior distribution specified by the Bayesian Network
 - ◆ Rejecting all that do not match the evidence
 - ◆ Estimating probability

Rejection Sampling

$\hat{P}(X|e)$ estimated from samples agreeing with e

```
function REJECTION-SAMPLING( $X, e, bn, N$ ) returns an estimate of  $P(X|e)$ 
  local variables:  $N$ , a vector of counts over  $X$ , initially zero
  for  $j = 1$  to  $N$  do
     $x \leftarrow$  PRIOR-SAMPLE( $bn$ )
    if  $x$  is consistent with  $e$  then
       $N[x] \leftarrow N[x] + 1$  where  $x$  is the value of  $X$  in  $x$ 
  return NORMALIZE( $N[X]$ )
```

72

Rejection Sampling *Example*

- Let us assume we want to estimate $P(\text{Rain}|\text{Sprinkler} = \text{true})$ with 100 samples
- 100 samples
 - ♦ 73 samples \Rightarrow Sprinkler = false
 - ♦ 27 samples \Rightarrow Sprinkler = true
 - 8 samples \Rightarrow Rain = true
 - 19 samples \Rightarrow Rain = false
- $P(\text{Rain}|\text{Sprinkler} = \text{true}) = \text{NORMALIZE}(\{8,19\}) = \{0.296,0.704\}$
- The true answer ist $\langle 0.3,0.7 \rangle$
- Problem
 - ♦ It rejects too many samples

Analysis of rejection sampling

$$\begin{aligned}\hat{P}(X|\mathbf{e}) &= \alpha N_{PS}(X, \mathbf{e}) && \text{(algorithm defn.)} \\ &= N_{PS}(X, \mathbf{e}) / N_{PS}(\mathbf{e}) && \text{(normalized by } N_{PS}(\mathbf{e})\text{)} \\ &\approx P(X, \mathbf{e}) / P(\mathbf{e}) && \text{(property of PRIORSAMPLE)} \\ &= P(X|\mathbf{e}) && \text{(defn. of conditional probability)}\end{aligned}$$

Hence rejection sampling returns consistent posterior estimates

Problem: hopelessly expensive if $P(\mathbf{e})$ is small

$P(\mathbf{e})$ drops off exponentially with number of evidence variables!

Likelihood Weighting

- Goal
 - ◆ Avoiding inefficiency of rejection sampling
- Idea
 - ◆ Generating only events consistent with evidence
 - ◆ Each event is weighted by likelihood that the event accords to the evidence

Likelihood weighting

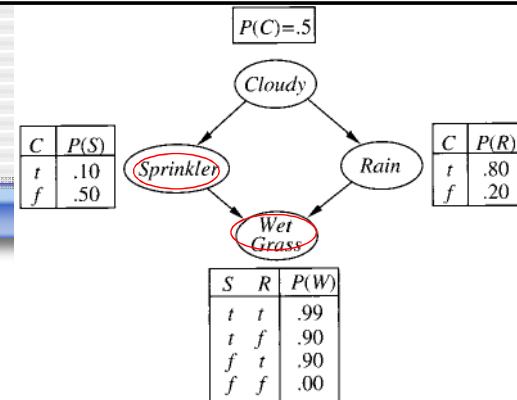
Idea: fix evidence variables, sample only nonevidence variables, and weight each sample by the likelihood it accords the evidence

```
function LIKELIHOOD-WEIGHTING( $X, e, bn, N$ ) returns an estimate of  $P(X|e)$ 
  local variables:  $\mathbf{W}$ , a vector of weighted counts over  $X$ , initially zero
  for  $j = 1$  to  $N$  do
     $\mathbf{x}, w \leftarrow$  WEIGHTED-SAMPLE( $bn, e$ )
     $\mathbf{W}[x] \leftarrow \mathbf{W}[x] + w$  where  $x$  is the value of  $X$  in  $\mathbf{x}$ 
  return NORMALIZE( $\mathbf{W}[X]$ )
```

```
function WEIGHTED-SAMPLE( $bn, e$ ) returns an event and a weight
   $\mathbf{x} \leftarrow$  an event with  $n$  elements;  $w \leftarrow 1$ 
  for  $i = 1$  to  $n$  do
    if  $X_i$  has a value  $x_i$  in  $e$ 
      then  $w \leftarrow w \times P(X_i = x_i \mid \text{parents}(X_i))$ 
      else  $x_i \leftarrow$  a random sample from  $\mathbf{P}(X_i \mid \text{parents}(X_i))$ 
  return  $\mathbf{x}, w$ 
```

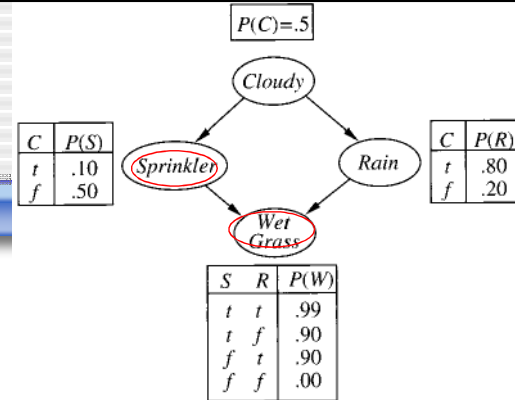
77

Likelihood Weighting Example



- $P(\text{Rain} | \text{Sprinkler}=\text{true}, \text{WetGrass} = \text{true})?$
- Sampling, start with weight=1
 - ♦ Sample from $P(\text{Cloudy}) = \{0.5, 0.5\} \Rightarrow \text{true}$
 - ♦ Sprinkler is an evidence variable with value **true**
 $w \leftarrow w * P(\text{Sprinkler}=\text{true} | \text{Cloudy} = \text{true}) = 0.1$
 - ♦ Sample from $P(\text{Rain} | \text{Cloudy}=\text{true}) = \{0.8, 0.2\} \Rightarrow \text{true}$
 - ♦ WetGrass is an evidence variable with value **true**
 $w \leftarrow w * P(\text{WetGrass}=\text{true} | \text{Sprinkler}=\text{true}, \text{Rain} = \text{true}) = 0.099$
 - ♦ [**true**, **true**, **true**, **true**] with weight 0.099

Likelihood Weighting Example



- $P(\text{Rain} | \text{Sprinkler}=\text{true}, \text{WetGrass} = \text{true})?$
- Sampling, start with weight=1
 - ♦ Sample from $P(\text{Cloudy}) = \{0.5, 0.5\} \Rightarrow \text{false}$
 - ♦ Sprinkler is an evidence variable with value **true**
 $w \leftarrow w * P(\text{Sprinkler}=\text{true} | \text{Cloudy} = \text{false}) = 0.5$
 - ♦ Sample from $P(\text{Rain} | \text{Cloudy}=\text{false}) = \{0.2, 0.8\} \Rightarrow \text{false}$
 - ♦ WetGrass is an evidence variable with value **true**
 $w \leftarrow w * P(\text{WetGrass}=\text{true} | \text{Sprinkler}=\text{true}, \text{Rain} = \text{false}) = 0.45$
 - ♦ [**true**, **true**, **true**, **true**] with weight 0.45
- Estimating
 - ♦ Accumulating weights to either Rain=true or Rain=false
 - ♦ Normalize

Likelihood analysis

Sampling probability for WEIGHTEDSAMPLE is

$$S_{WS}(\mathbf{z}, \mathbf{e}) = \prod_{i=1}^l P(z_i | \text{parents}(Z_i))$$

Note: pays attention to evidence in **ancestors** only

⇒ somewhere “in between” prior and posterior distribution

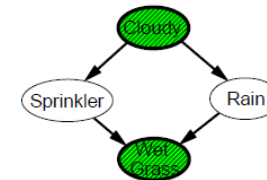
Weight for a given sample \mathbf{z}, \mathbf{e} is

$$w(\mathbf{z}, \mathbf{e}) = \prod_{i=1}^m P(e_i | \text{parents}(E_i))$$

Weighted sampling probability is

$$\begin{aligned} S_{WS}(\mathbf{z}, \mathbf{e}) w(\mathbf{z}, \mathbf{e}) &= \prod_{i=1}^l P(z_i | \text{parents}(Z_i)) \prod_{i=1}^m P(e_i | \text{parents}(E_i)) \\ &= P(\mathbf{z}, \mathbf{e}) \text{ (by standard global semantics of network)} \end{aligned}$$

Hence likelihood weighting returns consistent estimates
but performance still degrades with many evidence variables
because a few samples have nearly all the total weight



Markov Chain Monte Carlo

- Let's think of the network as being in a particular current state specifying a value for every variable
- MCMC generates each event by making a random change to the preceding event
- The next state is generated by randomly sampling a value for one of the non evidence variables X_i , **conditioned on the current values of the variables in the MarkovBlanket of X_i**
- Likelihood Weighting only takes into account the evidences of the parents.

Gibbs sampling

- Gibbs sampling is a MCMC method
 - ♦ State of the network => current assignment
 - ♦ Generate next state by sampling one non-evidence variable given Markov blanket
 - ♦ Sample each variable in turn (can choose it random)

```

function GIBBS-ASK( $X, e, bn, N$ ) returns an estimate of  $\mathbf{P}(X|e)$ 
  local variables:  $\mathbf{N}$ , a vector of counts for each value of  $X$ , initially zero
                    $\mathbf{Z}$ , the nonevidence variables in  $bn$ 
                    $\mathbf{x}$ , the current state of the network, initially copied from  $e$ 

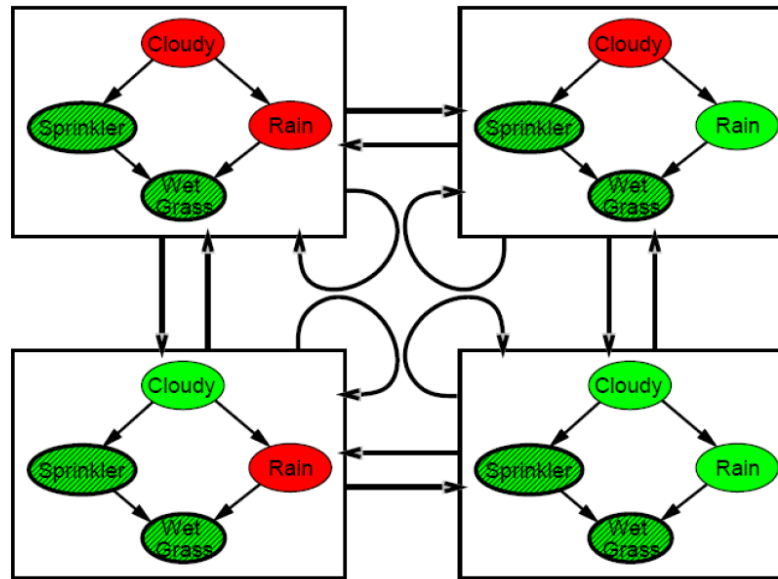
  initialize  $\mathbf{x}$  with random values for the variables in  $\mathbf{Z}$ 
  for  $j = 1$  to  $N$  do
    for each  $Z_i$  in  $\mathbf{Z}$  do
      set the value of  $Z_i$  in  $\mathbf{x}$  by sampling from  $\mathbf{P}(Z_i|mb(Z_i))$ 
       $\mathbf{N}[x] \leftarrow \mathbf{N}[x] + 1$  where  $x$  is the value of  $X$  in  $\mathbf{x}$ 
  return NORMALIZE( $\mathbf{N}$ )

```

Figure 14.16 The Gibbs sampling algorithm for approximate inference in Bayesian networks; this version cycles through the variables, but choosing variables at random also works.

Example

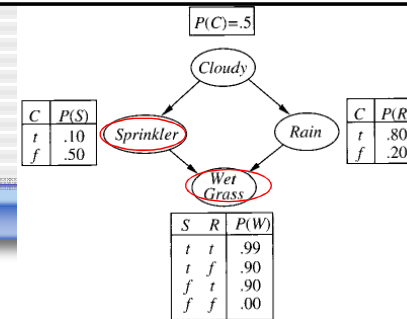
With *Sprinkler = true*, *WetGrass = true*, there are four states:



Wander about for a while, average what you see

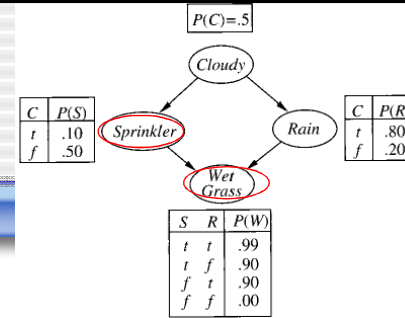
84

Gibbs Example



- Query $P(\text{Rain} \mid \text{Sprinkler} = \text{true}, \text{WetGrass} = \text{true})$
- Initial state is [true, true, false, true] [Cloudy, Sprinkler, Rain, WetGrass]
- The following steps are executed repeatedly:
 - ♦ Cloudy is sampled, given the current values of its Markov Blanket variables
So, we sample from $P(\text{Cloudy} \mid \text{Sprinkler} = \text{true}, \text{Rain} = \text{false})$
The result is Cloudy = false (???????)
 - ♦ Now current state is [false, true, false, true] and counts are updated
 - ♦ Rain is sampled, given the current values of its Markov Blanket variables
Sample from $P(\text{Rain} \mid \text{Cloudy} = \text{false}, \text{Sprinkler} = \text{true}, \text{WetGrass} = \text{true})$
First create the distribution we want to sample from.
→ Rain = true.
 - ♦ Current state is [false, true, true, true]
- After all the iterations, let's say the process visited 20 states where rain is true and 60 states where rain is false then the answer of the query is $\text{NORMALIZE}(\{20, 60\}) = \{0.25, 0.75\}$

Sample distribution



Want to sample *Cloudy*.

The current state is [*Cloudy?*, *true*, *false*, *true*]

What is the Markov blanket, the sampling distribution?

evidence *sampled*

$P(\text{Cloudy} \mid \text{Sprinkler}=\textit{true}, \text{Rain}=\textit{false}) =$

$$\alpha P(\text{Cloudy}) * P(\text{Sprinkler}=\textit{true} \mid \text{Cloudy}) P(\text{Rain}=\textit{false} \mid \text{Cloudy}) =$$

$$\alpha \langle 0.5, 0.5 \rangle * \langle 0.1, 0.5 \rangle * \langle 0.2, 0.8 \rangle =$$

$$\alpha \langle 0.5, 0.5 \rangle * \langle 0.1 * 0.2, 0.5 * 0.8 \rangle =$$

$$\alpha \langle 0.5, 0.5 \rangle * \langle 0.02, 0.4 \rangle =$$

$$\alpha \langle 0.01, 0.2 \rangle \approx \langle 0.05, 0.95 \rangle$$

[*false*, *true*, *false*, *true*] with probability 0,95

[*true*, *true*, *false*, *true*] with probability 0,05

86

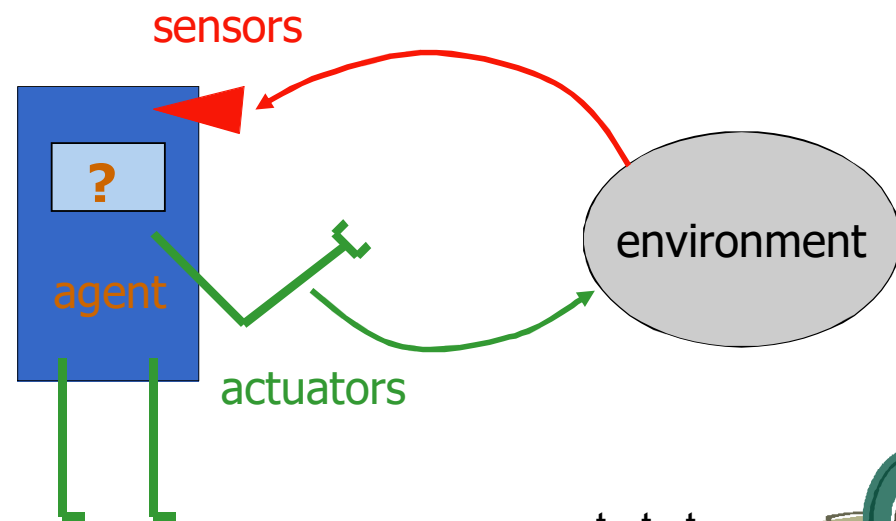
Summary

- Bayesian networks provide a natural representation for (causally induced) conditional independence
- Topology + CPTs = compact representation of joint distribution
- Generally easy for domain experts to construct (if not too big)
- Exact inference by variable elimination
 - ♦ polytime on polytrees, NP-hard on general graphs
 - ♦ space can be exponential as well
- Approximate inference based on sampling and counting help to overcome complexity of exact inference

**Intelligent Autonomous Agents
and Cognitive Robotics**
**Topic 6: Probabilistic Reasoning
over Time**
(Dynamic Bayesian Networks)

Ralf Möller, Rainer Marrone
Hamburg University of Technology

Temporal Probabilistic Agent



So far we only have taken care about one moment in time !!!!!

t_1, t_2, t_3, \dots



Time and Uncertainty

- The world changes over time, we need to track and predict it
- Examples:
diabetes management, localization, speech recognition, ...
- Basic idea: copy state and evidence variables for each time step
- \mathbf{X}_t – set of unobservable state variables at time t
 - ♦ e.g., BloodSugar _{t} , StomachContents _{t} , ...
- \mathbf{E}_t – set of evidence variables at time t
 - ♦ e.g., MeasuredBloodSugar _{t} , PulseRate _{t} , FoodEaten _{t} , ...
- Assumes discrete time steps

Dynamic Bayesian Networks

- How can we model *dynamic* situations with a Bayesian network?
- Example: *Is it raining today?*

$$X_t = \{R_t\}$$

$$E_t = \{U_t\}$$

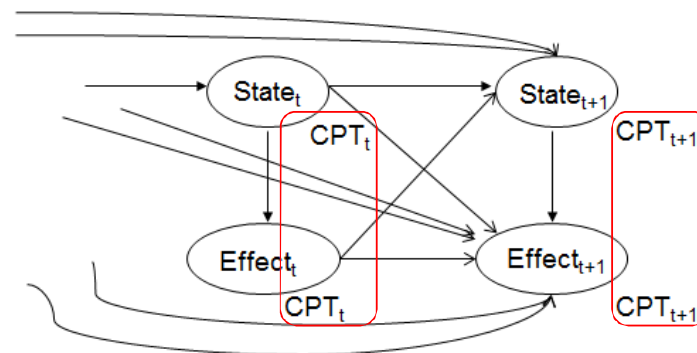
➡ next step: specify dependencies among the variables.

The term “dynamic” means we are modeling a dynamic system, not that the network structure changes over time.

4

DBN - Representation

- Problem:
 1. Necessity to specify an unbounded number of conditional probability tables, one for each variable in each slice,
 2. Each one might involve an unbounded number of parents.



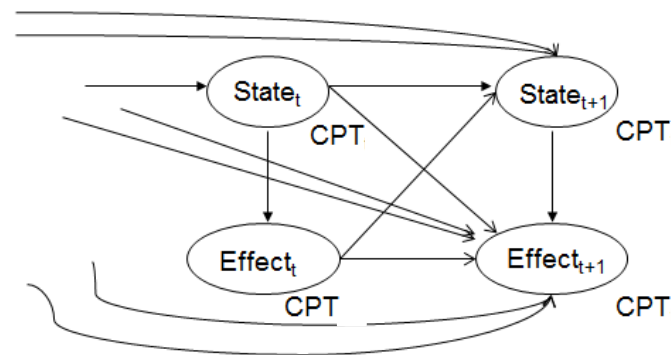
DBN - Representation

- Problem:
 1. Necessity to specify an unbounded number of conditional probability tables, one for each variable in each slice,
 2. Each one might involve an unbounded number of parents.
- Solution:
 1. Assume that changes in the world state are caused by a stationary process (the laws for a state change do not change over time).

$P(U_t / Parent(U_t))$ is the same for all t

DBN - Representation

- Problem:
 1. Necessity to specify an unbounded number of conditional probability tables, one for each variable in each slice \rightarrow solved
 2. Each one might involve an unbounded number of parents.



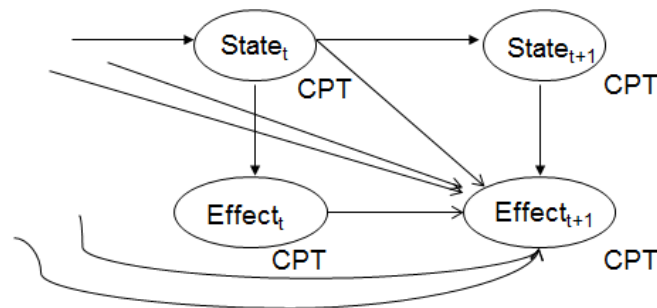
DBN - Representation

- Solution cont.:
 2. Use **Markov assumption** - The current state depends on only a finite history of previous states.

Using the first-order Markov process:

$$P(X_t / X_{0:t-1}) = P(X_t / X_{t-1})$$

Transition Model



DBN - Representation

- Solution cont.:
 2. Use **Markov assumption** - The current state depends on only a finite history of previous states.

Using the first-order Markov process:

$$P(X_t / X_{0:t-1}) = P(X_t / X_{t-1}) \quad \text{Transition Model}$$

In addition to restricting the parents of the state variable X_t , we must restrict the parents of the evidence variable E_t

$$P(E_t / X_{0:t}, E_{0:t-1}) = P(E_t / X_t) \quad \text{Sensor Model}$$

DBN - Representation

- Solution cont.:
 2. Use **Markov assumption** - The current state depends on only in a finite history of previous states.

Using the first-order Markov process:

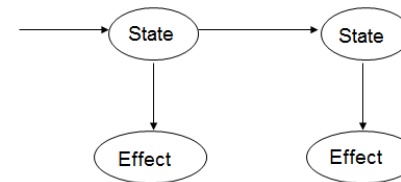
$$P(X_t / X_{0:t-1}) = P(X_t / X_{t-1})$$

Transition Model

In addition to restricting the parents of the state variable X_t , we must restrict the parents of the evidence variable E_t

Sensor Model

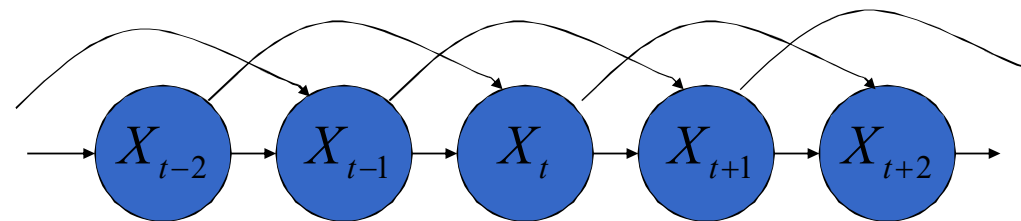
$$P(E_t / X_{0:t}, E_{0:t-1}) = P(E_t / X_t)$$



10

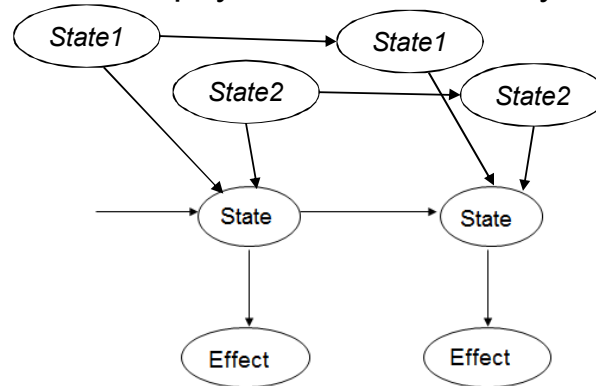
Dynamic Bayesian Networks

- There are two possible fixes if the approximation is too inaccurate:
 - ♦ Increasing the order of the Markov process model. For example, adding \mathbf{Rain}_{t-2} as a parent of \mathbf{Rain}_t , which might give slightly more accurate predictions.



Dynamic Bayesian Networks

- There are two possible fixes if the approximation is too inaccurate:
 - ♦ Increasing the set of state variables. For example, adding **Season_t** to allow to incorporate historical records of rainy seasons, or adding **Temperature_t**, **Humidity_t** and **Pressure_t** to allow to use a physical model of rainy conditions.

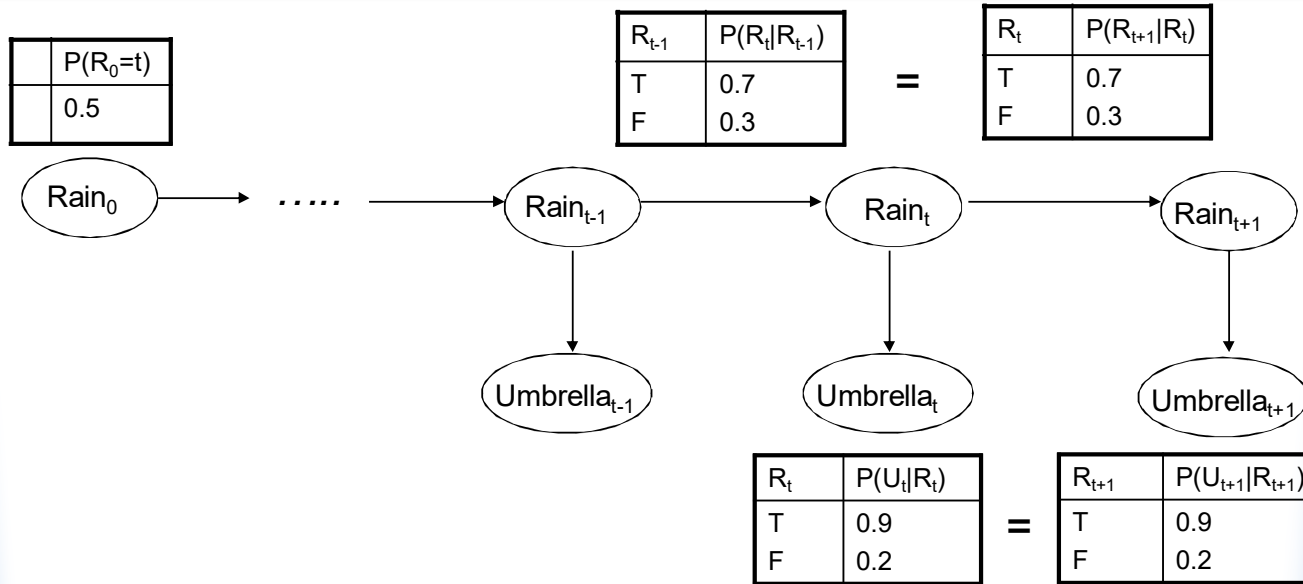


Complete Joint Distribution

- Given:
 - ♦ Transition model: $P(X_t|X_{t-1})$
 - ♦ Sensor model: $P(E_t|X_t)$
 - ♦ Prior probability: $P(X_0)$
- Then we can specify complete joint distribution:

$$P(X_0, X_1, \dots, X_t, E_1, \dots, E_t) = P(X_0) \prod_{i=1}^t P(X_i | X_{i-1}) P(E_i | X_i)$$

Simple Example



Inference Tasks: Examples

- **Filtering/State estimation:**
What is the probability that it is raining today, given all the umbrella observations up through today?
- **Prediction:**
What is the probability that it will rain the day after tomorrow, given all the umbrella observations up through today?
- **Smoothing:**
What is the probability that it rained yesterday, given all the umbrella observations through today?
- **Most likely explanation:**
If the umbrella appeared the first three days but not on the fourth, what is the most likely weather sequence to produce these umbrella sightings?

DBN – Basic Inference

- Filtering or Monitoring:

Compute the **belief state** - the posterior distribution over the *current* state, given all evidence to date.

$$P(X_t / e_{1:t})$$

Filtering is what a rational agent needs to do in order to keep track of the **current state** so that the rational decisions can be made.

DBN – Basic Inference

- Filtering cont.

$$P(B|A,C) = \alpha P(A|B,C) P(B|C)$$

Given the results of filtering up to time t , one can easily compute the result for $t+1$ from the new evidence e_{t+1}

$$P(X_{t+1} / e_{1:t+1}) = f(e_{t+1}, P(X_t / e_{1:t}))$$

(seeking for some recursive function f ?)

$$= P(X_{t+1} / e_{1:t}, e_{t+1})$$

(dividing up the evidence)

$$= \alpha P(e_{t+1} / X_{t+1}, e_{1:t}) P(X_{t+1} / e_{1:t})$$

(using Bayes' Theorem)

$$= \alpha P(e_{t+1} / X_{t+1}) P(X_{t+1} / e_{1:t})$$

(by the Markov property of evidence)

DBN – Basic Inference

- Filtering cont.

$P(X_{t+1} / e_{1:t})$ represents a one-step prediction

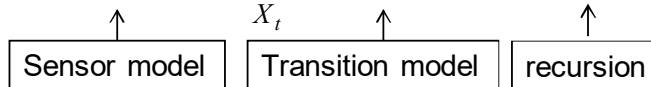
$P(e_{t+1} | X_{t+1})$ updates this with the new evidence

$$P(X_{t+1} / e_{1:t+1}) = \alpha P(e_{t+1} / X_{t+1}) P(X_{t+1} / e_{1:t})$$

$$P(X_{t+1} / e_{1:t+1}) = \alpha P(e_{t+1} / X_{t+1}) \sum_{X_t} P(X_{t+1} / x_t, e_{1:t}) P(x_t / e_{1:t})$$

(using the Markov property)

$$= \alpha P(e_{t+1} / X_{t+1}) \sum_{X_t} P(X_{t+1} / x_t) P(x_t / e_{1:t})$$



DBN – Basic Inference

For two steps in the Umbrella example: $= \alpha P(e_{t+1} / X_{t+1}) \sum_{X_t} P(X_{t+1} / x_t) P(x_t / e_{1:t})$

- On day 1, the umbrella appears so $U_1 = \text{true}$. The prediction from $t=0$ to $t=1$ is

$$P(R_1) = \sum_{r_0} P(R_1 / r_0) P(r_0)$$

and updating it with the evidence for $t=1$ gives

$$P(R_1 / u_1) = \alpha P(u_1 / R_1) P(R_1)$$

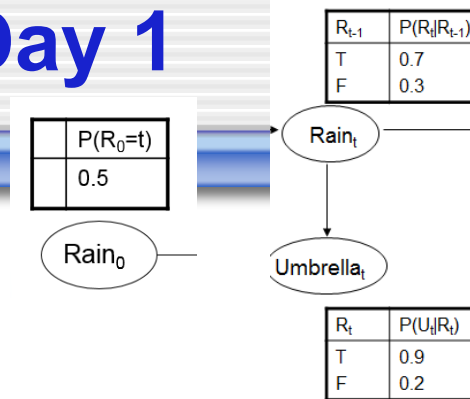
- On day 2, the umbrella appears so $U_2 = \text{true}$. The prediction from $t=1$ to $t=2$ is

$$P(R_2 / u_1) = \sum_{r_1} P(R_2 / r_1) P(r_1 / u_1)$$

and updating it with the evidence for $t=2$ gives

$$P(R_2 / u_1, u_2) = \alpha P(u_2 / R_2) P(R_2 / u_1)$$

Example: Day 1



$$P(R_1 | u_1) = P(u_1 / R_1) \sum_{r_0} P(R_1 / r_0) P(r_0)$$

evidence prediction

$$P(R_0) = \langle 0.5, 0.5 \rangle$$

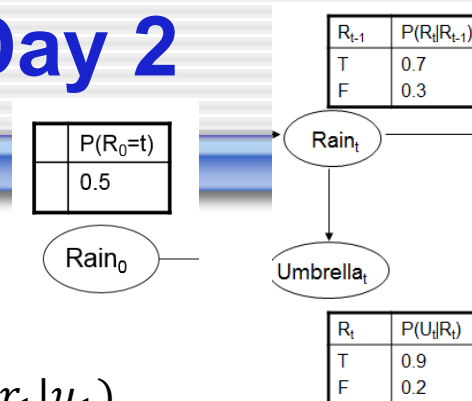
$$P(R_1) = \sum_{r_0} P(R_1 | r_0) P(r_0) = \langle 0.7, 0.3 \rangle 0.5 + \langle 0.3, 0.7 \rangle 0.5 = \langle 0.5, 0.5 \rangle$$

$$P(R_1 | u_1) = \alpha P(u_1 | R_1) P(R_1) = \alpha \langle 0.9, 0.2 \rangle \langle 0.5, 0.5 \rangle$$

$$= \alpha \langle 0.45, 0.1 \rangle \approx \langle 0.818, 0.182 \rangle$$

21

Example: Day 2



$$P(R_2|u_1, u_2) = \alpha P(u_2|R_2) \sum_{r_1} P(R_2|r_1) P(r_1|u_1)$$

evidence prediction

$$P(R_1|u_1) \approx \langle 0.818, 0.182 \rangle$$

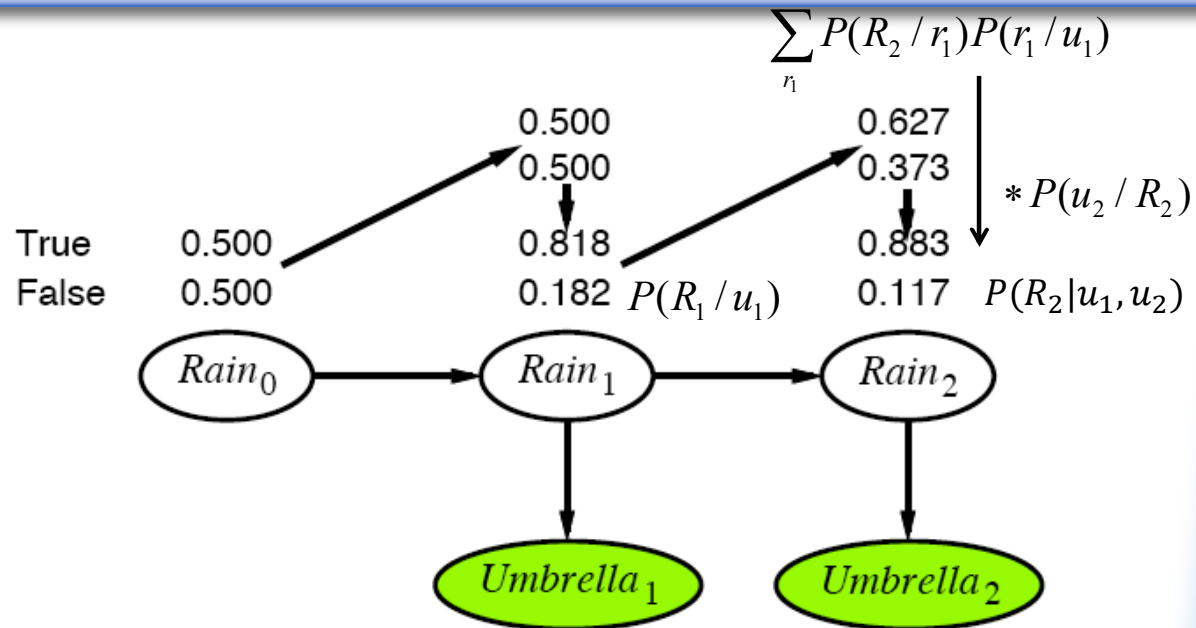
$$P(R_2|u_1) = \sum_{r_1} P(R_2|r_1) P(r_1|u_1) =$$

$$\langle 0.7, 0.3 \rangle \cdot 0.818 + \langle 0.3, 0.7 \rangle \cdot 0.182 \approx \langle 0.627, 0.373 \rangle$$

$$P(R_2|u_1, u_2) = \alpha P(u_2|R_2) P(R_2|u_1) = \alpha \langle 0.9, 0.2 \rangle \langle 0.627, 0.373 \rangle$$

$$= \alpha \langle 0.565, 0.075 \rangle \approx \langle 0.883, 0.117 \rangle$$

Example



23

DBN – Basic Inference

- Prediction:

Compute the posterior distribution over the *future* state, given all evidence to date.

$$P(X_{t+k+1} / e_{1:t}) = \sum_{x_{t+k}} P(X_{t+k+1} | x_{t+k}) P(x_{t+k} | e_{1:t})$$

for some $k > 0$

The task of **prediction** can be seen simply as filtering without the addition of new evidence.

DBN – Basic Inference

- Smoothing or hindsight:

Compute the posterior distribution over the *past* state, given all evidence up to the present.

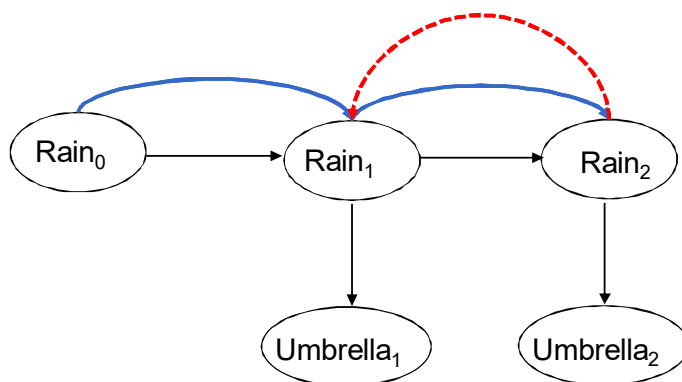
$$P(X_k / e_{1:t}) \quad \text{for some } k \text{ such that } 0 \leq k < t.$$

Hindsight provides a better estimate of the state than was available at the time, because it incorporates more evidence.

25

Smoothing

- Can I use future information to increase the accuracy of filtering for past states?



Umbrella₁=t Umbrella₂=t

Smoothing

Divide evidence $e_{1:t}$ into $e_{1:k}$, $e_{k+1:t}$:

$$\begin{aligned} P(\mathbf{X}_k | e_{1:t}) &= P(\mathbf{X}_k | e_{1:k}, e_{k+1:t}) \\ &= \alpha P(\mathbf{X}_k | e_{1:k}) P(e_{k+1:t} | \mathbf{X}_k, e_{1:k}) && \text{Bayes rule} \\ &= \alpha P(\mathbf{X}_k | e_{1:k}) P(e_{k+1:t} | \mathbf{X}_k) && \text{Markov} \\ &= \alpha f_{1:k} b_{k+1:t} \end{aligned}$$

Smoothing

Divide evidence $\mathbf{e}_{1:t}$ into $\mathbf{e}_{1:k}$, $\mathbf{e}_{k+1:t}$:

$$\begin{aligned}\mathbf{P}(\mathbf{X}_k | \mathbf{e}_{1:t}) &= \mathbf{P}(\mathbf{X}_k | \mathbf{e}_{1:k}, \mathbf{e}_{k+1:t}) \\ &= \alpha \mathbf{P}(\mathbf{X}_k | \mathbf{e}_{1:k}) \mathbf{P}(\mathbf{e}_{k+1:t} | \mathbf{X}_k, \mathbf{e}_{1:k}) \\ &= \alpha \mathbf{P}(\mathbf{X}_k | \mathbf{e}_{1:k}) \mathbf{P}(\mathbf{e}_{k+1:t} | \mathbf{X}_k) \\ &= \alpha \mathbf{f}_{1:k} \mathbf{b}_{k+1:t}\end{aligned}$$

Backward message computed by a backwards recursion:

$$\mathbf{P}(\mathbf{e}_{k+1:t} | \mathbf{X}_k) = \sum_{\mathbf{x}_{k+1}} \mathbf{P}(\mathbf{e}_{k+1:t} | \mathbf{X}_k, \mathbf{x}_{k+1}) \mathbf{P}(\mathbf{x}_{k+1} | \mathbf{X}_k)$$

Smoothing

Divide evidence $\mathbf{e}_{1:t}$ into $\mathbf{e}_{1:k}$, $\mathbf{e}_{k+1:t}$:

$$\begin{aligned} \mathbf{P}(\mathbf{X}_k | \mathbf{e}_{1:t}) &= \mathbf{P}(\mathbf{X}_k | \mathbf{e}_{1:k}, \mathbf{e}_{k+1:t}) \\ &= \alpha \mathbf{P}(\mathbf{X}_k | \mathbf{e}_{1:k}) \mathbf{P}(\mathbf{e}_{k+1:t} | \mathbf{X}_k, \mathbf{e}_{1:k}) \\ &= \alpha \mathbf{P}(\mathbf{X}_k | \mathbf{e}_{1:k}) \mathbf{P}(\mathbf{e}_{k+1:t} | \mathbf{X}_k) \\ &= \alpha \mathbf{f}_{1:k} \mathbf{b}_{k+1:t} \end{aligned}$$

Backward message computed by a backwards recursion:

$$\begin{aligned} \mathbf{P}(\mathbf{e}_{k+1:t} | \mathbf{X}_k) &= \sum_{\mathbf{x}_{k+1}} \mathbf{P}(\mathbf{e}_{k+1:t} | \mathbf{X}_k, \mathbf{x}_{k+1}) \mathbf{P}(\mathbf{x}_{k+1} | \mathbf{X}_k) \\ &= \sum_{\mathbf{x}_{k+1}} P(\mathbf{e}_{k+1:t} | \mathbf{x}_{k+1}) \mathbf{P}(\mathbf{x}_{k+1} | \mathbf{X}_k) \end{aligned}$$

Smoothing

Divide evidence $\mathbf{e}_{1:t}$ into $\mathbf{e}_{1:k}$, $\mathbf{e}_{k+1:t}$:

$$\begin{aligned}
 \mathbf{P}(\mathbf{X}_k | \mathbf{e}_{1:t}) &= \mathbf{P}(\mathbf{X}_k | \mathbf{e}_{1:k}, \mathbf{e}_{k+1:t}) \\
 &= \alpha \mathbf{P}(\mathbf{X}_k | \mathbf{e}_{1:k}) \mathbf{P}(\mathbf{e}_{k+1:t} | \mathbf{X}_k, \mathbf{e}_{1:k}) \\
 &= \alpha \mathbf{P}(\mathbf{X}_k | \mathbf{e}_{1:k}) \mathbf{P}(\mathbf{e}_{k+1:t} | \mathbf{X}_k) \\
 &= \alpha \mathbf{f}_{1:k} \mathbf{b}_{k+1:t}
 \end{aligned}$$

Backward message computed by a backwards recursion:

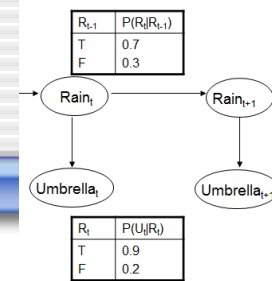
$$\begin{aligned}
 \mathbf{P}(\mathbf{e}_{k+1:t} | \mathbf{X}_k) &= \sum_{\mathbf{x}_{k+1}} \mathbf{P}(\mathbf{e}_{k+1:t} | \mathbf{X}_k, \mathbf{x}_{k+1}) \mathbf{P}(\mathbf{x}_{k+1} | \mathbf{X}_k) \\
 &= \sum_{\mathbf{x}_{k+1}} P(\mathbf{e}_{k+1:t} | \mathbf{x}_{k+1}) \mathbf{P}(\mathbf{x}_{k+1} | \mathbf{X}_k) \\
 &= \sum_{\mathbf{x}_{k+1}} \underset{\substack{\uparrow \\ \text{Sensor model}}}{P(\mathbf{e}_{k+1} | \mathbf{x}_{k+1})} \underset{\substack{\uparrow \\ \text{recursion}}}{P(\mathbf{e}_{k+2:t} | \mathbf{x}_{k+1})} \underset{\substack{\uparrow \\ \text{Transition model}}}{\mathbf{P}(\mathbf{x}_{k+1} | \mathbf{X}_k)}
 \end{aligned}$$

Sensor model

recursion

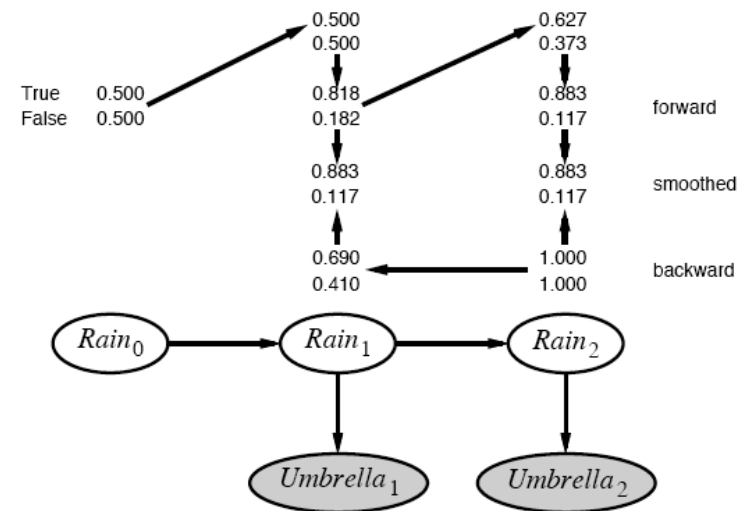
Transition model

Example



- Smoothed estimate for rain at $k=1$, given u_1, u_2 .
 $\mathbf{P}(R_1 | u_1, u_2) = \alpha \mathbf{P}(R_1 | u_1) \mathbf{P}(u_2 | R_1)$
- The first term is taken from the forward example
 $\langle 0.818, 0.182 \rangle$
- $\mathbf{P}(u_2 | R_1) = \sum_{r_2} P(u_2 | r_2) P(r_2) \mathbf{P}(r_2 | R_1)$
 $= (0.9 \times 1 \times \langle 0.7, 0.3 \rangle) + (0.2 \times 1 \times \langle 0.3, 0.7 \rangle)$
 $= \langle 0.69, 0.41 \rangle$
- $\mathbf{P}(R_1 | u_1, u_2) = \alpha \langle 0.818, 0.182 \rangle \times \langle 0.69, 0.41 \rangle$
 $\approx \langle 0.883, 0.117 \rangle$
- If we do it for each time slice $O(t^2)$!!!

Example contd.



Forward-backward algorithm: cache forward messages along the way
 Time linear in t (polytree inference), space $O(t|f|)$

Forward-Backward Algorithm

function FORWARD-BACKWARD(*ev*, *prior*) **returns** a vector of probability distributions

inputs: *ev*, a vector of evidence values for steps $1, \dots, t$

prior, the prior distribution on the initial state, $\mathbf{P}(\mathbf{X}_0)$

local variables: *fv*, a vector of forward messages for steps $0, \dots, t$

b, a representation of the backward message, initially all 1s

sv, a vector of smoothed estimates for steps $1, \dots, t$

fv[0] \leftarrow *prior*

for $i = 1$ **to** t **do**

fv[i] \leftarrow FORWARD(*fv*[$i - 1$], *ev*[i])

for $i = t$ **downto** 1 **do**

sv[i] \leftarrow NORMALIZE(*fv*[i] \times *b*)

b \leftarrow BACKWARD(*b*, *ev*[i])

return *sv*

DBN – Basic Inference

- Most likely explanation:

Compute the sequence of states that is most likely to have generated a given **sequence of observation**.

$$\operatorname{argmax}_{x_{1:t}} P(X_{1:t} / e_{1:t})$$

Algorithms for this task are useful in many applications, including, e.g., speech recognition. Can also be used to compare different temporal models that might have produced as sequence of events.

Most-likely explanation

Most likely path to each \mathbf{x}_{t+1}

= most likely path to **some** \mathbf{x}_t plus one more step

$$\begin{aligned} & \max_{\mathbf{x}_1 \dots \mathbf{x}_t} P(\mathbf{x}_1, \dots, \mathbf{x}_t, \mathbf{X}_{t+1} | \mathbf{e}_{1:t+1}) \\ & = P(\mathbf{e}_{t+1} | \mathbf{X}_{t+1}) \max_{\mathbf{x}_t} \left(P(\mathbf{X}_{t+1} | \mathbf{x}_t) \max_{\mathbf{x}_1 \dots \mathbf{x}_{t-1}} P(\mathbf{x}_1, \dots, \mathbf{x}_{t-1}, \mathbf{x}_t | \mathbf{e}_{1:t}) \right) \end{aligned}$$

Identical to filtering, except $\mathbf{f}_{1:t}$ replaced by

$$\mathbf{m}_{1:t} = \max_{\mathbf{x}_1 \dots \mathbf{x}_{t-1}} P(\mathbf{x}_1, \dots, \mathbf{x}_{t-1}, \mathbf{X}_t | \mathbf{e}_{1:t}),$$

I.e., $\mathbf{m}_{1:t}(i)$ gives the probability of the most likely path to state i .

Update has sum replaced by max, giving the Viterbi algorithm:

$$\mathbf{m}_{1:t+1} = P(\mathbf{e}_{t+1} | \mathbf{X}_{t+1}) \max_{\mathbf{x}_t} (P(\mathbf{X}_{t+1} | \mathbf{x}_t) \mathbf{m}_{1:t})$$

The occasionally dishonest casino

- A casino uses a fair die most of the time, but occasionally switches to a loaded one
 - ♦ Fair die: $\text{Prob}(1) = \dots = \text{Prob}(6) = 1/6$
 - ♦ Loaded die: $\text{Prob}(1) = \dots = \text{Prob}(5) = 1/10$, $\text{Prob}(6) = 1/2$
 - ♦ These are the **emission** probabilities
- **Transition probabilities**
 - ♦ $\text{Prob}(\text{Fair} \rightarrow \text{Loaded}) = 0.01$
 - ♦ $\text{Prob}(\text{Loaded} \rightarrow \text{Fair}) = 0.2$
 - ♦ Transitions between states modeled by a Markov process

The occasionally dishonest casino

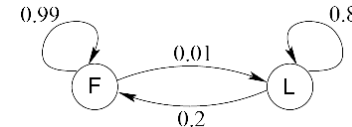
- Known:
 - ♦ The structure of the model
 - ♦ The transition probabilities
- Hidden: What the casino did
 - ♦ FFFFLLLLLLLLFFFF...
- Observable: The series of die tosses
 - ♦ 3415256664666153...
- What we must infer:
 - ♦ When was a fair/loading die used?
 - The answer is a sequence
FFFFFFFFLLLLLLLLFFF...

39

Making the inference

- Model assigns a probability to each explanation of the observation:

$$\begin{aligned}
 &P(326|\mathbf{FFL}) \\
 &= P(3|F) \cdot P(F \rightarrow F) \cdot P(2|F) \cdot P(F \rightarrow L) \cdot P(6|L) \\
 &= 1/6 \cdot 0.99 \cdot 1/6 \cdot 0.01 \cdot 1/2
 \end{aligned}$$



- Maximum Likelihood:** Determine which explanation is most likely
 - Find the path *most likely* to have produced the observed sequence
- Total probability:** Determine the probability that the observed sequence was produced by the model
 - Consider *all* paths that could have produced the observed sequence

Notation

- x is the sequence of *symbols/observations* emitted by the model
 - ♦ x_i is the symbol emitted at time i
- A **path**, π , is a sequence of *states*
 - ♦ The i -th state in π is π_i
- t_{kr} is the probability of making a transition from state k to state r :
$$t_{kr} = \Pr(\pi_i = r \mid \pi_{i-1} = k)$$
- $e_k(b)$ is the probability that symbol b is emitted when in state k
$$e_k(b) = \Pr(x_i = b \mid \pi_i = k)$$

41

The occasionally dishonest casino

$$x = \langle x_1, x_2, x_3 \rangle = \langle 6, 2, 6 \rangle$$

$$\pi^{(1)} = FFF$$

$$\begin{aligned} \Pr(x, \pi^{(1)}) &= t_{0F}e_F(6)t_{FF}e_F(2)t_{FF}e_F(6) \\ &= 0.5 \times \frac{1}{6} \times 0.99 \times \frac{1}{6} \times 0.99 \times \frac{1}{6} \\ &\approx 0.00227 \end{aligned}$$

$$\pi^{(2)} = LLL$$

$$\begin{aligned} \Pr(x, \pi^{(2)}) &= t_{0L}e_L(6)t_{LL}e_L(2)t_{LL}e_L(6) \\ &= 0.5 \times 0.5 \times 0.8 \times 0.1 \times 0.8 \times 0.5 \\ &= 0.008 \end{aligned}$$

$$\pi^{(3)} = LFL$$

$$\begin{aligned} \Pr(x, \pi^{(3)}) &= t_{0L}e_L(6)t_{LF}e_F(2)t_{FL}e_L(6) \\ &= 0.5 \times 0.5 \times 0.2 \times \frac{1}{6} \times 0.01 \times 0.5 \\ &\approx 0.0000417 \end{aligned}$$

43

The most likely path

The most likely path π^ satisfies*

$$\pi^* = \operatorname{argmax}_{\pi} \Pr(x, \pi)$$

To find π^ , consider all possible ways the last symbol of x could have been emitted*

Let

$p_k(i)$ = Prob. of path $\langle \pi_1, \dots, \pi_i \rangle$ most likely

to emit $\langle x_1, \dots, x_i \rangle$ such that $\pi_i = k$

Then

$$p_k(i) = e_k(x_i) \max_r (p_r(i-1)t_{rk})$$

The Viterbi Algorithm

- Initialization ($i = 0$)

$$p_0(0) = 1, \quad p_k(0) = 0 \text{ for } k > 0$$

- Recursion ($i = 1, \dots, L$): For each state k

$$p_k(i) = e_k(x_i) \max_r (p_r(i-1)t_{rk})$$

- Termination:

$$\Pr(x, \pi^*) = \max_k (p_k(\text{Length})t_{k-1,k})$$

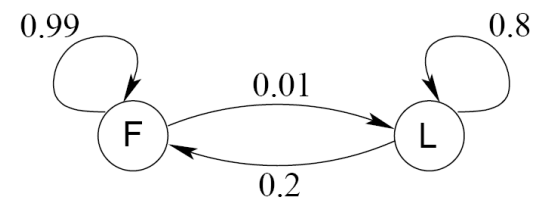
To find π^ , use trace-back, as in dynamic programming*

45

Viterbi: Example

		x		
		6	2	6
π	F	$(1/6) \times (1/2)$ $= 1/12$	$(1/6) \times \max\{(1/12) \times 0.99,$ $(1/4) \times 0.2\}$ $= 0.01375$	$(1/6) \times \max\{0.01375 \times 0.99,$ $0.02 \times 0.2\}$ $= 0.00226875$
	L	$(1/2) \times (1/2)$ $= 1/4$	$(1/10) \times \max\{(1/12) \times 0.01,$ $(1/4) \times 0.8\}$ $= 0.02$	$(1/2) \times \max\{0.01375 \times 0.01,$ $0.02 \times 0.8\}$ $= 0.08$

$$p_k(i) = e_k(x_i) \max_r (p_r(i-1) t_{rk})$$



Dynamic Bayesian Networks

- In addition to the discussed tasks, methods are needed for *learning* the transition and sensor models from observation.
- Learning can be done by inference, where inference provides an estimate of what transitions actually occurred and of what states generated the sensor readings. These estimates can be used to update the models.
- The updated models provide new estimates, and the process iterates to convergence.

DBN – Special Cases

- Hidden Markov Model (HMMs):

Temporal probabilistic model in which the state of the process is described by a single discrete random variable. (The simplest kind of DBN)

- Kalman Filter Models (KFMs):

Estimate the state (continuous) of a physical system from noisy observations over time. Also known as linear dynamical systems (LDSs).

Hidden Markov Models

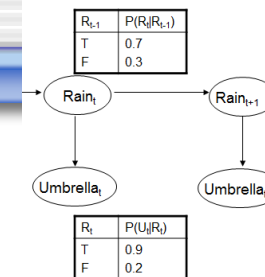
X_t is a single, discrete variable (usually E_t is too)

Domain of X_t is $\{1, \dots, S\}$

Transition matrix $\mathbf{T}_{ij} = P(X_t = j | X_{t-1} = i)$, e.g., $\begin{pmatrix} 0.7 & 0.3 \\ 0.3 & 0.7 \end{pmatrix}$

Sensor matrix \mathbf{O}_t for each time step, diagonal elements $P(e_t | X_t = i)$

e.g., with $U_1 = true$, $\mathbf{O}_1 = \begin{pmatrix} 0.9 & 0 \\ 0 & 0.2 \end{pmatrix}$



Hidden Markov Models

X_t is a single, discrete variable (usually E_t is too)
 Domain of X_t is $\{1, \dots, S\}$

Transition matrix $T_{ij} = P(X_t = j | X_{t-1} = i)$, e.g., $\begin{pmatrix} 0.7 & 0.3 \\ 0.3 & 0.7 \end{pmatrix}$

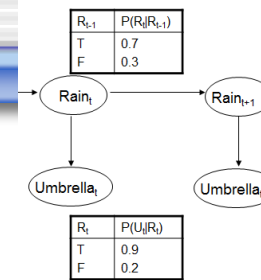
Sensor matrix O_t for each time step, diagonal elements $P(e_t | X_t = i)$
 e.g., with $U_1 = true$, $O_1 = \begin{pmatrix} 0.9 & 0 \\ 0 & 0.2 \end{pmatrix}$

Forward and backward messages as column vectors:

$$\mathbf{f}_{1:t+1} = \alpha \mathbf{O}_{t+1} \mathbf{T}^\top \mathbf{f}_{1:t}$$

$$\mathbf{b}_{k+1:t} = \mathbf{T} \mathbf{O}_{k+1} \mathbf{b}_{k+2:t}$$

Forward-backward algorithm needs time $O(S^2t)$ and space $O(St)$

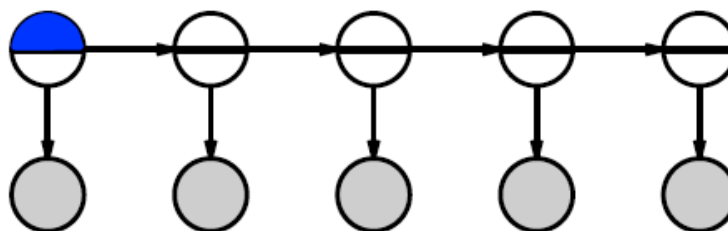


Country Dance Algorithm

Can avoid storing all forward messages in smoothing by running forward algorithm backwards:

$$\mathbf{f}_{1:t+1} = \alpha \mathbf{O}_{t+1} \mathbf{T}^\top \mathbf{f}_{1:t}$$

Algorithm: forward pass computes \mathbf{f}_t , backward pass does $\mathbf{f}_i, \mathbf{b}_i$

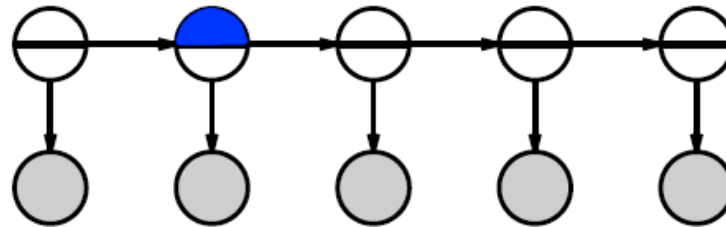


Country Dance Algorithm

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$$\mathbf{f}_{1:t+1} = \alpha \mathbf{O}_{t+1} \mathbf{T}^\top \mathbf{f}_{1:t}$$

Algorithm: forward pass computes \mathbf{f}_t , backward pass does $\mathbf{f}_i, \mathbf{b}_i$



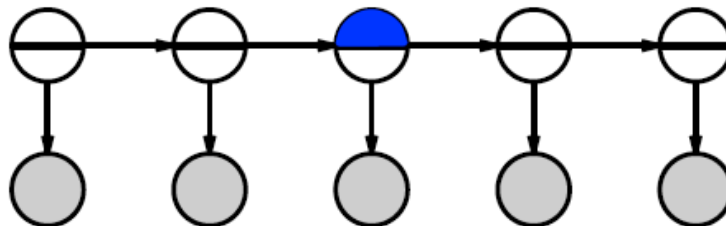
53

Country Dance Algorithm

Can avoid storing all forward messages in smoothing by running forward algorithm backwards:

$$\mathbf{f}_{1:t+1} = \alpha \mathbf{O}_{t+1} \mathbf{T}^\top \mathbf{f}_{1:t}$$

Algorithm: forward pass computes \mathbf{f}_t , backward pass does $\mathbf{f}_i, \mathbf{b}_i$

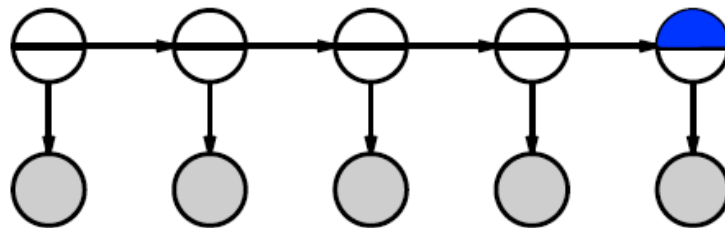


Country Dance Algorithm

Can avoid storing all forward messages in smoothing by running forward algorithm backwards:

$$\mathbf{f}_{1:t+1} = \alpha \mathbf{O}_{t+1} \mathbf{T}^\top \mathbf{f}_{1:t}$$

Algorithm: forward pass computes \mathbf{f}_t , backward pass does $\mathbf{f}_i, \mathbf{b}_i$



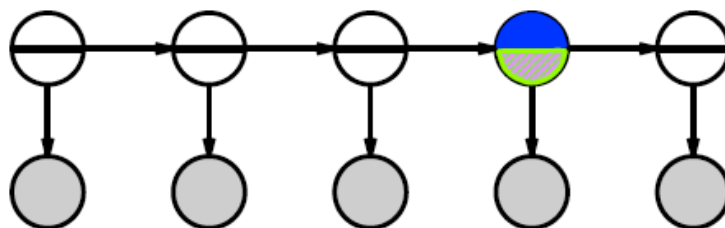
55

Country Dance Algorithm

Can avoid storing all forward messages in smoothing by running forward algorithm backwards:

$$\begin{aligned} \mathbf{f}_{1:t+1} &= \alpha \mathbf{O}_{t+1} \mathbf{T}^\top \mathbf{f}_{1:t} \\ \mathbf{O}_{t+1}^{-1} \mathbf{f}_{1:t+1} &= \alpha \mathbf{T}^\top \mathbf{f}_{1:t} \\ \alpha' (\mathbf{T}^\top)^{-1} \mathbf{O}_{t+1}^{-1} \mathbf{f}_{1:t+1} &= \mathbf{f}_{1:t} \end{aligned}$$

Algorithm: forward pass computes \mathbf{f}_t , backward pass does $\mathbf{f}_i, \mathbf{b}_i$

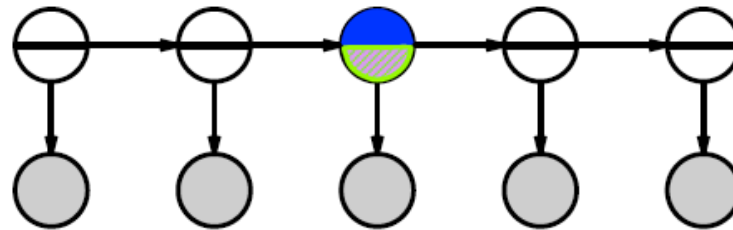


Country Dance Algorithm

Can avoid storing all forward messages in smoothing by running forward algorithm backwards:

$$\begin{aligned} \mathbf{f}_{1:t+1} &= \alpha \mathbf{O}_{t+1} \mathbf{T}^\top \mathbf{f}_{1:t} \\ \mathbf{O}_{t+1}^{-1} \mathbf{f}_{1:t+1} &= \alpha \mathbf{T}^\top \mathbf{f}_{1:t} \\ \alpha' (\mathbf{T}^\top)^{-1} \mathbf{O}_{t+1}^{-1} \mathbf{f}_{1:t+1} &= \mathbf{f}_{1:t} \end{aligned}$$

Algorithm: forward pass computes \mathbf{f}_t , backward pass does $\mathbf{f}_i, \mathbf{b}_i$



Applications

- Speech recognition
- Robot localization
- ...

○	○	○	○
■	■	○	○
■	○	○	○
○	○	■	○

One non-deterministic operation MOVE.

$$P(X_{t+1} = j | X_t = i) = \mathbf{T}_{ij} = (1/N(i) \text{ if } j \in \text{NEIGHBORS}(i) \text{ else } 0)$$

E_t has 16 possible values, each a four-bit sequence giving the presence or absence of an obstacle: NSWE.

ϵ is the error rate. All four bits right $(1 - \epsilon)^4$. All wrong ϵ^4 .

d_{it} is the number of bits that are different between the true values for square i and the actual reading e_t , then the probability that a robot in square i would receive a sensor reading e_t is:

$$P(E_t = e_t | X_t = i) = \mathbf{O}_{t_{ii}} = (1 - \epsilon)^{4-d_{it}} \epsilon^{d_{it}}$$

○	○	○	○
■	■	○	○
■	○	○	○
○	○	■	○

Cell numbers: start in top row, left to right

Matrix for **NSW**

	1	2	3	...	12
1	$(1-\varepsilon)^4$				
2		$(1-\varepsilon)^3 \varepsilon$			
3			$(1-\varepsilon)^2 \varepsilon^2$		
...				...	
12					$(1-\varepsilon)^2 \varepsilon^2$

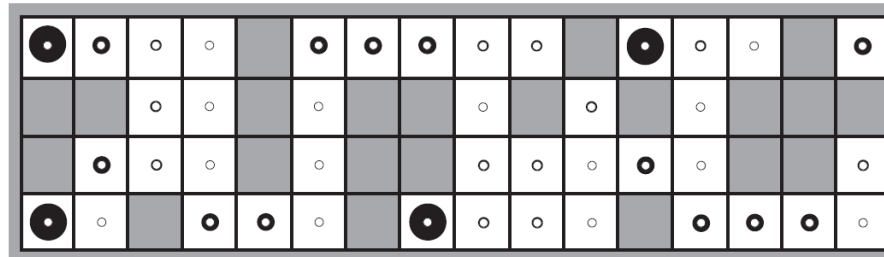
●	○	○	○
■	■	○	○
■	●	○	○
●	○	■	○

$E_1 = \text{NSW}$

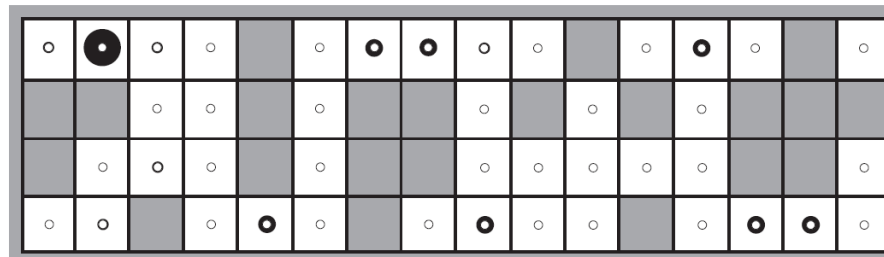
○	●	○	○
■	■	○	○
■	○	○	○
○	○	■	○

$E_2 = \text{NS}$

Example AIMA



(a) Posterior distribution over robot location after $E_1 = \text{NSW}$



(b) Posterior distribution over robot location after $E_1 = \text{NSW}, E_2 = \text{NS}$

Performance

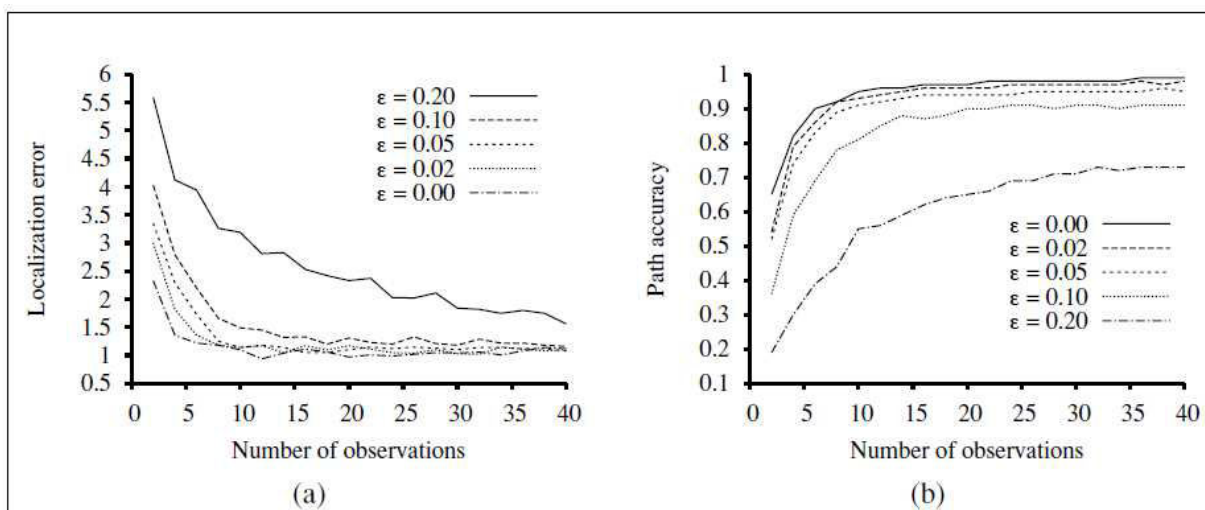


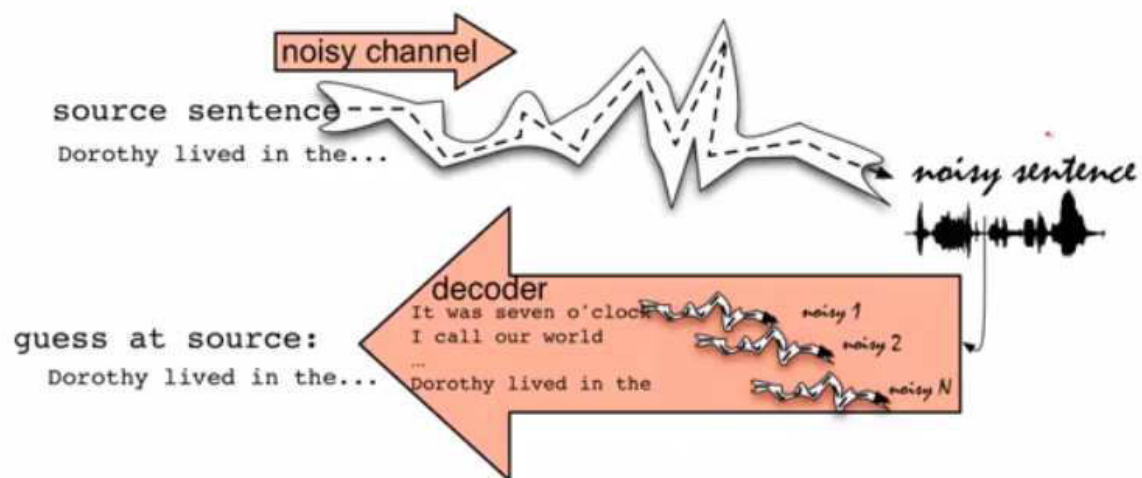
Figure 15.8 Performance of HMM localization as a function of the length of the observation sequence for various different values of the sensor error probability ϵ ; data averaged over 400 runs. (a) The localization error, defined as the Manhattan distance from the true location. (b) The Viterbi path accuracy, defined as the fraction of correct states on the Viterbi path.

Last time

- Filtering
- Prediction
- Smoothing
- Viterbi for *most likely path/state sequence* for given observation
- HMM
 - Only one state variable
 - Efficient computation because of matrix operations

$$\begin{aligned} \mathbf{f}_{1:t+1} &= \alpha \mathbf{O}_{t+1} \mathbf{T}^\top \mathbf{f}_{1:t} \\ \mathbf{O}_{t+1}^{-1} \mathbf{f}_{1:t+1} &= \alpha \mathbf{T}^\top \mathbf{f}_{1:t} \\ \alpha' (\mathbf{T}^\top)^{-1} \mathbf{O}_{t+1}^{-1} \mathbf{f}_{1:t+1} &= \mathbf{f}_{1:t} \end{aligned}$$

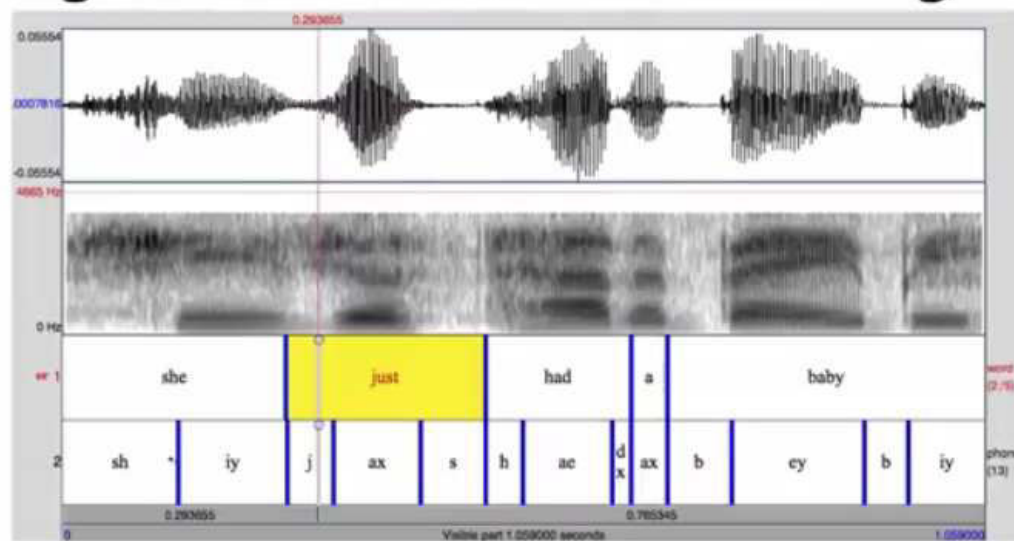
Speech recognition



Dan Jurafsky, Stanford

Daphne Koller

Segmentation of Acoustic signals



Dan Jurafsky, Stanford

Daphne Keller

Phonetic alphabet

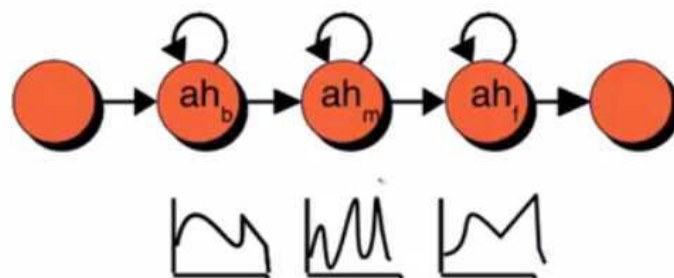
• AA	odd	AAD	• G	green	GRIYN	• R	read	RIYD
• AE	at	AET	• HH	he	HHIY	• S	sea	SIY
• AH	hut	HH AHT	• IH	it	IHT	• SH	she	SHIY
• AO	ought	AOT	• IY	eat	IYT	• T	tea	TIY
• AW	cow	KAW	• JH	gee	JHIY	• TH	theta	THEY T AH
• AY	hide	HH AYD	• K	key	KIY	• UH	hood	HH UHD
• B	be	BIY	• L	lee	LIY	• UW	two	TUW
• CH	cheese	CH IY Z	• M	me	MIY	• V	vee	VIY
• D	dee	DIY	• N	knee	NIY	• W	we	WIY
• DH	thee	DHIY	• NG	ping	PIHNG	• Y	yield	YIYL D
• EH	Ed	EHD	• OW	oat	OWT	• Z	zee	ZIY
• ER	hurt	HHERT	• OY	toy	TOY	• ZH	seizure	SIY ZH ER
• EY	ate	EYT	• P	pee	PIY			
• F	fee	FIY						

<http://www.speech.cs.cmu.edu/cgi-bin/cmudict>



The CMU Pronouncing Dictionary

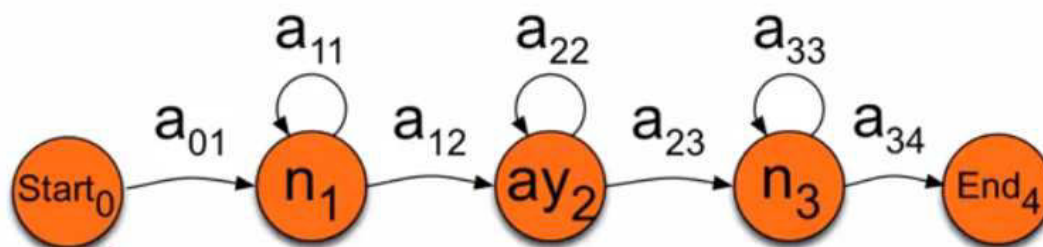
HMM ah



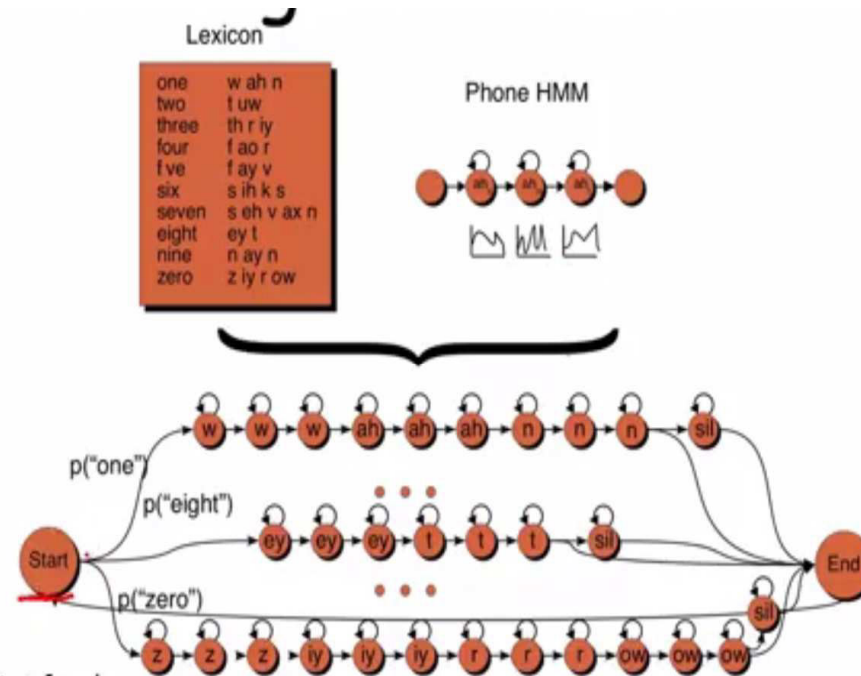
Dan Jurafsky, Stanford

Daphne Koller

Word HMM: nine



Recognition HMM



Dan Jurafsky, Stanford

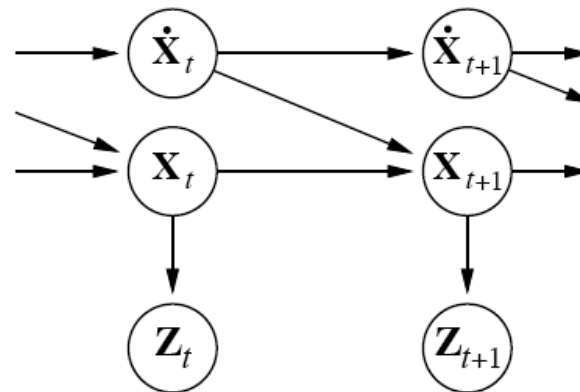
Daphne Koller

Kalman Filters

Modelling systems described by a set of continuous variables,

e.g., tracking a bird flying— $\mathbf{X}_t = X, Y, Z, \dot{X}, \dot{Y}, \dot{Z}$.

Airplanes, robots, ecosystems, economies, chemical plants, planets, ...



Gaussian prior, linear Gaussian transition model and sensor model

Updating Gaussian Distributions

Prediction step: if $\mathbf{P}(\mathbf{X}_t|\mathbf{e}_{1:t})$ is Gaussian, then prediction

$$\mathbf{P}(\mathbf{X}_{t+1}|\mathbf{e}_{1:t}) = \int_{\mathbf{x}_t} \mathbf{P}(\mathbf{X}_{t+1}|\mathbf{x}_t)P(\mathbf{x}_t|\mathbf{e}_{1:t}) d\mathbf{x}_t$$

is Gaussian. If $\mathbf{P}(\mathbf{X}_{t+1}|\mathbf{e}_{1:t})$ is Gaussian, then the updated distribution

$$\mathbf{P}(\mathbf{X}_{t+1}|\mathbf{e}_{1:t+1}) = \alpha\mathbf{P}(\mathbf{e}_{t+1}|\mathbf{X}_{t+1})\mathbf{P}(\mathbf{X}_{t+1}|\mathbf{e}_{1:t})$$

is Gaussian

Hence $\mathbf{P}(\mathbf{X}_t|\mathbf{e}_{1:t})$ is multivariate Gaussian $N(\boldsymbol{\mu}_t, \boldsymbol{\Sigma}_t)$ for all t

Simple 1-D Example

Prior $P(x_0) = \alpha e^{-\frac{1}{2} \left(\frac{(x_0 - \mu_0)^2}{\sigma_0^2} \right)}$

Transition model $P(x_{t+1}|x_t) = \alpha e^{-\frac{1}{2} \left(\frac{(x_{t+1} - x_t)^2}{\sigma_x^2} \right)}$

Sensor model $P(z_t|x_t) = \alpha e^{-\frac{1}{2} \left(\frac{(z_t - x_t)^2}{\sigma_z^2} \right)}$

Prediction
$$P(x_1) = \int_{-\infty}^{\infty} P(x_1|x_0)P(x_0) dx_0 = \alpha \int_{-\infty}^{\infty} e^{-\frac{1}{2} \left(\frac{(x_1 - x_0)^2}{\sigma_x^2} \right)} e^{-\frac{1}{2} \left(\frac{(x_0 - \mu_0)^2}{\sigma_0^2} \right)} dx_0$$

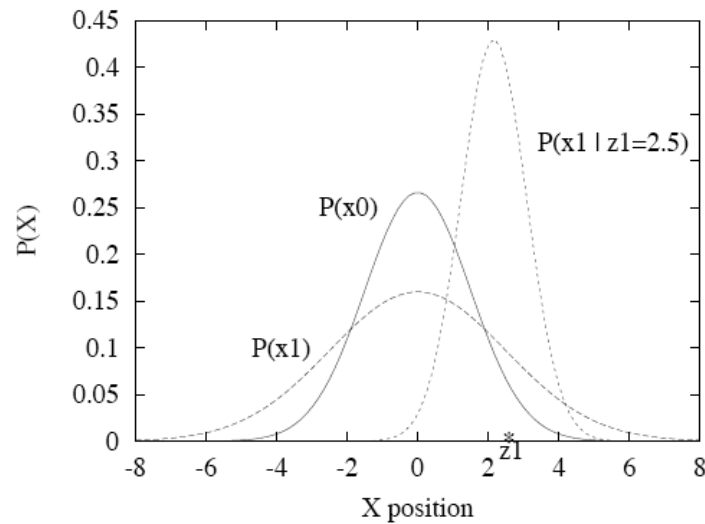
$$= \alpha \int_{-\infty}^{\infty} e^{-\frac{1}{2} \left(\frac{\sigma_0^2(x_1 - x_0)^2 + \sigma_x^2(x_0 - \mu_0)^2}{\sigma_0^2 \sigma_x^2} \right)} dx_0$$

$$= \alpha e^{-\frac{1}{2} \left(\frac{(x_1 - \mu_0)^2}{\sigma_0^2 + \sigma_x^2} \right)} \quad \text{(by using completing the square. Not discussed here)}$$

Simple 1-D Example

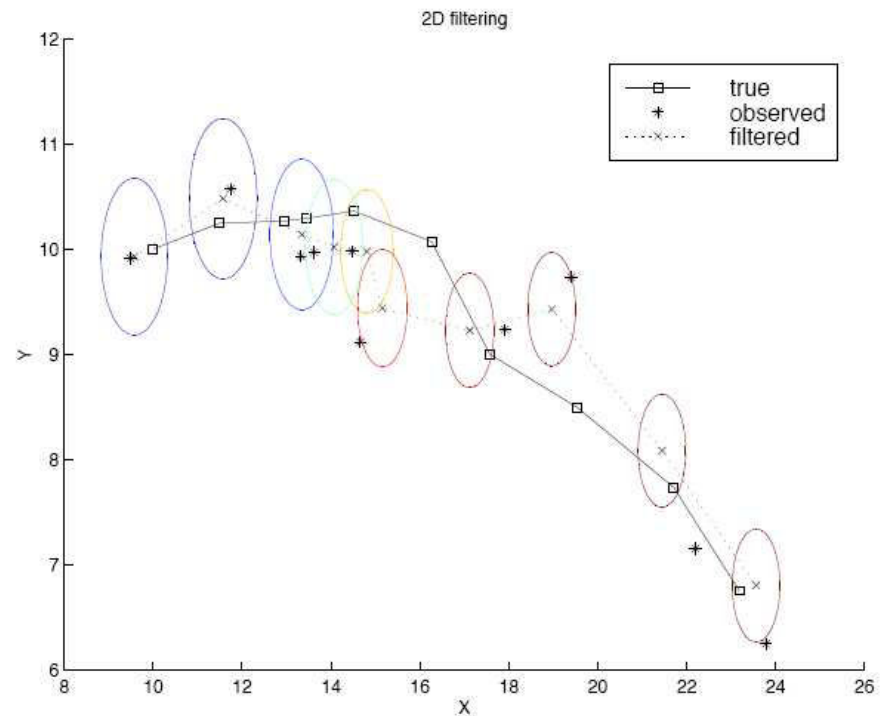
Gaussian random walk on X -axis, s.d. σ_x , sensor s.d. σ_z

$$\mu_{t+1} = \frac{(\sigma_t^2 + \sigma_x^2)z_{t+1} + \sigma_z^2\mu_t}{\sigma_t^2 + \sigma_x^2 + \sigma_z^2} \quad \sigma_{t+1}^2 = \frac{(\sigma_t^2 + \sigma_x^2)\sigma_z^2}{\sigma_t^2 + \sigma_x^2 + \sigma_z^2}$$

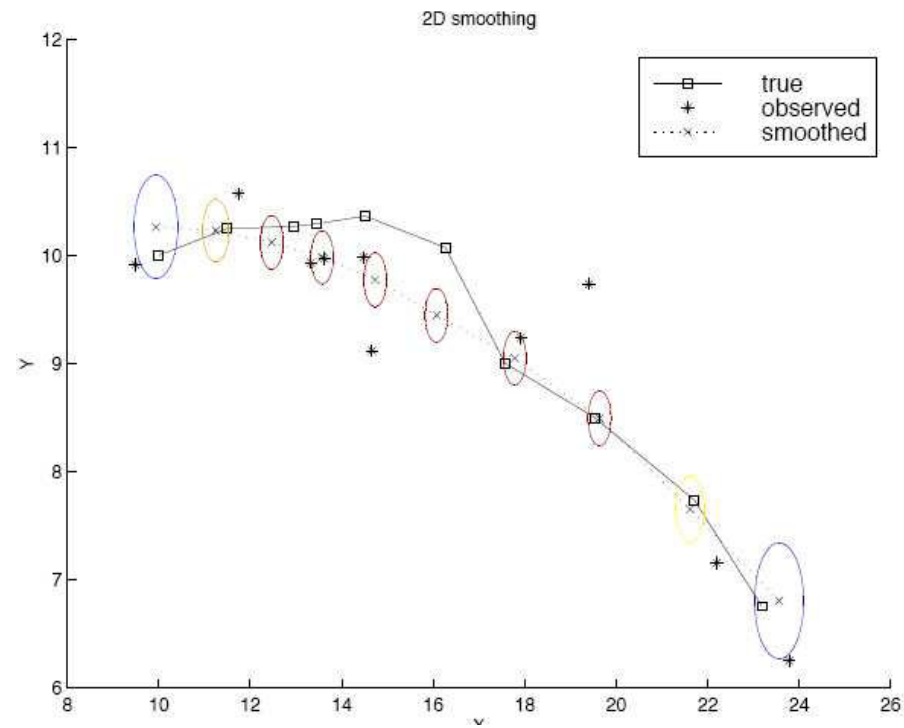


75

2-D Tracking: Filtering

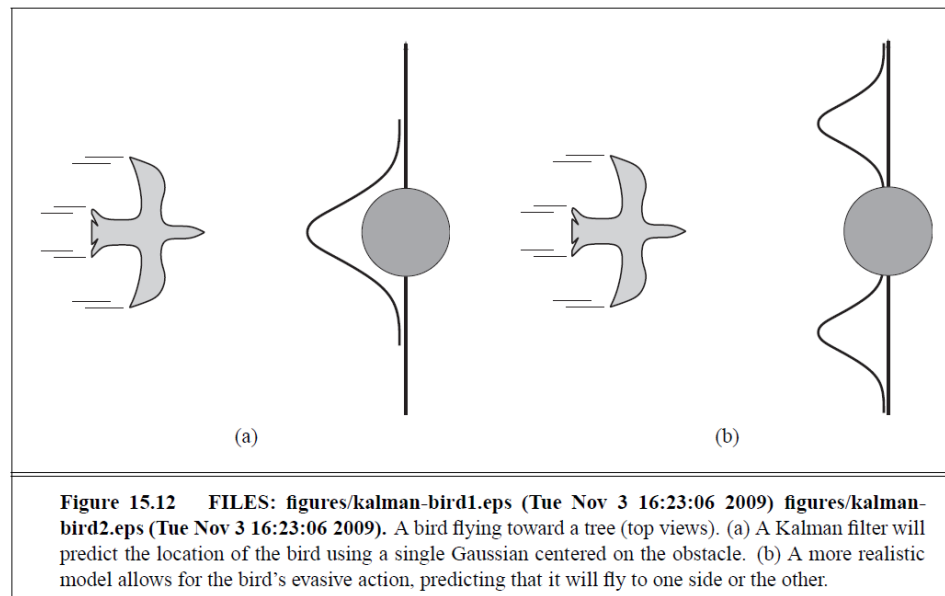


2-D Tracking: Smoothing



78

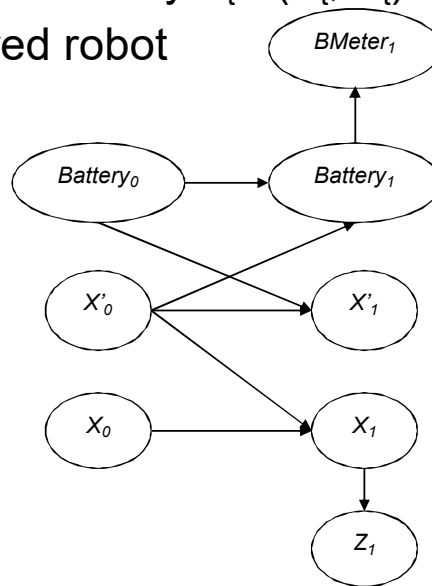
Where it breaks



One solution → switching kalman filters

Creating DBNs with failures

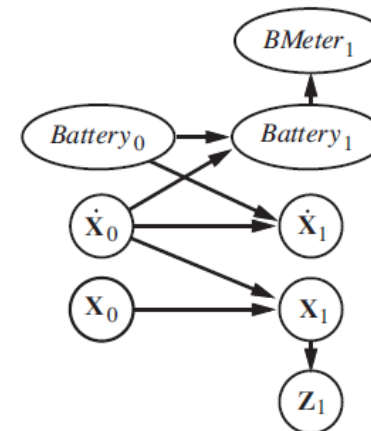
- $\mathbf{X}'_t = (X_t, Y_t)$ for velocity $\mathbf{X}_t = (X_t, Y_t)$ for position
- Battery powered robot



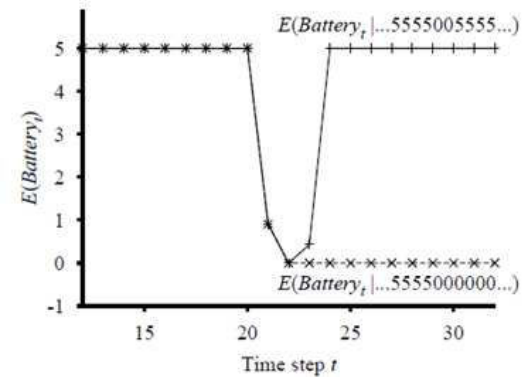
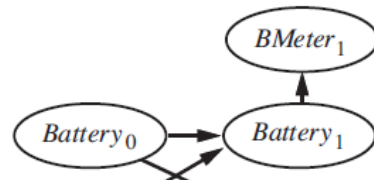
81

Failure of sensors

- Sensor measurements are noisy
- Real sensors can fail
- May use a Gaussian error model for *discrete variables*
- Transient failure
- Persistent failure



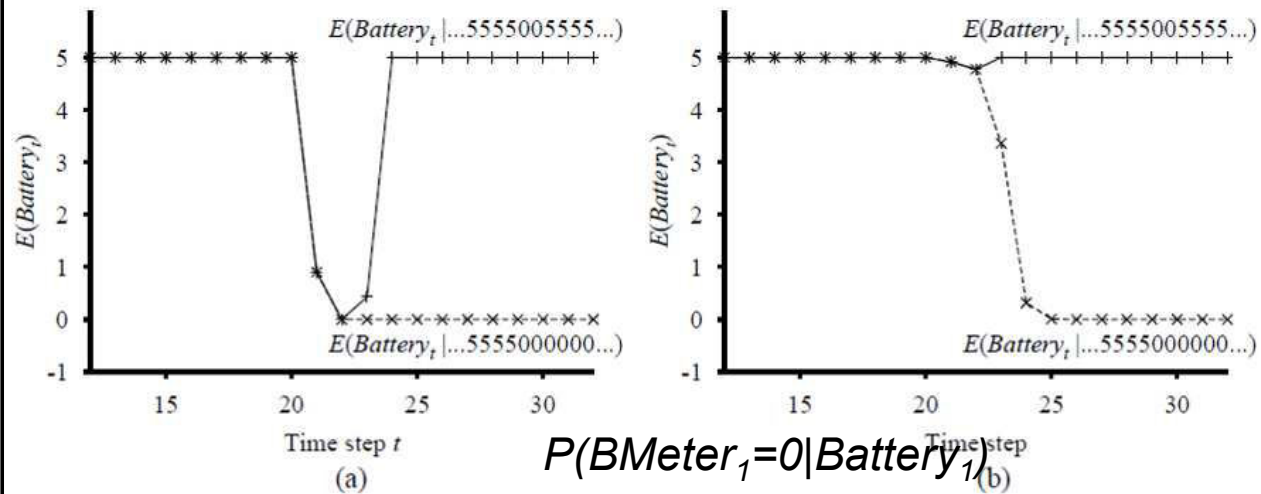
Transient failure model



$$\begin{aligned}
 & P(BMeter_1=0|Battery_1) P(Battery_1) \\
 P(Battery_1|BMeter_1=0) &= \alpha \langle 0.99, 0.006, 0.004 \rangle \langle 0.05, 0.05, 0.9 \rangle \\
 &= \langle 0,92178771, 0,01117318, 0,06703911 \rangle
 \end{aligned}$$

83

Transient failure model



$$P(\text{Battery}_1 | \text{BMeter}_1=0) = \alpha \langle 0.8, 0.1, 0.1 \rangle \langle 0.05, 0.05, 0.9 \rangle$$

$$\approx \langle 0.44, 0.06, 0.5 \rangle$$

Persistent failure model

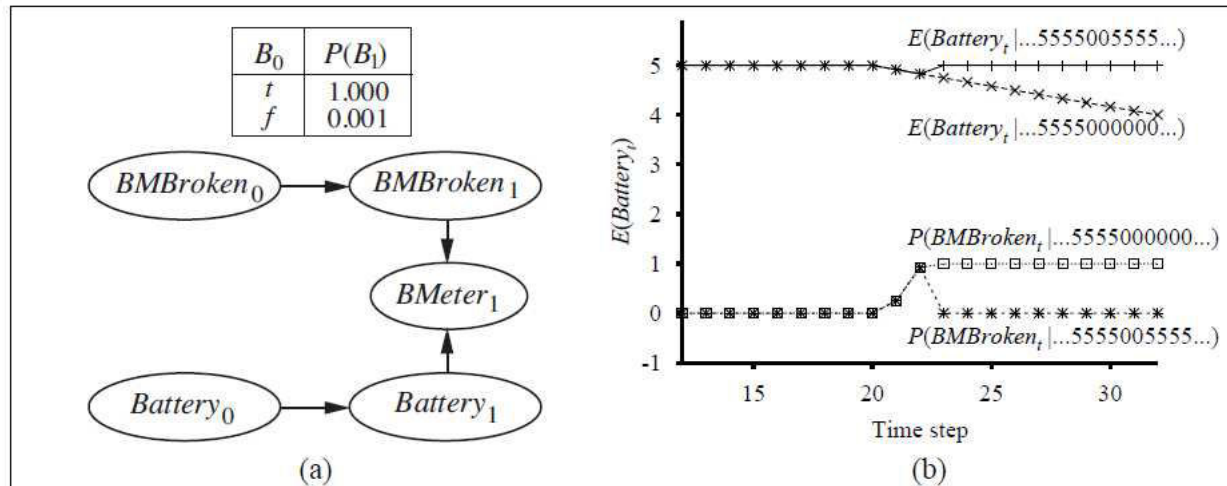
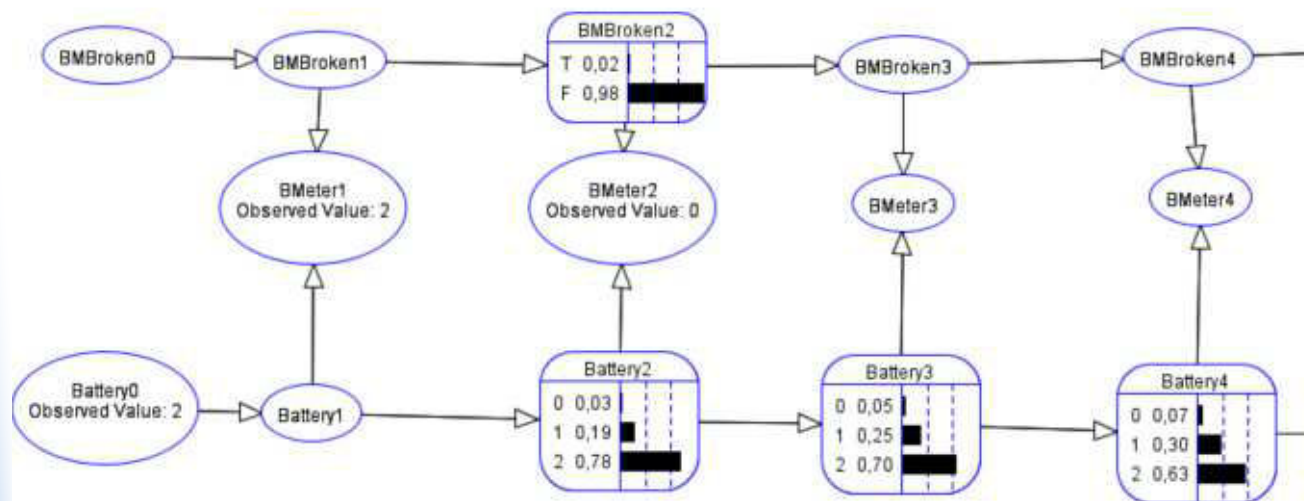


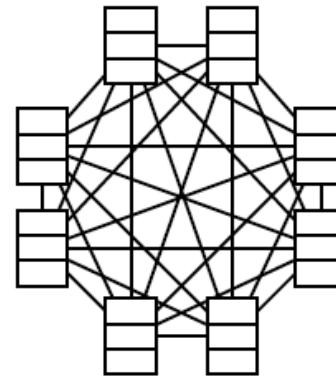
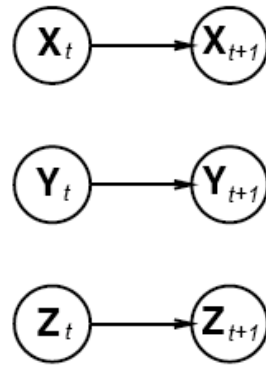
Figure 15.15 FILES: figures/battery-persistence.eps (Tue Nov 3 16:22:26 2009). (a) A DBN fragment showing the sensor status variable required for modeling persistent failure of the battery sensor. (b) Upper curves: trajectories of the expected value of $Battery_t$ for the “transient failure” and “permanent failure” observations sequences. Lower curves: probability trajectories for $BMBroken$ given the two observation sequences.

Example



DBNs vs. HMMs

Every HMM is a single-variable DBN; every discrete DBN is an HMM



Consider the transition model

Sparse dependencies \Rightarrow exponentially fewer parameters;

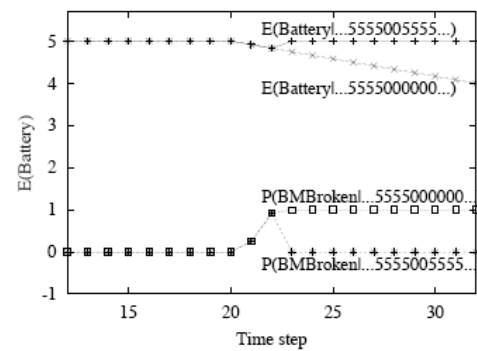
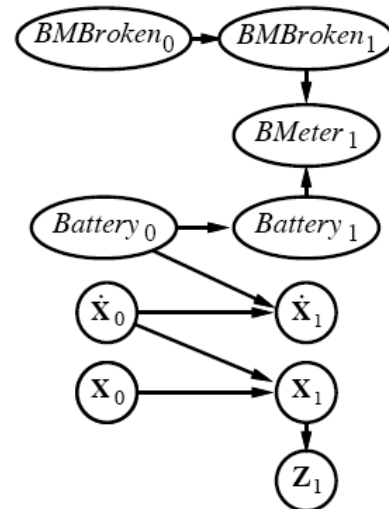
e.g., 20 state variables, three parents each

DBN has $20 \times 2^3 = 160$ parameters, HMM has $2^{20} \times 2^{20} \approx 10^{12}$

DBNs vs. Kalman Filters

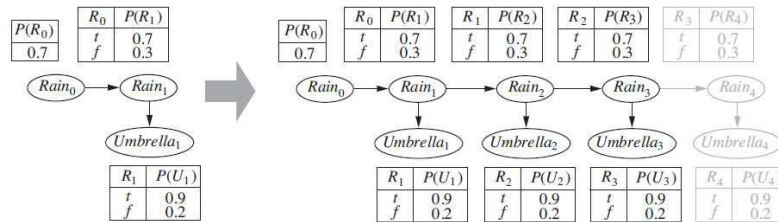
Every Kalman filter model is a DBN, but few DBNs are KFs;
real world requires non-Gaussian posteriors

E.g., where are bin Laden and my keys? What's the battery charge?

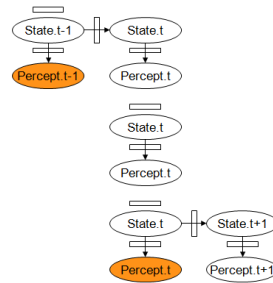


Exact Inference in DBNs

Naive:



Rollup filtering:



$O(d^{n+k})$ largest factor

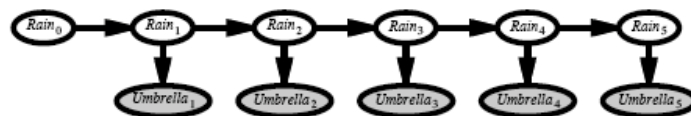
20 state variables (4 values)
mean $4^{20+1} = 4.398.046.511.104$

d = possible values for variables
 n = number of states
 k = number of parents

90

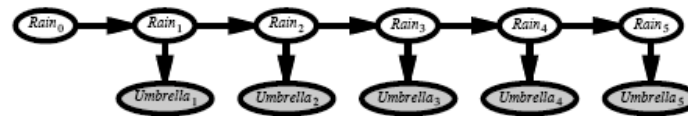
Approximate inference: Likelihood Weighting

Set of weighted samples approximates the belief state



Likelihood Weighting

Set of weighted samples approximates the belief state

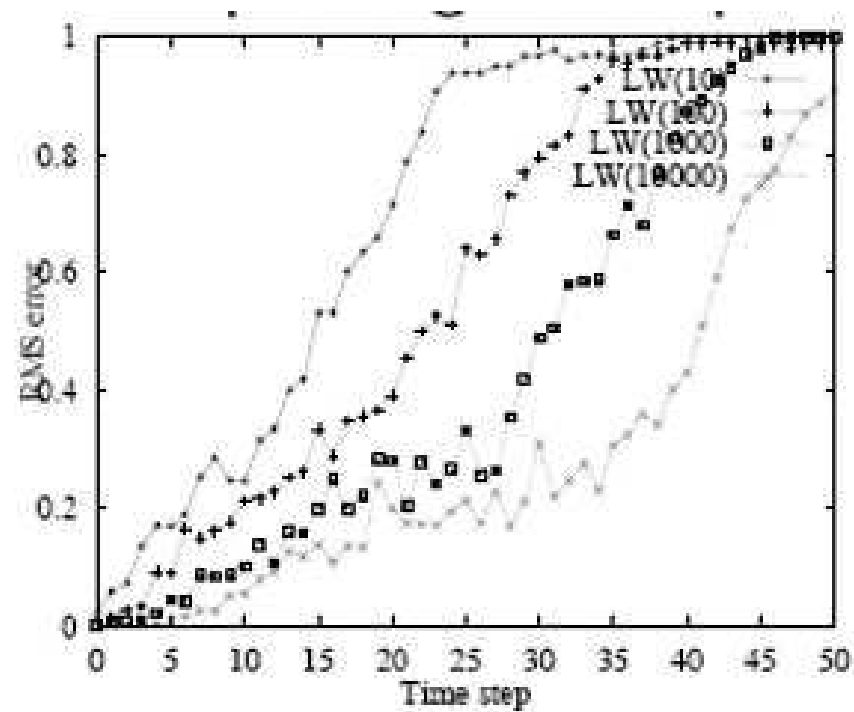


LW samples pay no attention to the evidence!

⇒ fraction “agreeing” falls exponentially with t

⇒ number of samples required grows exponentially with t

Likelihood Weighting



Solution

- Instead of running one example at a time run N .
 - ◆ The N samples also represent an approximate representation of the current state distribution.
- Instead of using initial examples throw low weighted ones away.
 - ◆ Must add new examples else lose to much.

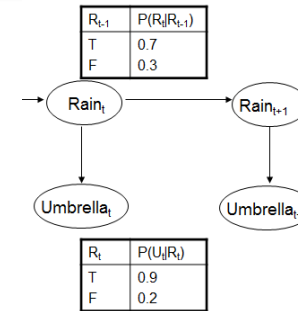
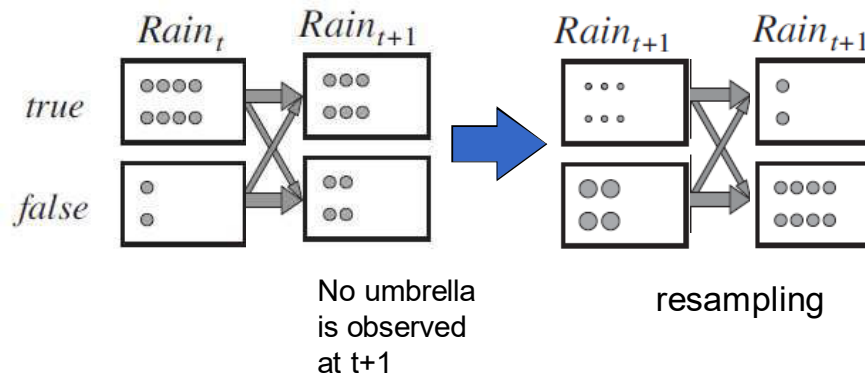
Idea: Particle filtering

A population of N initial-state samples is created sampling from $P(X_0)$

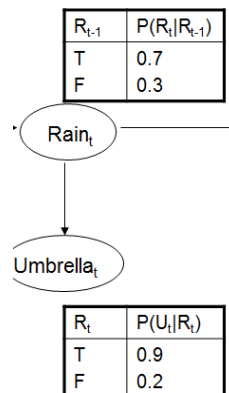
1. Based on the transition matrix propagate examples forward. $P(X_{t+1}|x_t)$
2. Each sample is weighted by the likelihood it assigns to the new evidence $P(e_{t+1}|x_{t+1})$.
3. **Resample examples based on it's weight.**

Particle Filtering

Our current particles, 10 Propagate forward



Example



$$N(r_{t+1}|e) = \sum_{x_t} P(x_{t+1}|x_t) N(x_t|e)$$

$$\text{For rain} = 0.7 * 8 + 0.3 * 2 = 6.2 \Rightarrow 6$$

$$\text{For not rain} = 0.3 * 8 + 0.7 * 2 = 3.8 \Rightarrow 4$$

Suppose no umbrella for t+1

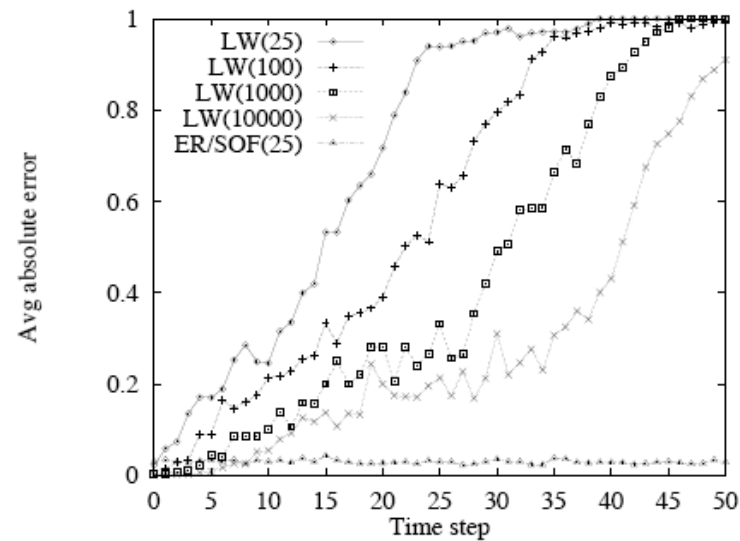
$$\text{total weight(rain particles)} = 0.1 * 6 = 0.6$$

$$\text{total weight(not rain)} = 0.8 * 4 = 3.2$$

$$\text{Normalized} = \langle 0.17, 0.83 \rangle$$

Particle Filtering: Performance

Approximation error of particle filtering remains bounded over time, at least empirically—theoretical analysis is difficult



99

Summary

Temporal models use state and sensor variables replicated over time

Markov assumptions and stationarity assumption, so we need

- transition model $P(\mathbf{X}_t | \mathbf{X}_{t-1})$
- sensor model $P(\mathbf{E}_t | \mathbf{X}_t)$

Tasks are filtering, prediction, smoothing, most likely sequence;
all done recursively with constant cost per time step

Hidden Markov models have a single discrete state variable; used for speech recognition

Kalman filters allow n state variables, linear Gaussian, $O(n^3)$ update

Dynamic Bayes nets subsume HMMs, Kalman filters; exact update intractable

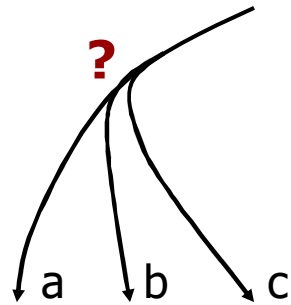
Particle filtering is a good approximate filtering algorithm for DBNs

Intelligent Autonomous Agents and Cognitive Robotics

Topic 7: Decision-Making under Uncertainty Simple Decisions

Ralf Möller, Rainer Marrone
Hamburg University of Technology

Non-Deterministic vs. Probabilistic Uncertainty

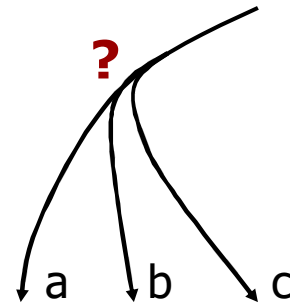


$\{a,b,c\}$

→ decision that is
best for worst case

Non-deterministic model

~ Adversarial search



$\{a(p_a),b(p_b),c(p_c)\}$

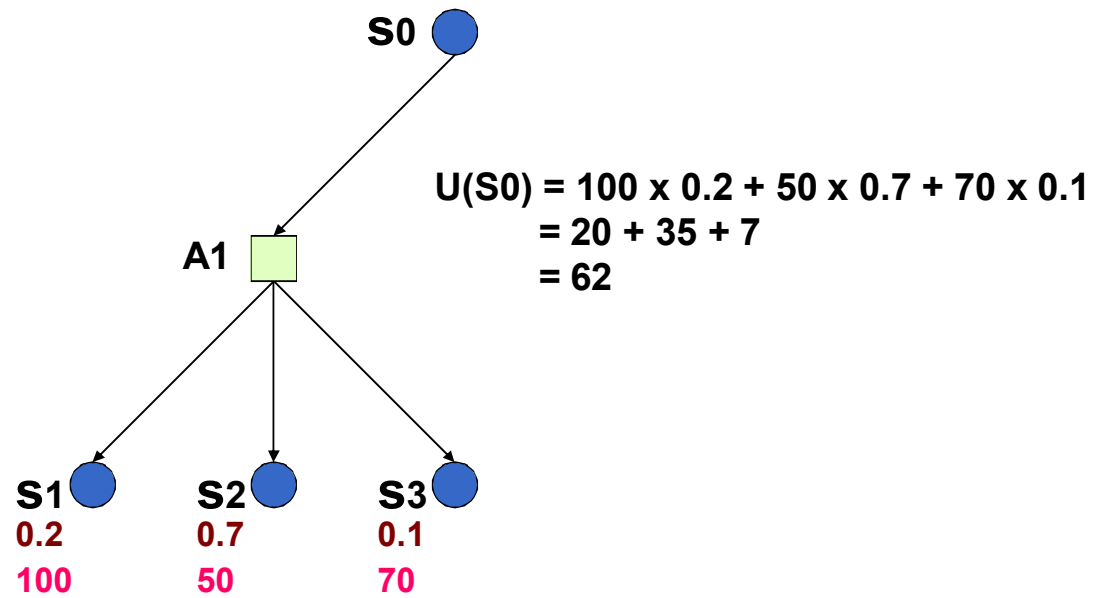
→ decision that maximizes
expected utility value

Probabilistic model

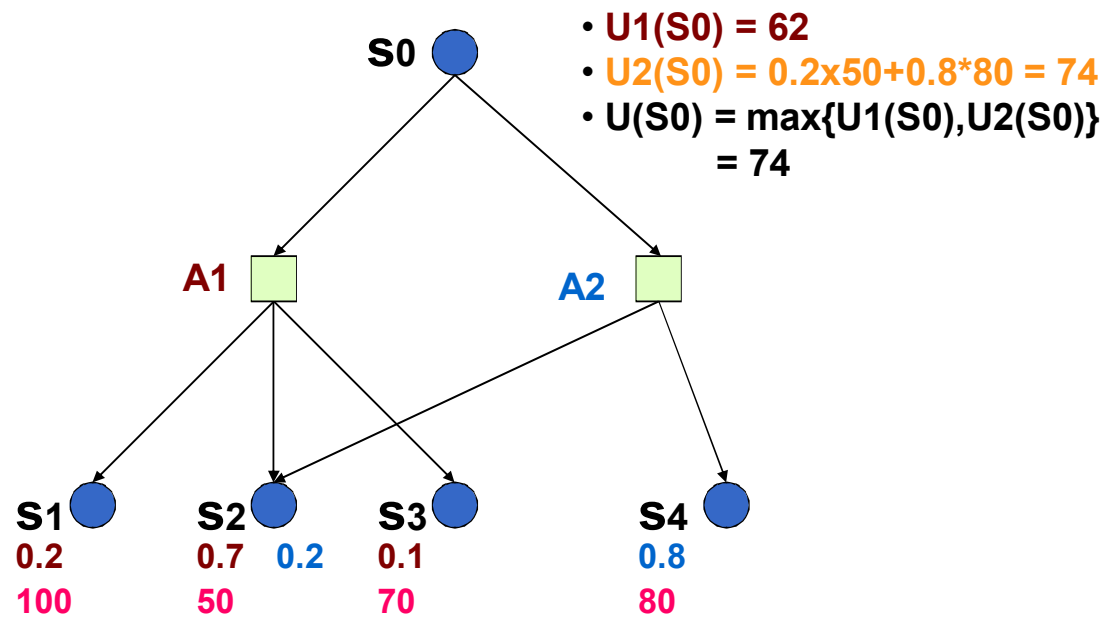
Expected Utility

- Random variable X with n values x_1, \dots, x_n and distribution (p_1, \dots, p_n)
 \mathbf{X} is the state reached after doing an action \mathbf{A} under uncertainty
- Function \mathbf{U} of \mathbf{X} : \mathbf{U} is the utility of a state
- The **expected utility** of \mathbf{A} is
$$EU[A] = \sum_{i=1, \dots, n} p(x_i | A) U(x_i)$$
$$\mathbf{MEU} = \operatorname{argmax}_A EU[A]$$

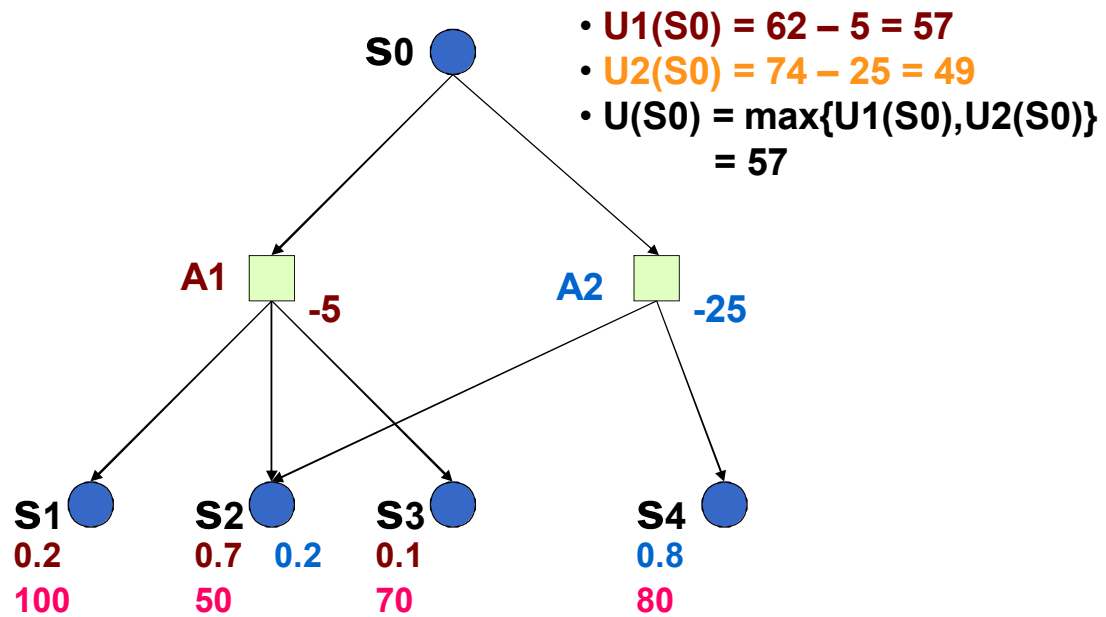
One State/One Action Example



One State/Two Actions Example



Introducing Action Costs



MEU Principle

- A **rational agent** should choose the action that maximizes agent's expected utility
- This is the basis of the field of **decision theory**
- The MEU principle provides a **normative criterion** for rational choice of action

But ...

- Must have **complete** model of:
 - ♦ Actions
 - ♦ Utilities
 - ♦ States
- Even if you have a complete model, it might be computationally **intractable**
- In fact, a truly rational agent takes into account the utility of reasoning as well---**bounded rationality**
- Nevertheless, great progress has been made in this area recently, and we are able to solve much more complex decision-theoretic problems than ever before

We'll look at

- Decision-Theoretic Planning
 - ◆ Simple decision making (ch. 16)
 - ◆ Sequential decision making (ch. 17)

Rational preferences

Idea: preferences of a rational agent must obey constraints.

Rational preferences →

behavior describable as maximization of expected utility

MEU is not the only possible solution:

minimize worst case

only preferences without numeric values

...

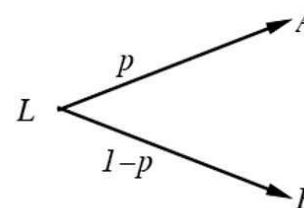
Why should a utility function with numerical values exist?

10

Basis of utility theory: constraints on preferences

- An agent chooses among prizes (A, B, \dots) and lotteries, i.e., situations with uncertain prizes.

Lottery $L = [p, A ; (1-p), B]$



$A \succ B$ the agent prefers A over B .

$A \sim B$ the agent is indifferent between A and B .

A and B can be lotteries again: Prizes are special lotteries: $[1, X; 0, \text{not } X]$

Axioms of Utility Theory

- **Orderability:** Given any two states, the rational agent prefers one of them, else the two as equally preferable.

$$(A \succ B) \vee (B \succ A) \vee (A \sim B)$$

- **Transitivity:** Given any three states, if an agent prefers A to B and prefers B to C , agent must prefer A to C .

$$(A \succ B) \wedge (B \succ C) \Rightarrow (A \succ C)$$

- **Continuity:** If some state B is between A and C in preference, then there is a p for which the rational agent will be indifferent between state B and the lottery in which A comes with probability p , C with probability $(1-p)$.

$$(A \succ B \succ C) \Rightarrow \exists p [p : A; (1-p) : C] \sim B$$

Rational preferences contd.

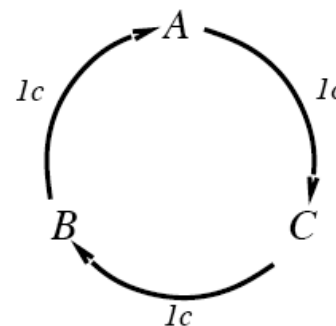
Violating the constraints leads to self-evident irrationality

For example: an agent with intransitive preferences can be induced to give away all its money

If $B \succ C$, then an agent who has C would pay (say) 1 cent to get B

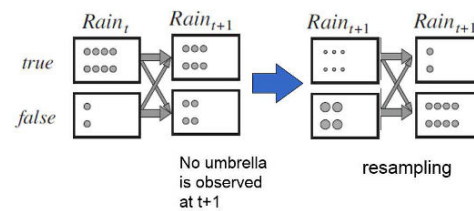
If $A \succ B$, then an agent who has B would pay (say) 1 cent to get A

If $C \succ A$, then an agent who has A would pay (say) 1 cent to get C



Last time

- Kalman filters
- Failure models for DBN: transient, persistent
- Approximate inference in DBNs: Particle filtering



- Utility theory
 - ♦ Lotteries and axioms for preferences

Axioms of Utility Theory

- **Substitutability:** If an agent is indifferent between two lotteries, A and B , then there is a more complex lottery in which A can be substituted with B . This also holds for \succ

$$(A \sim B) \Rightarrow [p : A; (1-p) : C] \sim [p : B; (1-p) : C]$$

- **Monotonicity:** If an agent prefers A to B , then the agent must prefer the lottery in which A occurs with a higher probability

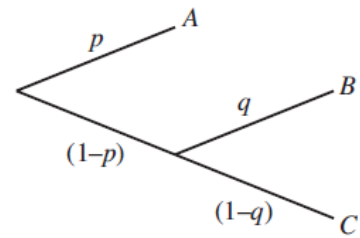
$$(A \succ B) \Rightarrow (p > q \Leftrightarrow [p : A; (1-p) : B] \succ [q : A; (1-q) : B])$$

- **Decomposability:** Compound lotteries can be reduced to simpler lotteries using the laws of probability.

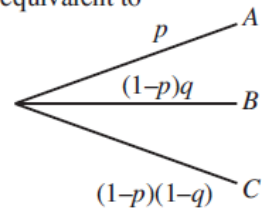
$$[p : A; (1-p) : [q : B; (1-q) : C]] \Rightarrow [p : A; (1-p)q : B; (1-p)(1-q) : C]$$

No fun in gambling

Decomposabilty



is equivalent to



And then there was utility

- Theorem by Neumann and Morgenstern, 1944
Given preferences satisfying the constraints there exists a real-valued function U such that

$$U(A) \geq U(B) \Leftrightarrow A \succsim B$$
$$U([p_1, S_1; \dots; p_n, S_n]) = \sum_i p_i U(S_i)$$

MEU principle:

Choose the action that maximizes expected utility

Allais Paradox

A : 80% chance of \$4000
B : 100% chance of \$3000

When presented with a choice between A and B, most people would choose the sure thing B.

C : 20% chance of \$4000
D : 25% chance of \$3000

When presented with a choice between C and D, most people would choose the C, with higher expected utility (800 vs. 750).

These choices together are inconsistent

$$1 \cdot U(3000) > 0.8 \cdot U(4000)$$

$$0.25 \cdot U(3000) < 0.2 \cdot U(4000)$$
$$1 \cdot U(3000) < 0.8 \cdot U(4000)$$

Utilities

Utilities map states to real numbers. Which numbers?

Standard approach to assessment of human utilities:

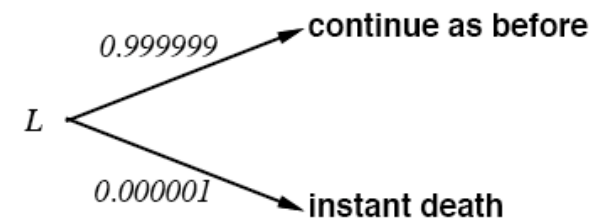
compare a given state A to a standard lottery L_p that has

“best possible prize” u_{\top} with probability p

“worst possible catastrophe” u_{\perp} with probability $(1 - p)$

adjust lottery probability p until $A \sim L_p$

pay \$30 \sim
-and-continue
-as-before



Utility scales

Normalized utilities: $u_{\top} = 1.0$, $u_{\perp} = 0.0$ $U(\text{pay } \$30\dots) = 0.999999$

Micromorts: one-millionth chance of death
useful for Russian roulette, paying to reduce product risks, etc.

QALYs: quality-adjusted life years
useful for medical decisions involving substantial risk

Note: behavior is **invariant** w.r.t. +ve linear transformation

$$U'(x) = k_1 U(x) + k_2 \quad \text{where } k_1 > 0$$

With deterministic prizes only (no lottery choices), only **ordinal utility** can be determined, i.e., total order on prizes

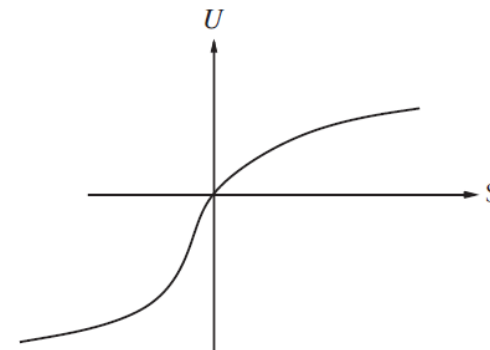
20

Value Functions

- Provides a ranking of alternatives, but not a meaningful metric scale
- Also known as an “ordinal utility function”
- Remember the expectiminimax example:
 - ♦ Sometimes, only relative judgments (value functions) are necessary
 - ♦ At other times, absolute judgments (utility functions) are required

Money Versus Utility

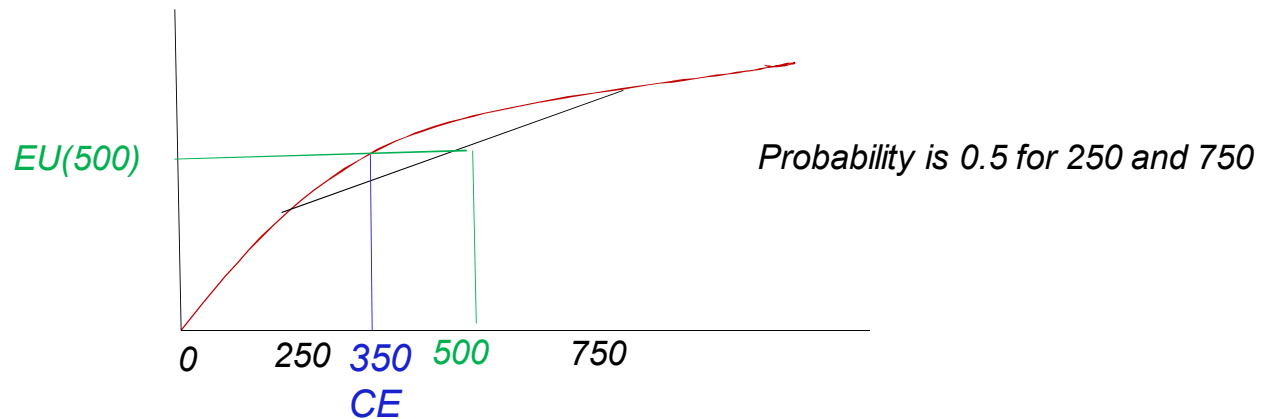
- Money \leftrightarrow Utility
 - ♦ More money is better, but not always in a linear relationship to the amount of money
- Expected Monetary Value
- Risk-averse – $U(L) < U(S_{EMV(L)})$
- Risk-seeking – $U(L) > U(S_{EMV(L)})$
- Risk-neutral – $U(L) = U(S_{EMV(L)})$



Two Concepts

- The ***certainty equivalent of a lottery***: the sum of money, X , which, if received with certainty will yield the same utility as the gamble
 X is **CE** if $u(X) = EU = p_G \times u(c_G) + p_B \times u(c_B)$
- The ***risk premium associated with a lottery*** is the maximum amount a person is prepared to pay to avoid the gamble
RP = EMV - CE

Risk averse



CE is the utility one get for sure when not choosing the lottery.

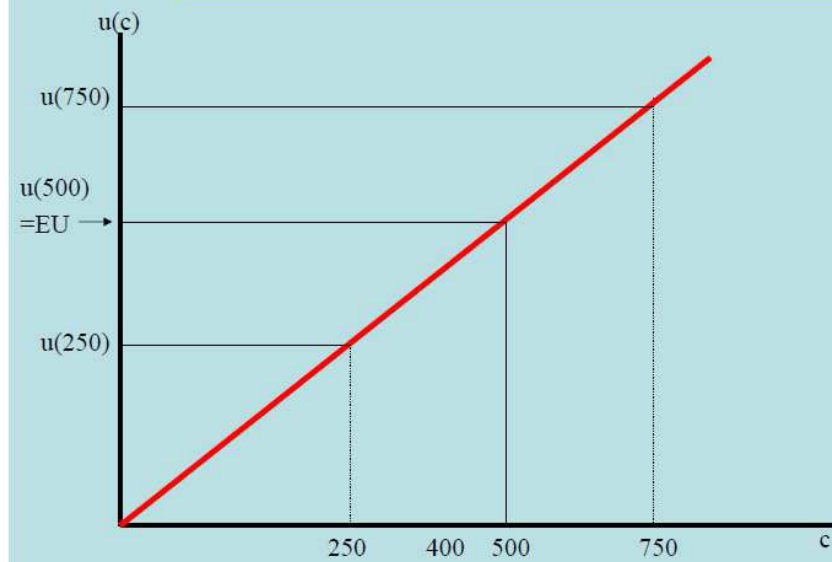
In our case 350.

The RP is the money someone pays for not participating in the lottery and getting the sure thing.

The risk premium is $150=500-350$.

Risk Neutral

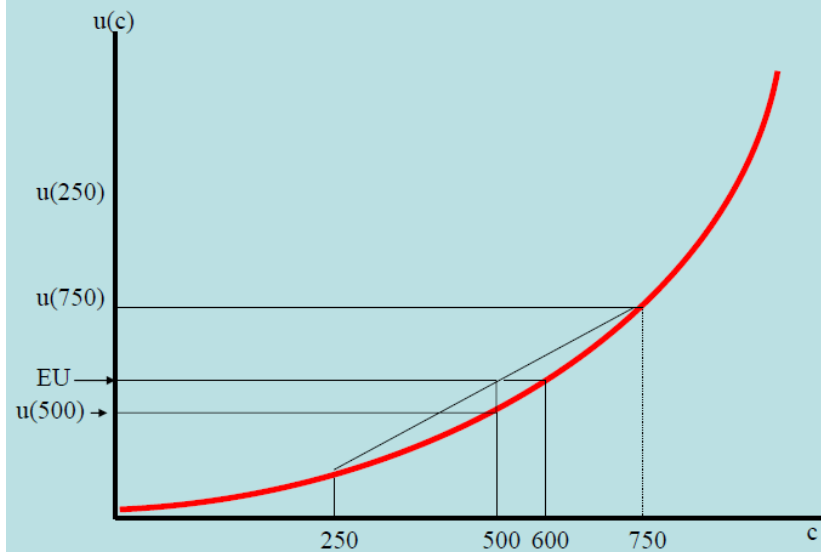
A risk neutral person / with constant MU / with a linear utility function will be indifferent between accepting/rejecting a fair gamble



The certainty equivalent of the gamble is \$500; the risk premium is \$0

Risk Seeking

A risk loving person / with increasing MU / with a convex utility function will accept a fair gamble



The certainty equivalent of the gamble is \$600; the risk premium is -\$100

29

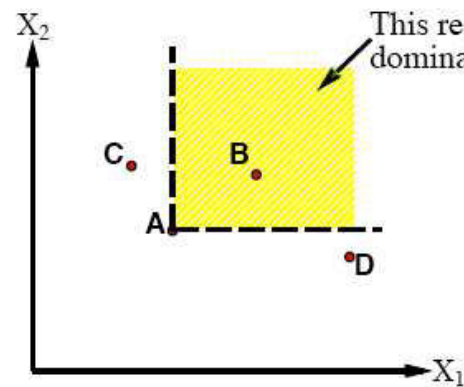
Multiattribute Utility Theory

- A given state may have multiple utilities
 - ◆ ...because of multiple evaluation criteria
 - ◆ ...because of multiple agents (interested parties) with different utility functions

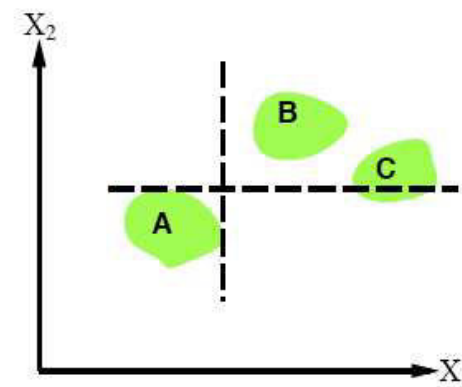
Strict dominance

Typically define attributes such that U is monotonic in each

Strict dominance: choice B strictly dominates choice A iff
 $\forall i X_i(B) \geq X_i(A)$ (and hence $U(B) \geq U(A)$)



Deterministic attributes



Uncertain attributes

Stochastic Dominance

- Introduced by Rothschild and Stiglitz (1970)
- When distribution $F(\cdot)$ yields unambiguously higher returns than $G(\cdot)$?
 - When **every** expected utility maximizer (who values more money over less) prefers $F(\cdot)$ to $G(\cdot)$
 - When for every amount of money x the probability of getting at least x is higher under $F(\cdot)$ than under $G(\cdot)$
- Fortunately, these two definitions are equivalent

Stochastic dominance

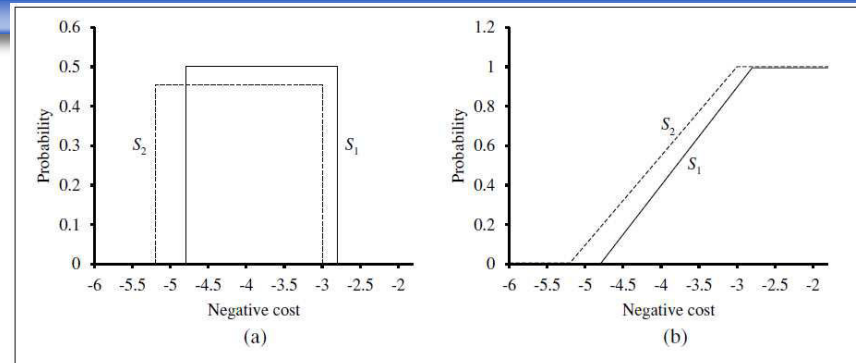


Figure 16.5 Stochastic dominance. (a) S_1 stochastically dominates S_2 on cost. (b) Cumulative distributions for the negative cost of S_1 and S_2 .

If two actions S_1 and S_2 lead to probability distributions $p_1(x)$ and $p_2(x)$ on attribute X , then S_1 stochastically dominates S_2 on X if:

$$\forall x \int_{-\infty}^x p_1(x') dx' \leq \int_{-\infty}^x p_2(x') dx'$$

For any monotonically non-decreasing utility function $U(x)$, the expected utility of S_1 is at least as high as the expected utility of S_2 . Hence, S_2 can be discarded.

Stochastic dominance contd.

Stochastic dominance can often be determined without exact distributions using **qualitative** reasoning

E.g., construction cost increases with distance from city

S_1 is closer to the city than S_2

$\Rightarrow S_1$ stochastically dominates S_2 on cost

E.g., injury increases with collision speed

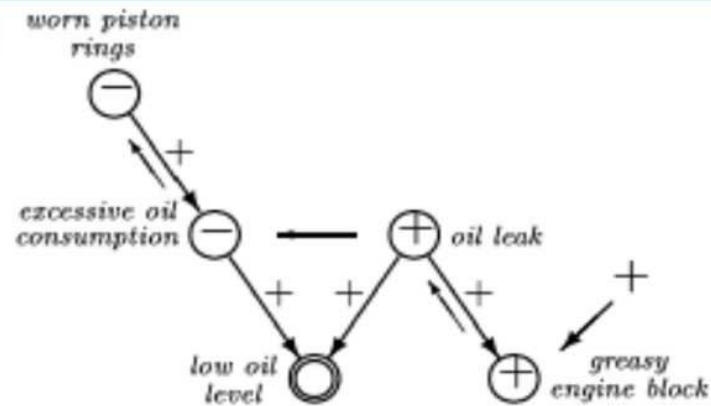
Can annotate belief networks with stochastic dominance information:

$X \xrightarrow{+} Y$ (X positively influences Y) means that

For every value \mathbf{z} of Y 's other parents \mathbf{Z}

$\forall x_1, x_2 \quad x_1 \geq x_2 \Rightarrow \mathbf{P}(Y|x_1, \mathbf{z})$ stochastically dominates $\mathbf{P}(Y|x_2, \mathbf{z})$

Example



Qualitative influence of greasy engine block on worn piston rings:

Greasy engine block is evidence of oil leak.

Oil leak and excessive oil consumption can each cause low oil level.

Oil leak explains low oil level and so is evidence against excessive oil consumption.

Decreased likelihood of excessive oil consumption is evidence against worn piston rings.

Therefore, greasy engine block is evidence against worn piston rings.

37

Preference structure: Deterministic

X_1 and X_2 preferentially independent of X_3 iff
 preference between $\langle x_1, x_2, x_3 \rangle$ and $\langle x'_1, x'_2, x_3 \rangle$
 does not depend on x_3

E.g., $\langle \text{Noise}, \text{Cost}, \text{Safety} \rangle$:

$\langle 20,000 \text{ suffer}, \$4.6 \text{ billion}, 0.06 \text{ deaths/mpm} \rangle$ vs.
 $\langle 70,000 \text{ suffer}, \$4.2 \text{ billion}, 0.06 \text{ deaths/mpm} \rangle$

Theorem (Leontief, 1947): if every pair of attributes is P.I. of its complement, then every subset of attributes is P.I. of its complement: **mutual P.I.**

Theorem (Debreu, 1960): mutual P.I. $\Rightarrow \exists$ additive value function:

$$V(S) = \sum_i V_i(X_i(S))$$

Hence assess n single-attribute functions; often a good approximation

Multi-attribute utility functions

- Multi-dimensional or multi-attribute utility theory deals with expressing such utilities
- Example: you are made a set of job offers, how do you decide?

$$u(\text{job-offer}) = u(\text{salary}) + u(\text{location}) + \\ u(\text{pension package}) + u(\text{career opportunities})$$

$$u(\text{job-offer}) = 0.4u(\text{salary}) + 0.1u(\text{location}) + \\ 0.3u(\text{pension package}) + 0.2u(\text{career opportunities})$$

But if there are interdependencies between attributes, then additive utility functions do not suffice. Multiplicative utility function:

$$u(x,y) = w_x u(x) + w_y u(y) + w_x w_y u(x)u(y)$$

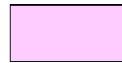
Decision Networks/ Influence diagrams

- Extend BNs to handle actions and utilities
- Also called *influence diagrams*
- Use BN inference methods
- Perform *Value of Information* calculations

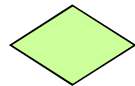
Decision Networks cont.



- Chance nodes: random variables, as in BNs: $X=\{x_1, \dots, x_n\}$



- Decision nodes: actions that decision maker can take: $A=\{a_1, \dots, a_n\}$



- Utility function nodes: the utility of the outcome state: $U(X,A)$

Expected Utility in DN/ID

$$EU[D(a)] = \sum_x P(x|a)U(x, a)$$

- Want to choose action a that maximizes the expected utility

$$a^* = \operatorname{argmax}_a EU[D(a)]$$

Simple example

m^0	m^1	m^2
0.5	0.3	0.2
poor	mid	great

Market

Found

U

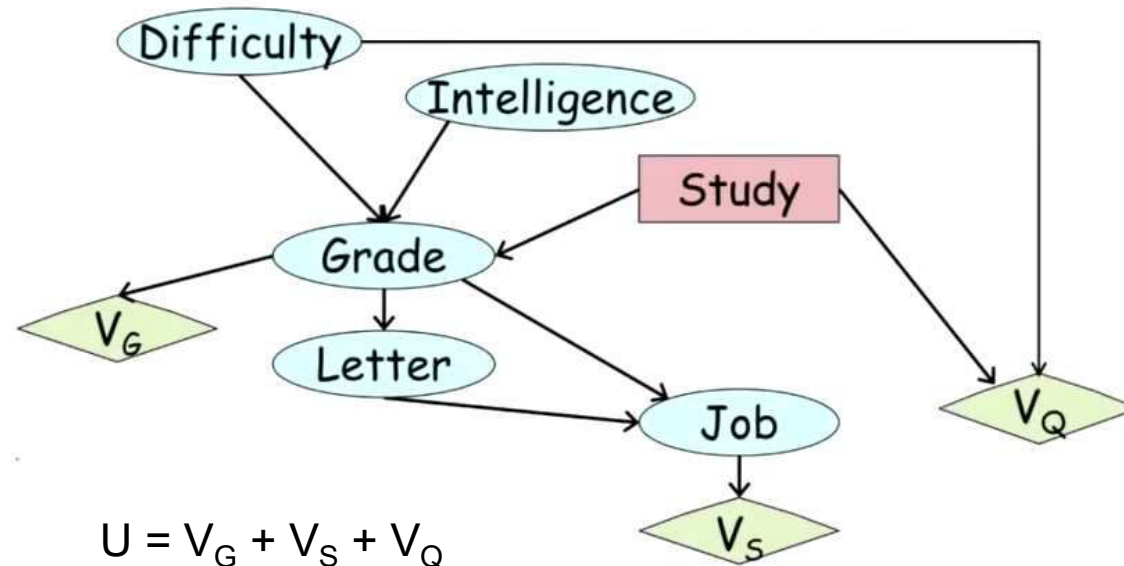
	f^0	f^1
m^0	0	-7
m^1	0	5
m^2	0	20

$$EU(f^0) = 0$$

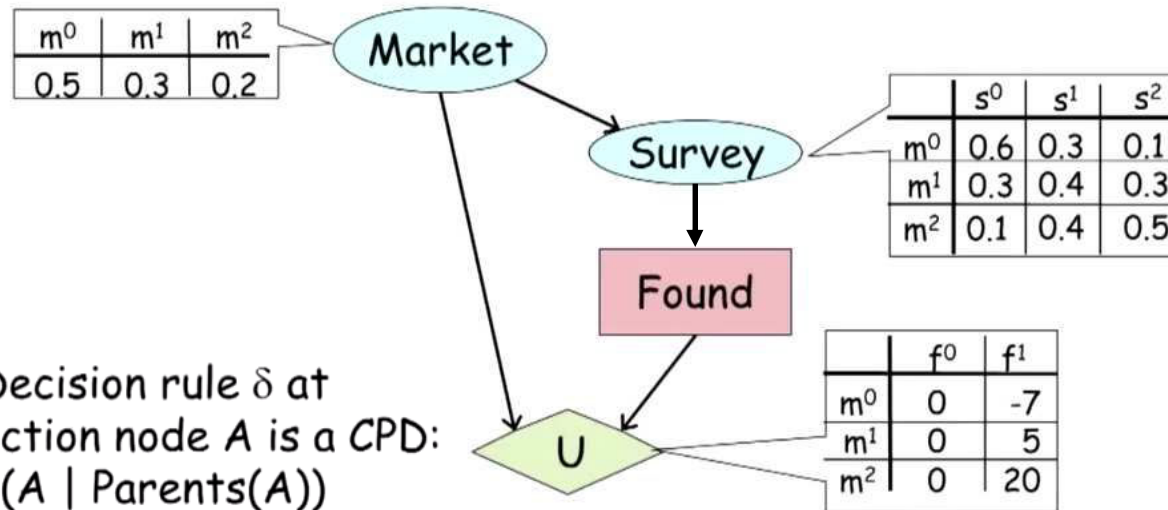
$$EU(f^1) = 0.5 \times -7 + 0.3 \times 5 + 0.2 \times 10 = \underline{2}$$

43

A more complex network



Information edges



45

Finding MEU Decision rules

$$\sum_{S,F} \delta_F(F | S) \sum_M P(M)P(S | M)U(F, M)$$

$$= \sum_{S,F} \delta_F(F | S)\mu(F, S)$$

m ⁰	m ¹	m ²
0.5	0.3	0.2

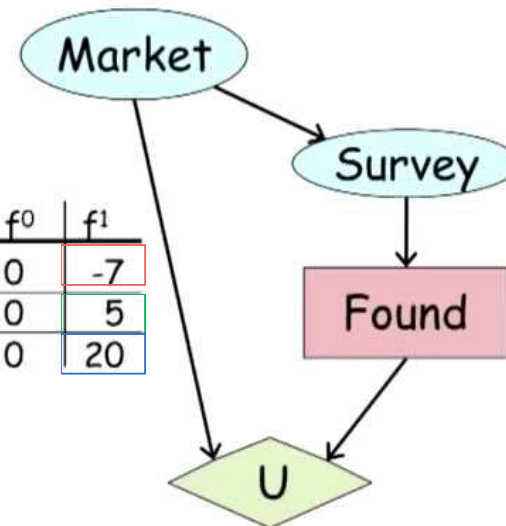
	s ⁰	s ¹	s ²
m ⁰	0.6	0.3	0.1
m ¹	0.3	0.4	0.3
m ²	0.1	0.4	0.5

	f ⁰	f ¹
m ⁰	0	-7
m ¹	0	5
m ²	0	20

$0.5 * 0.6 * -7 = -2.1$
 $0.3 * 0.3 * 5 = 0.45$
 $0.2 * 0.1 * 20 = 0.4$

	f ⁰	f ¹
s ⁰	0	
s ¹	0	
s ²	0	

Summing leads to -1.25



Finding MEU Decision rules

$$\sum_{S,F} \delta_F(F | S) \sum_M P(M)P(S | M)U(F, M)$$

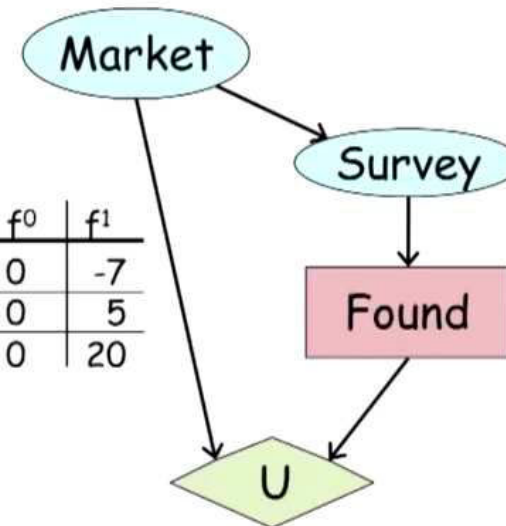
$$= \sum_{S,F} \delta_F(F | S)\mu(F, S)$$

m ⁰	m ¹	m ²
0.5	0.3	0.2

	s ⁰	s ¹	s ²
m ⁰	0.6	0.3	0.1
m ¹	0.3	0.4	0.3
m ²	0.1	0.4	0.5

	f ⁰	f ¹
m ⁰	0	-7
m ¹	0	5
m ²	0	20

	f ⁰	f ¹
s ⁰	0	-1.25
s ¹	0	1.15
s ²	0	2.1



$$\text{MEU} = 0 + 1.15 + 2.1 = 3.25$$

47

More Generally

$$\begin{aligned}
 \text{EU}[\mathcal{D}[\delta_A]] &= \sum_{\mathbf{x}, a} P_{\delta_A}(\mathbf{x}, a) U(\mathbf{x}, a) & \mathbf{Z} &= \mathbf{Pa}_A \\
 & & \mathbf{W} &= \{X_1, \dots, X_n\} - \mathbf{Z} \\
 &= \sum_{X_1, \dots, X_n, A} \left(\left(\prod_i P(X_i | \mathbf{Pa}_{X_i}) \right) U(\mathbf{Pa}_U) \delta_A(A | \mathbf{Z}) \right) \\
 &= \sum_{\mathbf{Z}, A} \delta_A(A | \mathbf{Z}) \sum_{\mathbf{W}} \left(\left(\prod_i P(X_i | \mathbf{Pa}_{X_i}) \right) U(\mathbf{Pa}_U) \right) \\
 &= \sum_{\mathbf{Z}, A} \delta_A(A | \mathbf{Z}) \mu(A, \mathbf{Z})
 \end{aligned}$$

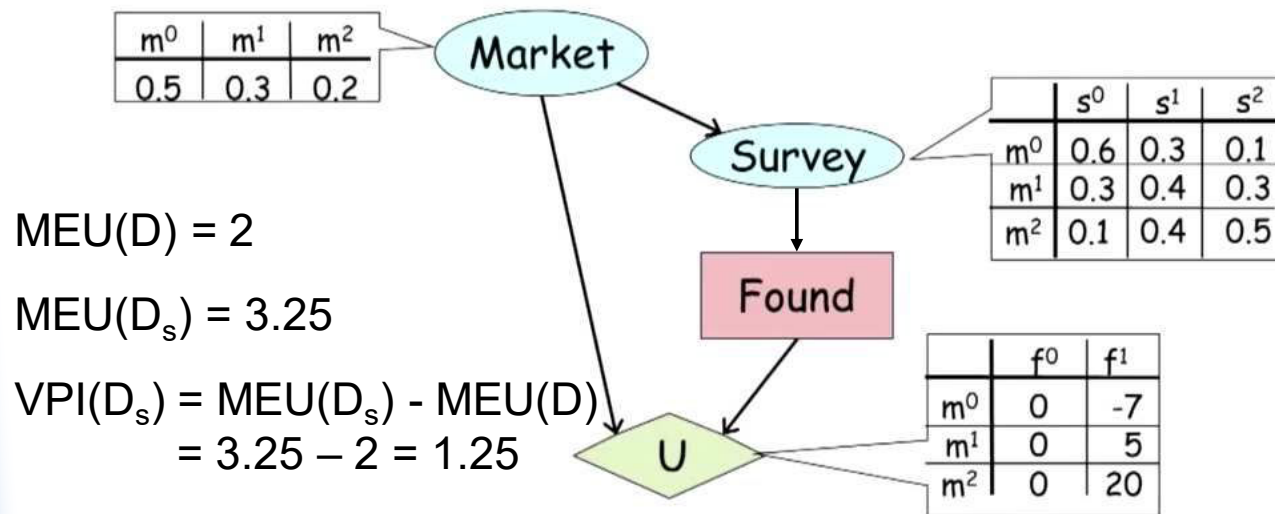
$$\delta_A^*(a | \mathbf{z}) = \begin{cases} 1 & a = \operatorname{argmax}_A \mu(A, \mathbf{z}) \\ 0 & \text{otherwise} \end{cases}$$

MEU Summary

- To compute MEU & optimize decision at A :
 - Treat A as random variable with arbitrary CPD
 - Introduce utility factor with scope Pa_U
 - Eliminate all variables except A, \mathbf{Z} (A 's parents) to produce factor $\mu(A, \mathbf{Z})$.
 - For each \mathbf{z} , set:

$$\delta_A^*(a | \mathbf{z}) = \begin{cases} 1 & a = \operatorname{argmax}_A \mu(A, \mathbf{z}) \\ 0 & \text{otherwise} \end{cases}$$

Value of Perfect Information



Value of Perfect Information

Current evidence \mathbf{E} , current best action a

Possible actions outcomes \mathbf{S}_i , potential new evidence \mathbf{E}_j

$$MEU(a|E) = \max_a \sum_i U(S_i)P(S_i|E, a)$$

Suppose we knew \mathbf{E}_j , we would choose $a_{e_{jk}}$

$$MEU(a_{e_{jk}}|E, E_j = e_{jk}) = \max_a \sum_i U(S_i)P(S_i|E, a, E_j = e_{jk})$$

\mathbf{E}_j is not known. Must compute expected gain.

$$VPI(E_j) = \left(\sum_k P(E_j|E) MEU(a_{e_{jk}}|E, E_j = e_{jk}) \right) - MEU(a|E)$$

51

Properties of VPI

Non negative

$$\forall j, E \ VPI_E(E_j) \geq 0$$

Non additive

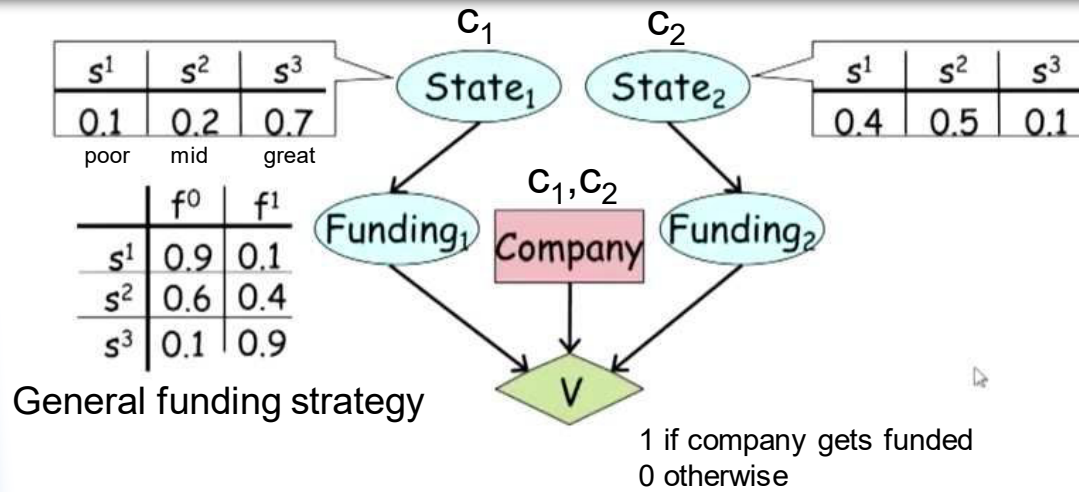
$$VPI_E(E_j, E_k) \neq VPI_E(E_j) + VPI_E(E_k)$$

Order-independent

$$VPI_E(E_j, E_k) = VPI_E(E_j) + VPI_{E, E_j}(E_k) = VPI_E(E_k) + VPI_{E, E_k}(E_j)$$

When is information useful?

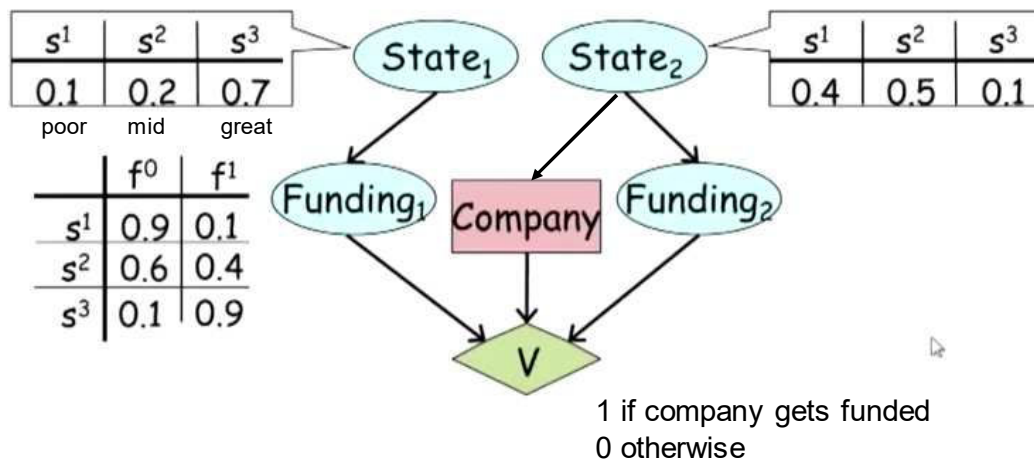
Example 1



$$EU(D[c_1]) = 0.1 \cdot 0.1 + 0.2 \cdot 0.4 + 0.7 \cdot 0.9 = 0.72$$

$$EU(D[c_2]) = 0.4 \cdot 0.1 + 0.5 \cdot 0.4 + 0.1 \cdot 0.9 = 0.33$$

Example 1



$$EU(D[c_1]) = 0.72$$

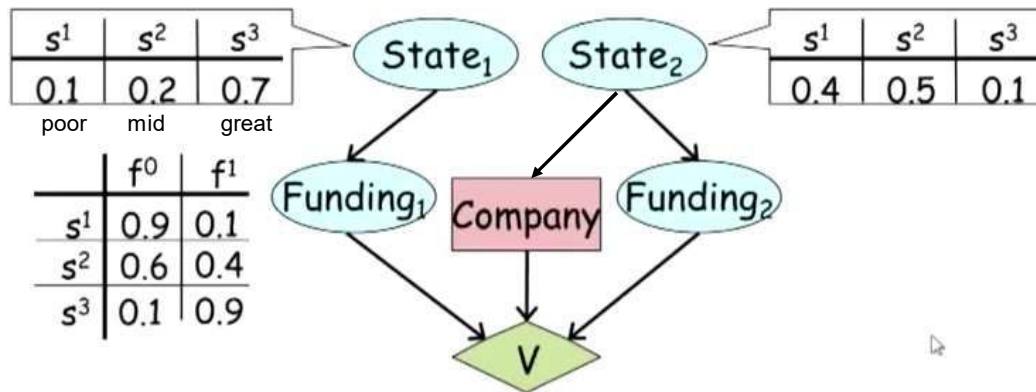
$$EU(D[c_2]) = 0.33$$

If c_2 is in state s_1 , the utility is 0.1

If c_2 is in state s_2 , the utility is 0.4

If c_2 is in state s_3 , the utility is 0.9

Example 1



1 if company gets funded
0 otherwise

$$MEU(D_{State2}) = 0.4 * 0.72 = 0.288$$

$$0.5 * 0.72 = 0.36$$

$$0.1 * 0.9 = 0.09$$

$$\Sigma 0.738$$

$$EU(D[c_1]) = 0.72$$

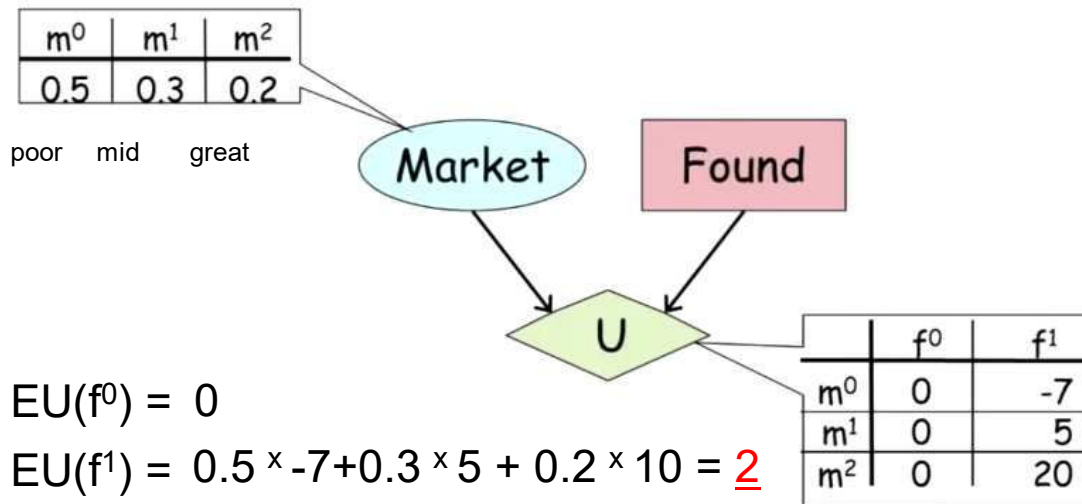
$$EU(D[c_2]) = 0.33$$

55

Last time

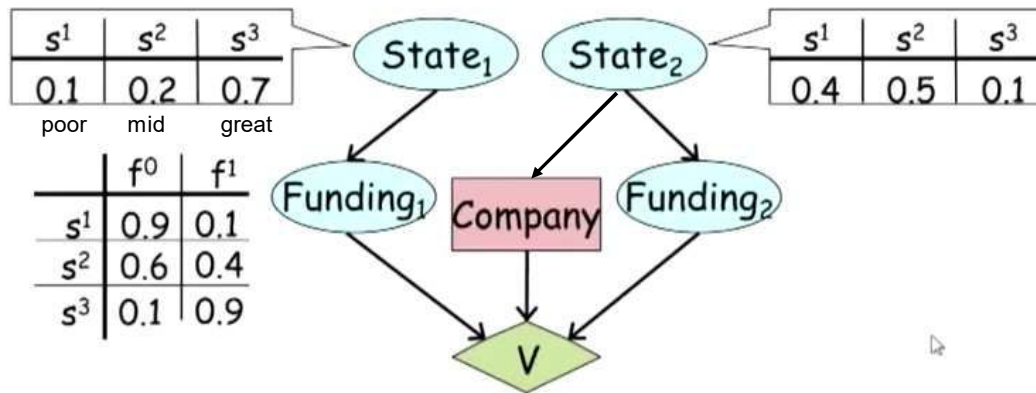
- Existence of a utility function
 - ◆ Additive vs multiplicative utility function
 - ◆ Stochastic dominance
- Risk profiles
 - ◆ Risk averse
 - ◆ Risk neutral
 - ◆ Risk seeking

Last time: Decision networks



57

Example 1



s^1	s^2	s^3
0.1	0.2	0.7
poor	mid	great

s^1	s^2	s^3
0.4	0.5	0.1

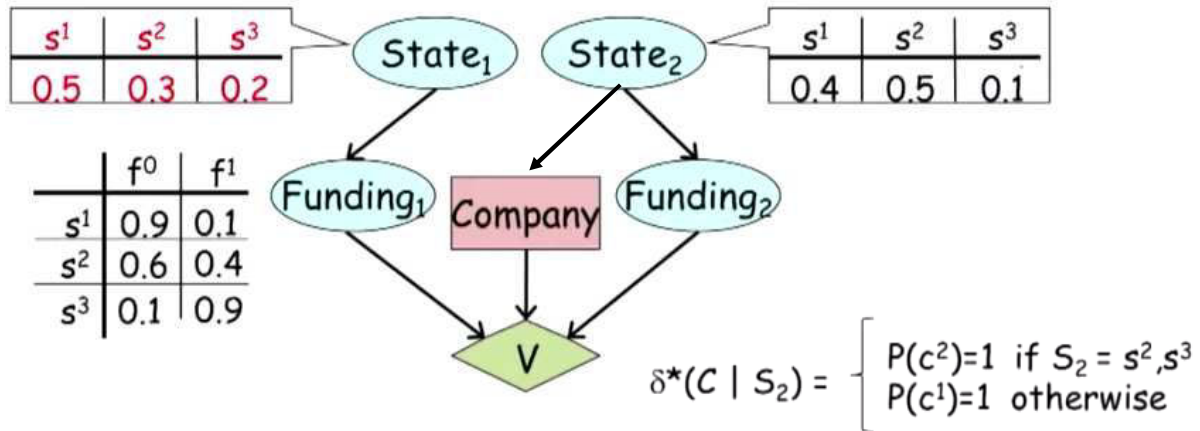
	f^0	f^1
s^1	0.9	0.1
s^2	0.6	0.4
s^3	0.1	0.9

1 if company gets funded
 0 otherwise
 Select c_2 if $State_2 = s_3$, c_1 otherwise

$EU(D[c_1]) = 0.72$
 $EU(D[c_2]) = 0.33$

$MEU(D_{State_2}) = 0.738$
 $VPI(D_{State_2}) = 0.738 - 0.72 = 0.018$

Example 2



$$EU(D[c_1]) = 0.35$$

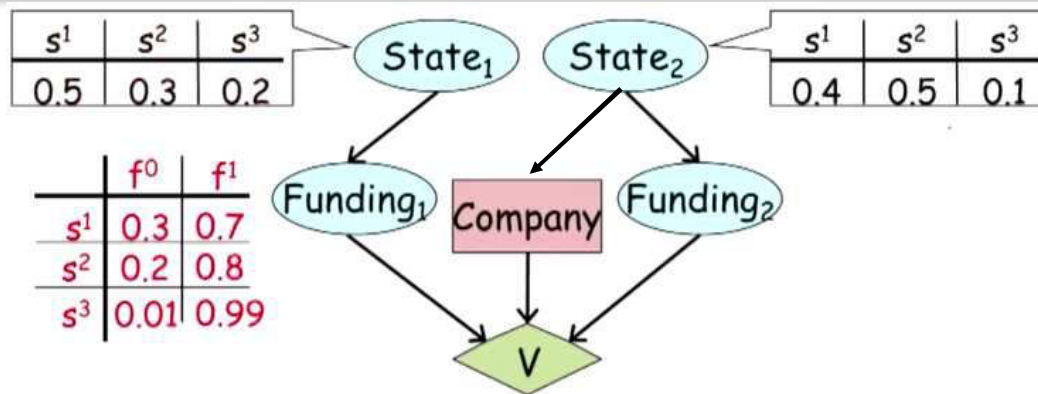
$$EU(D[c_2]) = 0.33$$

$$MEU(D_{S_2 \rightarrow c}) = 0.43$$

$$VPI = 0.43 - 0.35 = 0.08$$

59

Example 3



$$EU(D[c_1]) = 0.788$$

$$EU(D[c_2]) = 0.779$$

$$\delta^*(C | S_2) = \begin{cases} P(c^2)=1 & \text{if } S_2 = s^2, s^3 \\ P(c^1)=1 & \text{otherwise} \end{cases}$$

$$MEU(D_{S_2 \rightarrow C}) = 0.8142$$

$$VPI = 0.8142 - 0.788 = 0.0262$$

Summary

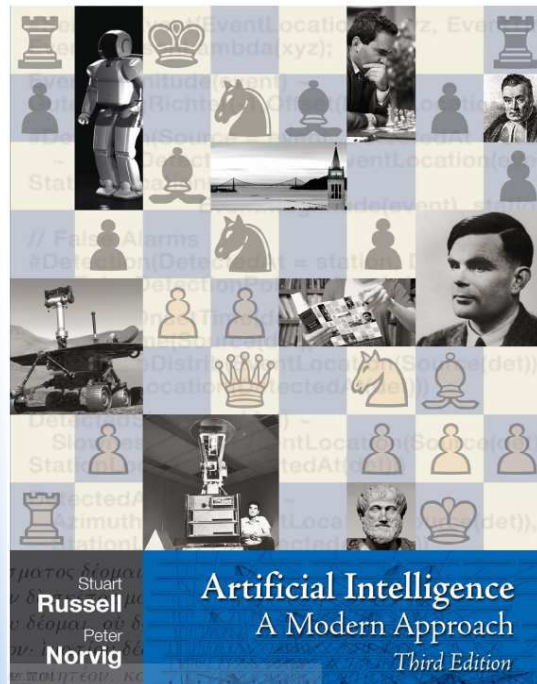
- Influence diagrams provide clear and coherent semantics for the value of making an observation
 - ♦ $VPI = P(\text{new observation}) * MEU(\text{new observation}) - MEU(\text{with current observations})$
- Information is valuable if and only if it induces a change in action in at least one context, and with (significant) higher MEU.

Intelligent Autonomous Agents and Cognitive Robotics

Topic 8: Decision-Making under Uncertainty Complex Decisions

Ralf Möller, Rainer Marrone
Hamburg University of Technology

Literature



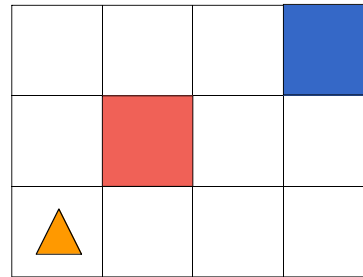
- Chapter 17

Material from Lise Getoor, Jean-Claude Latombe, Daphne Koller, and Stuart Russell

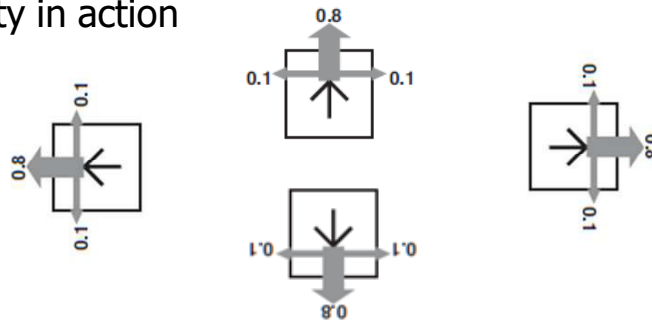
Sequential Decision Making

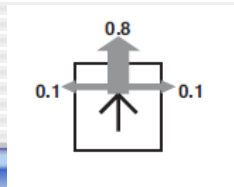
- Finite Horizon
 - ◆ Fixed time N after that nothing happens
- Infinite Horizon
 - ◆ N not fixed

Simple Robot Navigation Problem

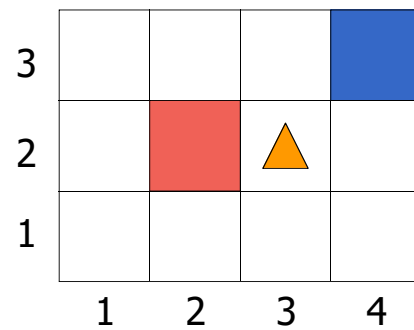


- In each state, the possible actions are **U, D, R, L**
Uncertainty in action





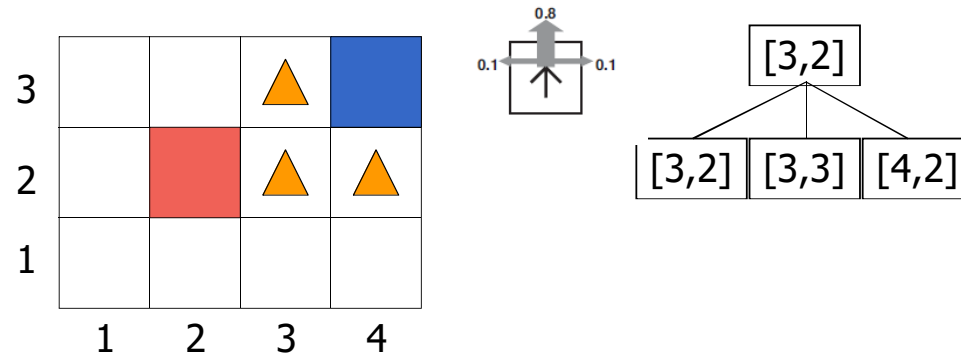
Sequence of Actions



[3,2]

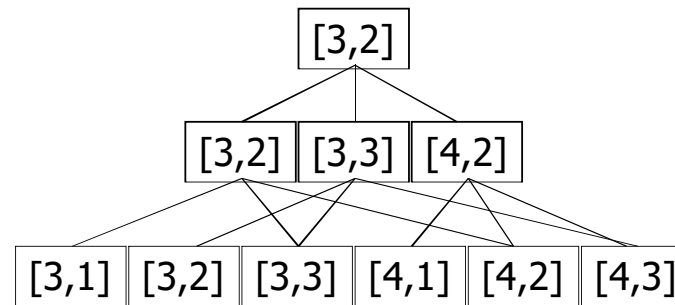
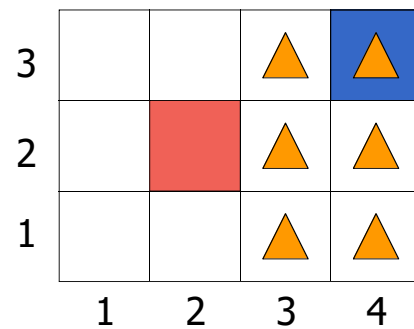
- Planned sequence of actions: (U, R)

Sequence of Actions



- Planned sequence of actions: (U, R)
- U is executed

Histories



- Planned sequence of actions: (U, R)
- U has been executed
- R is executed
- There are 9 possible sequences of states – called **histories** – and 6 possible final states for the robot!

Probability of Reaching the Goal

Note importance of Markov property
in this derivation

3			▲	▲
2		■	▲	▲
1			▲	▲
	1	2	3	4

$$\begin{aligned} \mathbf{P}([4,3] \mid (U,R).[3,2]) &= \mathbf{P}([3,3] \mid U.[3,2]) \times \mathbf{P}([4,3] \mid R.[3,3]) \\ &+ \mathbf{P}([4,2] \mid U.[3,2]) \times \mathbf{P}([4,3] \mid R.[4,2]) \\ &= \mathbf{0.8} \times \mathbf{0.8} + \mathbf{0.1} \times \mathbf{0.1} \\ &= 0.65 \end{aligned}$$

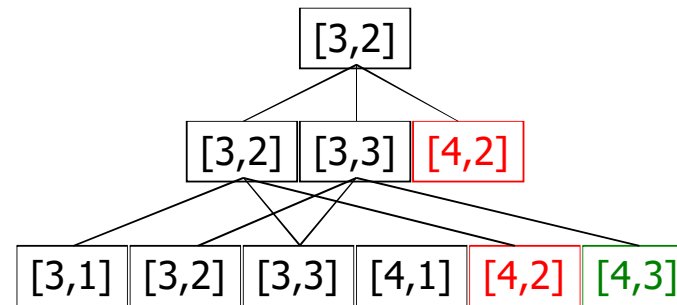
Utility of a History

3				+1
2				-1
1				
	1	2	3	4

- [4,3] provides power supply
- [4,2] is a sand area from which the robot cannot escape
- The robot needs to recharge its batteries
- [4,3] or [4,2] are terminal states
- The **utility of a history** is defined by the utility of the last state (+1 or -1) minus $n/25$, where n is the number of moves

Utility of an Action Sequence

3				+1
2				-1
1				
	1	2	3	4



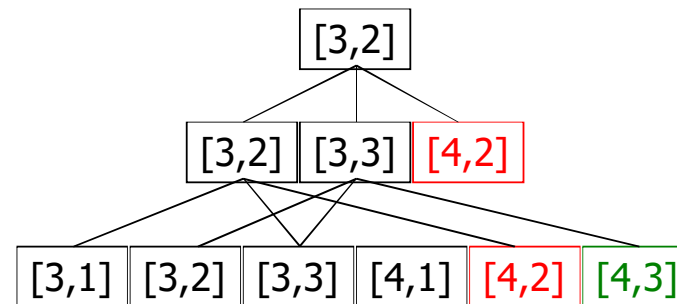
- Consider the action sequence (U,R) from [3,2]
- A run produces one among 7 possible histories, each with some probability
- The **utility of the sequence** is the expected utility of the histories:

$$u = \sum_h u_h \mathbf{P}(h)$$

10

Optimal Action Sequence

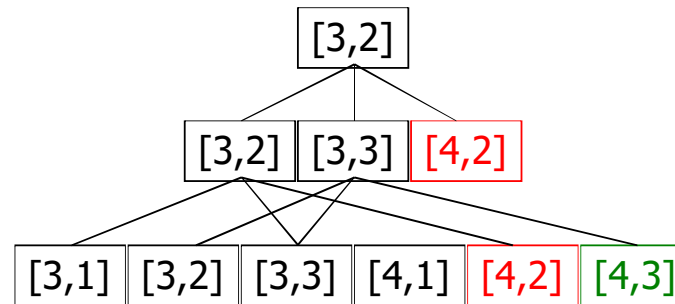
3				+1
2				-1
1				
	1	2	3	4



- Consider the action sequence (U,R) from [3,2]
- A run produces one among 7 possible histories, each with some probability
- The utility of the sequence is the expected utility of the histories
- The **optimal sequence** is the one with maximal utility

Optimal Action Sequence

3				+1
2				-1
1				
	1	2	3	4



- Consider the action sequence (U,R) from [3,2]
- A run prod **only if the sequence is executed blindly!** me probability
- The utility of the sequence is the expected utility of the histories
- The **optimal sequence** is the one with maximal utility
- **But is the optimal action sequence what we want to compute?**

12

Policy (Reactive/Closed-Loop Strategy)

3	→	→	→	+1
2	↑		↑	-1
1	↑	←	←	←
	1	2	3	4

- A **policy Π** is a complete mapping from states to actions

Reactive Agent Algorithm

Repeat:

- ◆ $s \leftarrow$ sensed state
- ◆ If s is terminal then exit
- ◆ $a \leftarrow \Pi(s)$
- ◆ Perform a

Optimal Policy

3	→	→	→	+1
2	↑		↑	-1
1	↑	←	←	←
	1	2	3	4

- A **policy** Π is a complete mapping from states to actions
- The **optimal policy** Π^* is the one that always yields a history (ending at a terminal state) with maximal expected utility

Optimal Policy

3	→	→	→	+1
2	↑		↑	-1
1	↑	←	←	←
	1	2	3	4

- A **policy** Π is a complete mapping from states to actions
- The **optimal policy** Π^* is the one that always yields a history with maximal expected utility

How to compute Π^* ?

This problem is called a
Markov Decision Problem
(MDP)

16

Additive Utility: Stationarity

- History $H = (s_0, s_1, \dots, s_n)$
- The utility of H is **additive** iff:

$$U(s_0, s_1, \dots, s_n) = R(0) + U(s_1, \dots, s_n) = \sum R(i)$$



Reward

The diagram consists of two blue arrows pointing upwards from the word 'Reward' to the terms $R(0)$ and $U(s_1, \dots, s_n)$ in the equation above. This illustrates that the total utility is composed of the immediate reward plus the utility of the remaining history.

Additive Utility

- History $H = (s_0, s_1, \dots, s_n)$
- The utility of H is **additive** iff:
$$U(s_0, s_1, \dots, s_n) = R(0) + U(s_1, \dots, s_n) = \sum R(i)$$
- Robot navigation example:
 - ♦ $R(n) = +1$ if $s_n = [4, 3]$
 - ♦ $R(n) = -1$ if $s_n = [4, 2]$
 - ♦ $R(i) = -1/25$ if $i = 0, \dots, n-1$

Principle of Max Expected Utility

- History $H = (s_0, s_1, \dots, s_n)$
- Utility of H : $\mathcal{U}(s_0, s_1, \dots, s_n) = \sum \mathcal{R}(i)$

3				+1
2				-1
1				
	1	2	3	4

→	→	→	+1
↑		↑	-1
↑	←	←	←

First-step analysis →

- $\mathcal{U}(i) = \mathcal{R}(i) + \max_a \sum_j \mathbf{P}(j | a, i) \mathcal{U}(j)$
- $\Pi^*(i) = \arg \max_a \sum_j \mathbf{P}(k | a, i) \mathcal{U}(j)$

Value Iteration

- Initialize the utility of each non-terminal state s_i to $U_0(i) = 0$
- For $t = 0, 1, 2, \dots$, do:

$$U_{t+1}(i) \leftarrow R(i) + \max_a \sum_k \mathbf{P}(k | a, i) U_t(k) \quad (\text{Bellman equation})$$

3				+1
2				-1
1				
	1	2	3	4

20

Initialization

3	0	0	0	+1
2	0		0	-1
1	0	0	0	0
	1	2	3	4

Iteration 1

3	0	0	0	+1
2	0		0	-1
1	-0.04	0	0	0
	1	2	3	4

$$U^{(1)}(1,1) = -0.04 + 1 * \max \begin{bmatrix} 0.8U^{(0)}(1,2) + 0.1U^{(0)}(2,1) + 0.1U^{(0)}(1,1) & UP \\ 0.9U^{(0)}(1,1) + 0.1U^{(0)}(1,2) & LEFT \\ 0.9U^{(0)}(1,1) + 0.1U^{(0)}(2,1) & DOWN \\ 0.8U^{(0)}(2,1) + 0.1U^{(0)}(1,2) + 0.1U^{(0)}(1,1) & RIGHT \end{bmatrix}$$

$$U^{(1)}(1,1) = -0.04 + \max \begin{bmatrix} 0 & UP \\ 0 & LEFT \\ 0 & DOWN \\ 0 & RIGHT \end{bmatrix}$$

Initialization

3	0	0	0	+1
2	0		0	-1
1	0	0	0	0
	1	2	3	4

Iteration 1

3	0	0	0.76	+1
2	0		0	-1
1	-0.04	0	0	0
	1	2	3	4

$$U^{(1)}(3,3) = -0.04 + 1 * \max \begin{bmatrix} 0.8U^{(0)}(3,3) + 0.1U^{(0)}(2,3) + 0.1U^{(0)}(4,3) & UP \\ 0.8U^{(0)}(2,3) + 0.1U^{(0)}(3,3) + 0.1U^{(0)}(3,2) & LEFT \\ 0.8U^{(0)}(3,2) + 0.1U^{(0)}(2,3) + 0.1U^{(0)}(4,3) & DOWN \\ 0.8U^{(0)}(4,3) + 0.1U^{(0)}(3,3) + 0.1U^{(0)}(3,2) & RIGHT \end{bmatrix}$$

$$U^{(1)}(3,3) = -0.04 + \max \begin{bmatrix} 0.1 & UP \\ 0 & LEFT \\ 0.1 & DOWN \\ 0.8 & RIGHT \end{bmatrix}$$

22

After a Full Iteration

- Only the state one step away from a positive reward (3,3) has gained value, all the others are losing value because of the cost of moving

Iteration 1

3	-.04	-.04	0.76	+1
2	-.04		-.04	-1
1	-.04	-.04	-.04	-.04
	1	2	3	4

Value Iteration: from state utilities to *policy*

- Now the agent can choose the action that implements the MEU principle: maximize the expected utility of the subsequent state

$$\pi^*(s) = \arg \max_a \sum_{s'} P(s'|s,a)U(s')$$

states reachable from s by doing a

Probability of getting to s' from s via a

expected value of following policy π^* in s'

Example

$$\pi^*(s) = \arg \max_a \sum_{s'} P(s'|s,a)U(s')$$

3	0.812	0.868	0.912	+1
2	0.762		0.660	-1
	0.705	0.655	0.611	0.388
	1	2	3	4

Green arrows point from the cell (1,1) to its neighbors: UP, LEFT, and DOWN.

➤ To find the best action in (1,1)

$$\pi^*(1,1) = \arg \max \begin{bmatrix} 0,7456 & UP \\ 0,7107 & LEFT \\ 0.9U(1,1) + 0.1U(2,1) & DOWN \\ 0.8U(2,1) + 0.1U(1,2) + 0.1U(1,1) & RIGHT \end{bmatrix}$$

➤ We have to do this for all fields!!!!

Example

$$\pi^*(s) = \arg \max_a \sum_{s'} P(s'|s,a)U(s')$$

3	0.812	0.868	0.912	+1
2	0.762		0.660	-1
1	0.705	0.655	0.611	0.388
	1	2	3	4

➤ To find the best action in (1,1)

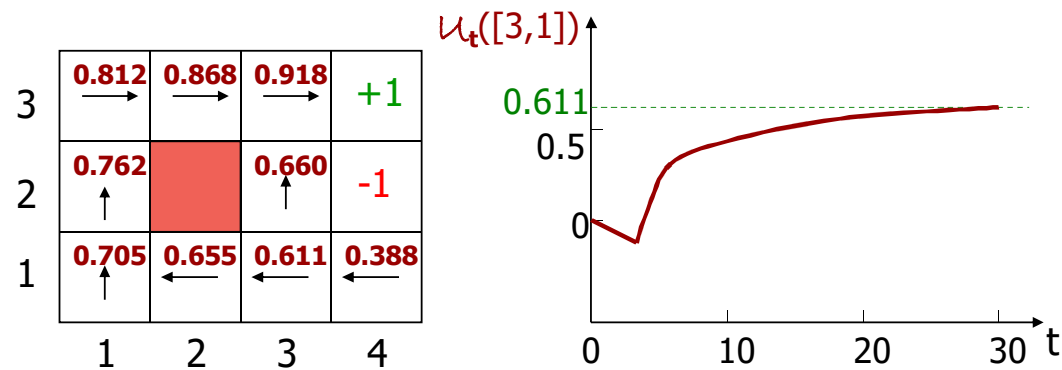
$$\pi^*(1,1) = \arg \max \begin{cases} 0,7456 \\ 0,7107 \\ 0,7 \\ 0,6707 \end{cases} \begin{cases} \text{UP} \\ \text{LEFT} \\ \text{DOWN} \\ \text{RIGHT} \end{cases}$$

➤ give *Up* as best action

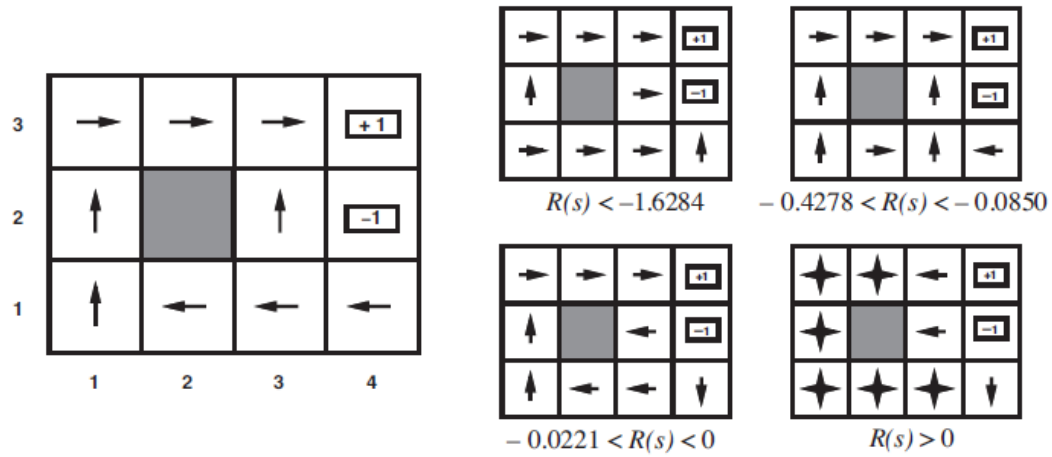
Value Iteration: the result

- Initialize the utility of each non-terminal state s_i to $U_0(i) = 0$
- For $t = 0, 1, 2, \dots$, do:

$$U_{t+1}(i) \leftarrow R(i) + \max_a \sum_k P(k | a, i) U_t(k)$$



The Reward is important



Infinite Horizon

In many problems, e.g., the robot navigation example, histories are potentially unbounded and the same state can be reached many times

3				+1
2				-1
1				
	1	2	3	4

$$U(i) = R(i) + \gamma \max_a \sum_j P(j | a, i) U(j)$$

One trick:
Use discounting to make infinite
Horizon problem mathematically
tractable

Value Iteration (finite and non-finite)

function VALUE-ITERATION(mdp, ϵ) **returns** a utility function
inputs: mdp , an MDP with states S , actions $A(s)$, transition model $P(s' | s, a)$,
rewards $R(s)$, discount γ
 ϵ , the maximum error allowed in the utility of any state
local variables: U, U' , vectors of utilities for states in S , initially zero
 δ , the maximum change in the utility of any state in an iteration

repeat
 $U \leftarrow U'; \delta \leftarrow 0$
 for each state s **in** S **do**
 $U'[s] \leftarrow R(s) + \gamma \max_{a \in A(s)} \sum_{s'} P(s' | s, a) U[s']$
 if $|U'[s] - U[s]| > \delta$ **then** $\delta \leftarrow |U'[s] - U[s]|$
until $\delta < \epsilon(1 - \gamma)/\gamma$
return U

31

Bellmann eq. is a contraction

- two important properties of contractions:
 - ◆ A contraction has **only one fixed point**; if there were two fixed points they would not get closer together when the function was applied, so it would not be a contraction.
 - ◆ **When the function is applied to any argument, the value must get closer to the fixed point**, so repeated application of a contraction always reaches the fixed point in the limit.

Value iteration

- Let U_i denote the vector of utilities for all the states at the i th iteration. Then the Bellman update equation can be written as

$$U_{i+1} \leftarrow BU_i$$

Value iteration

- use the **max norm**, which measures the “length” of a vector by the absolute value of its biggest component:

$$\|U\| = \max_s |U(s)|$$

- Let U_i and U'_i be *any two utility vectors*. Then we have

$$\|BU_i - BU'_i\| \leq \gamma \|U_i - U'_i\| \quad \boxed{17.7}$$

Value iteration

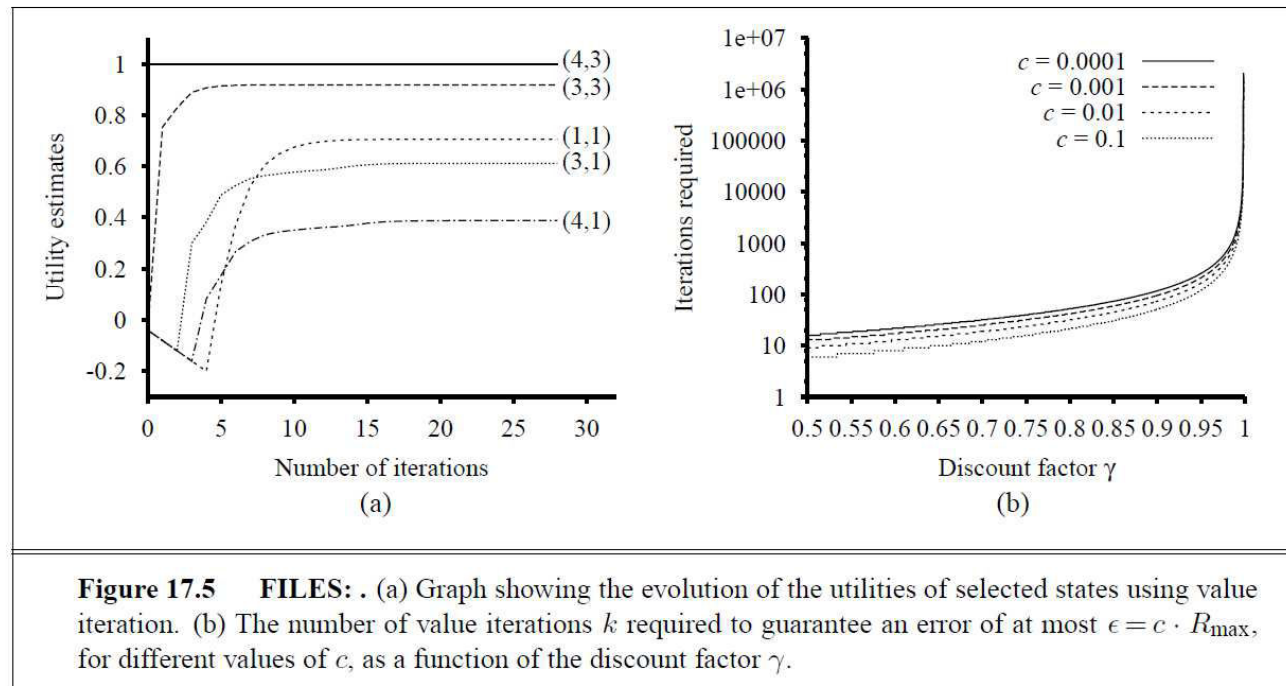


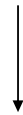
Figure 17.5 FILES: . (a) Graph showing the evolution of the utilities of selected states using value iteration. (b) The number of value iterations k required to guarantee an error of at most $\epsilon = c \cdot R_{\max}$, for different values of c , as a function of the discount factor γ .

6

Value iteration

- From the contraction, it can be shown that if the update is small (i.e., no state's utility changes by much), then the error, compared with the true utility function, also is small. More precisely,

$$\text{if } \|U_{i+1} - U_i\| < \varepsilon(1-\gamma)/\gamma \text{ then } \|U_{i+1} - U\| \leq \varepsilon \quad (17.8)$$



This is the stopping criteria for value iteration

Value iteration

- But the crucial question is!!!! How well will I do using this utility function?

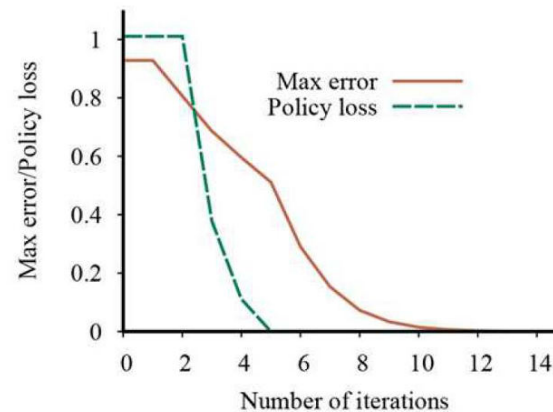
- **policy loss**

$U^{\pi_i}(s)$ is the utility obtained if π_i is executed starting in s ,
policy loss $\|U^{\pi_i} - U\|$ is the most the agent can lose by executing π_i *instead of the optimal* policy π^*

Value iteration

- The **policy loss** of π_i is connected to the error in U_i by the following inequality:

$$\text{if } \|U_i - U\| < \varepsilon \text{ then } \|U^{\pi_i} - U\| < 2\varepsilon \quad (17.9)$$



The maximum error $\|U_i - U\|$ of the utility estimates and the policy loss $\|U^{\pi_i} - U\|$, as a function of the number of iterations of value iteration on the 4×3 world.

Policy Iteration

- Pick a policy Π at random

Policy Iteration

- Pick a policy Π at random
- Repeat:
 - ♦ Policy evaluation
Compute the utility of each state for Π

$$U_t(i) \leftarrow R(i) + \sum_k \mathbf{P}(k | \Pi(i).i) U_t(k)$$

Policy Iteration

- Pick a policy Π at random
- Repeat:
 - ♦ Policy evaluation
Compute the utility of each state for Π
$$U_{\Pi}(i) \leftarrow R(i) + \sum_k \mathbf{P}(k | \Pi(i), i) U_{\Pi}(k)$$
 - ♦ Policy improvement:
Compute the policy Π' given these utilities
$$\Pi'(i) = \arg \max_a \sum_k \mathbf{P}(k | a, i) U_{\Pi}(k)$$

Policy Iteration

- Pick a policy Π at random
- Repeat:
 - ♦ Policy evaluation:
Compute the utility of each state for Π :
$$U_t(i) \leftarrow R(i) + \sum_k \mathbf{P}(k | \Pi(i), i) U_t(k)$$
 - ♦ Policy improvement:
Compute the policy Π' given these utilities
$$\Pi'(i) = \arg \max_a \sum_k \mathbf{P}(k | a, i) U(k)$$
 - ♦ If $\Pi' = \Pi$ then return Π

Policy Iteration

```
function POLICY-ITERATION(mdp) returns a policy
inputs: mdp, an MDP with states  $S$ , actions  $A(s)$ , transition model  $P(s' | s, a)$ 
local variables:  $U$ , a vector of utilities for states in  $S$ , initially zero
                   $\pi$ , a policy vector indexed by state, initially random

repeat
   $U \leftarrow$  POLICY-EVALUATION( $\pi, U, mdp$ )
  unchanged?  $\leftarrow$  true
  for each state  $s$  in  $S$  do
    if  $\max_{a \in A(s)} \sum_{s'} P(s' | s, a) U[s'] > \sum_{s'} P(s' | s, \pi[s]) U[s']$  then do
       $\pi[s] \leftarrow \operatorname{argmax}_{a \in A(s)} \sum_{s'} P(s' | s, a) U[s']$ 
      unchanged?  $\leftarrow$  false
until unchanged?
return  $\pi$ 
```


Linear equations

- By removing the max operator (Value Iteration) we can *also* solve the set of linear equations:

$$u(i) = \mathcal{R}(i) + \sum_k \mathbf{P}(k | \Pi(i).i) u(k)$$

(often a sparse system)

- Suppose we have $\Pi(1).1 = U_p$ $\Pi(2).2 = U_p$
 $U(1,1) = -0.04 + 0.8U(1,2) + 0.1U(1,1) + 0.1U(2,1)$
 $U(1,2) = -0.04 + 0.8U(1,3) + 0.2U(1,2)$
 ...
- Can be solved in $O(n^3)$ by standard linear algebra methods
- For large state spaces we can mix value iteration and policy iteration

Further optimization

- All algorithms require updating the utility or policy for all states at once.
- At each step we can also select a subset for updating
asynchronous policy iteration/mod. value iter.
(can show it will converge if some conditions for initial policy and utility function hold)
- Leads to heuristic algorithms that concentrate on states that are likely to be reached by a good policy.
 - ♦ “if one has no intention of throwing oneself off a cliff, one should not spend time worrying about the exact value of the resulting state”

46

Summary

- Decision making under uncertainty
- Sequential decision making
 - ◆ Utility function
 - ◆ Value iteration
 - ◆ Policy iteration

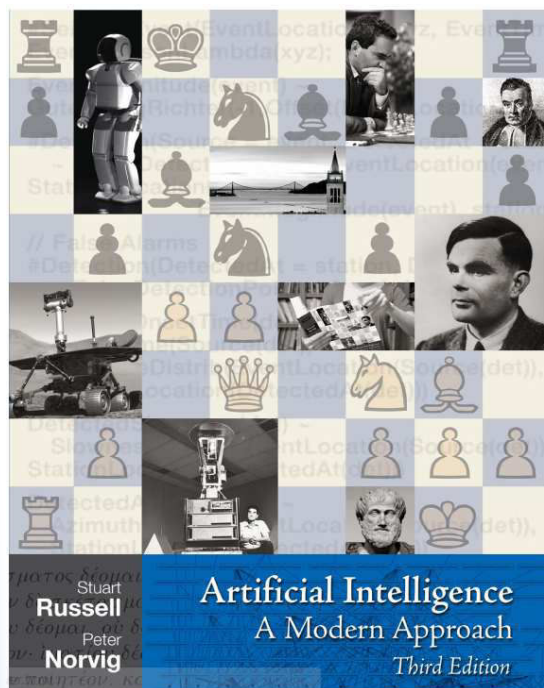
Intelligent Autonomous Agents

Agents and Rational Behavior

Topic 9: Decision-Making under Uncertainty Decision-Theoretic Agent Design

Ralf Möller, Rainer Marrone
Hamburg University of Technology

Literature



- Chapter 17

Last time

- Sequential decision making (uncertain actions)
 - ◆ Need a policy -> best action for each possible state
- Finding the best policy

- ◆ Value iteration

```

repeat
   $U \leftarrow U'; \delta \leftarrow 0$ 
  for each state  $s$  in  $S$  do
     $U'[s] \leftarrow R(s) + \gamma \max_{a \in A(s)} \sum_{s'} P(s'|s,a) U[s']$ 
    if  $|U'[s] - U[s]| > \delta$  then  $\delta \leftarrow |U'[s] - U[s]|$ 
until  $\delta < \epsilon(1-\gamma)/\gamma$ 

```

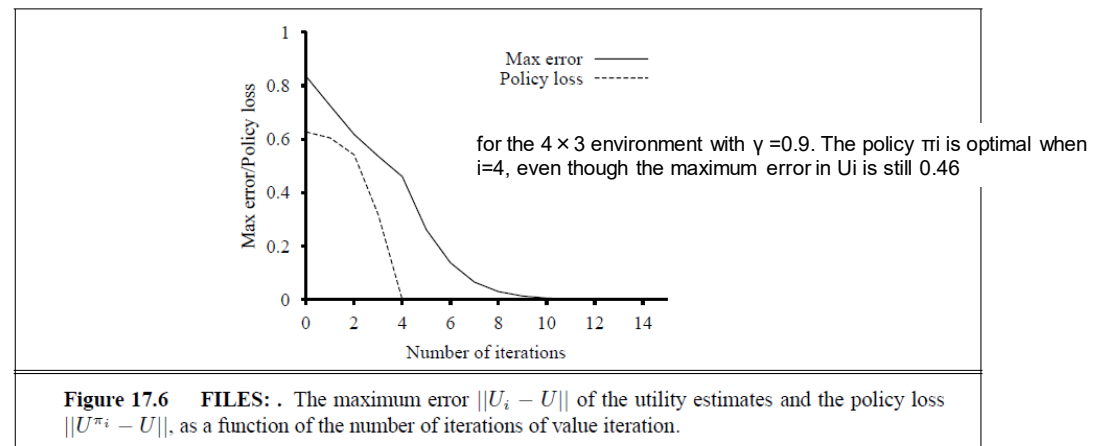
$$\pi^*(s) = \arg \max_a \sum_{s'} P(s'|s,a) U(s')$$

- ◆ Bellman update is a contraction →
Lead to the definition of when to stop value iteration.

Policy Loss

- The **policy loss** of π_i is connected to the error in U_i by the following inequality:

$$\text{if } \|U_i - U\| < \varepsilon \text{ then } \|U^{\pi_i} - U\| < 2\varepsilon \quad (17.9)$$



Last time: Policy iteration

- Create a random policy

Repeat:

- ◆ Value determination

$$U_t(i) \leftarrow R(i) + \sum_k \mathbf{P}(k | \Pi(i)) U_t(k)$$

- ◆ Policy Update:

$$\Pi'(i) = \arg \max_a \sum_k \mathbf{P}(k | a, i) U(k)$$

- ◆ If $\Pi' = \Pi$ then return Π

- We can combine Value- and Policy Iteration to get the best of both

Further optimization

- All algorithms require updating the utility or policy for all states at once.
- At each step we can also select a subset for updating
asynchronous policy iteration/mod. value iter.
(can show it will converge if some conditions for initial policy and utility function hold)
- Leads to heuristic algorithms that concentrate on states that are likely to be reached by a good policy.
 - ♦ “if one has no intention of throwing oneself off a cliff, one should not spend time worrying about the exact value of the resulting state”

Summary

- Decision making under uncertainty
- Sequential decision making
 - ◆ Utility of histories
 - ◆ Value iteration
 - ◆ Policy iteration

Jumping-off Point

- Let us assume again that the agent lives in the 4x3 environment
- The agent knows the environment
- BUT
 - ◆ Agent has no or very unreliable sensors
 - ◆ It does not make sense to determine the optimal policy wrt. a single state
 - ◆ $\Pi^*(s)$ is not well defined

POMDP: Uncertainty

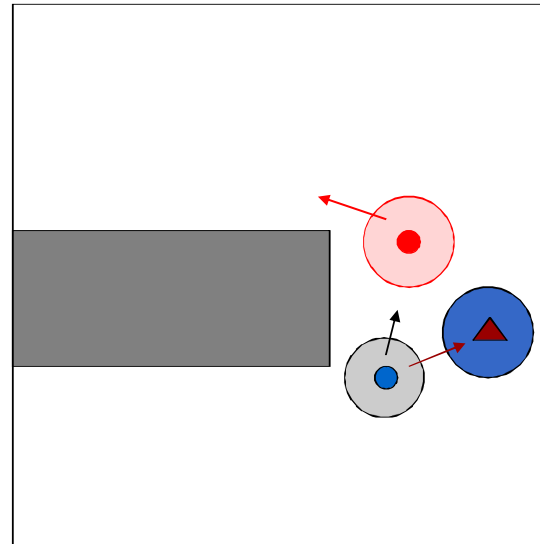
- Uncertainty about the action outcome
- Uncertainty about the world state due to imperfect (partial) information

Example: Target Tracking

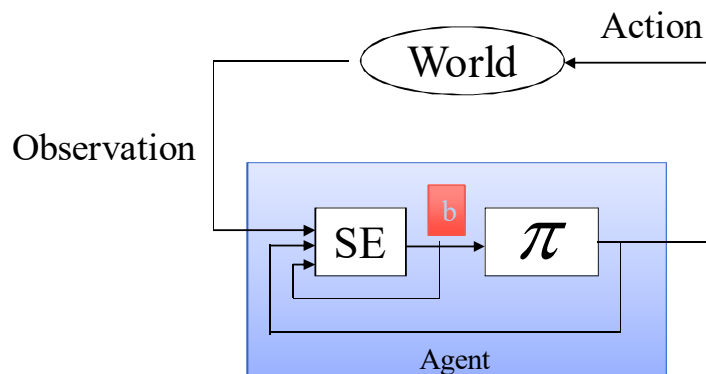
There is uncertainty in the robot's and target's positions; this uncertainty grows with further motion

There is a risk that the target may escape behind the corner, requiring the robot to move appropriately

But there is a positioning landmark nearby. Should the robot try to reduce its position uncertainty?



Decision cycle of a POMDP agent



- Given the current belief state b , execute the action $a = \pi^*(b)$
- Receive observation o
- Set the current belief state to $FORWARD(b, a, o)$ and repeat

Example Scenario

The agent has no sensors!!!

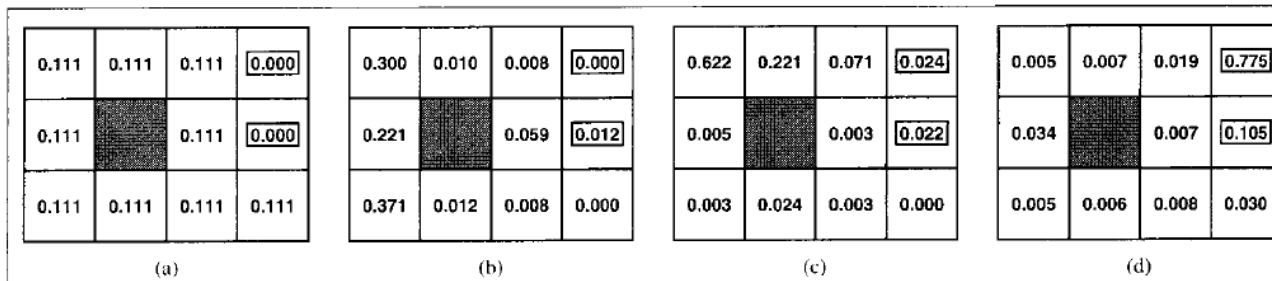


Figure 17.8 (a) The initial probability distribution for the agent's location. (b) After moving *Left* five times. (c) After moving *Up* five times. (d) After moving *Right* five times.

Belief state

- $b(s)$ is the probability assigned to the actual state s by belief state b .

0.111	0.111	0.111	<u>0.000</u>
0.111		0.111	<u>0.000</u>
0.111	0.111	0.111	0.111

if \mathbf{a} is executed in \mathbf{b} and observation is \mathbf{e} the belief in \mathbf{s}' is?

$$\left(\frac{1}{9}, \frac{1}{9}, \frac{1}{9}, \frac{1}{9}, \frac{1}{9}, \frac{1}{9}, \frac{1}{9}, \frac{1}{9}, \frac{1}{9}, \frac{1}{9}, 0, 0\right)$$

$$b'(s') = P(e|s') \sum_s P(s'|s, a) b(s) \text{ this is Filtering}$$

$$b' = \alpha \text{ FORWARD}(b, a, e)$$

Outcome of actions

- Probability of an observation \mathbf{e} given that \mathbf{a} was performed in \mathbf{b}
$$P(\mathbf{e}|\mathbf{a},\mathbf{b}) = \sum_{s'} P(\mathbf{e}|\mathbf{a},s',\mathbf{b}) P(s'|\mathbf{a},\mathbf{b})$$
$$= \sum_{s'} P(\mathbf{e}|s') P(s'|\mathbf{a},\mathbf{b}) \quad \text{markov assumption}$$
$$= \sum_{s'} P(\mathbf{e}|s') \sum_s P(s'|s,\mathbf{a}) b(s)$$
- Probability of reaching \mathbf{b}' from \mathbf{b} , given action \mathbf{a} not knowing \mathbf{e}
$$P(\mathbf{b}'|\mathbf{a},\mathbf{b}) = \sum_e P(\mathbf{b}'|\mathbf{e},\mathbf{a},\mathbf{b}) P(\mathbf{e}|\mathbf{a},\mathbf{b})$$
$$= \sum_e P(\mathbf{b}'|\mathbf{e},\mathbf{a},\mathbf{b}) \sum_{s'} P(\mathbf{e}|s') \sum_s P(s'|s,\mathbf{a}) b(s)$$

Where $P(\mathbf{b}'|\mathbf{e},\mathbf{a},\mathbf{b}) = 1$ if $\text{FORWARD}(\mathbf{b}, \mathbf{a}, \mathbf{e}) = \mathbf{b}'$ and $P(\mathbf{b}'|\mathbf{b}, \mathbf{a}, \mathbf{e}) = 0$ otherwise
- A new reward function for belief states: $\rho(\mathbf{b}) = \sum_s b(s)R(s)$
- $P(\mathbf{b}'|\mathbf{b},\mathbf{a})$ and $\rho(\mathbf{b})$ define an *observable* MDP on the space of belief states.

Belief MDP

- A belief MDP is a tuple $\langle B, A, \rho, E \rangle$:

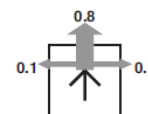
B = infinite set of belief states

E = percepts

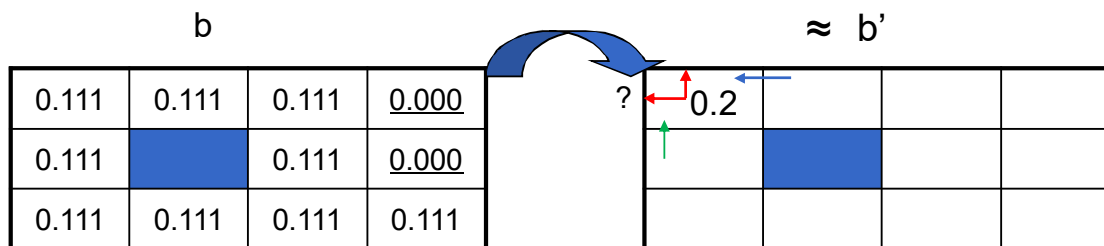
A = finite set of actions

$$\rho(b) = \sum_s b(s)R(s) \quad \text{(reward function)}$$

$$P(b'|b, a) = \sum_e P(b'|e, a, b) \sum_{s'} P(e|s') \sum_s P(s'|s, a) b(s) \quad \text{(transition function)}$$



Move left once, without observations



$$0.8 \cdot 0.111 + 0.1 \cdot 0.111 + 0.1 \cdot 0.111 + 0.8 \cdot 0.111 = 0.2$$

Belief MDP

- A belief MDP is a tuple $\langle B, A, \rho, E \rangle$:

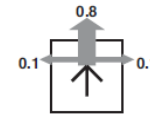
B = infinite set of belief states

E = percepts

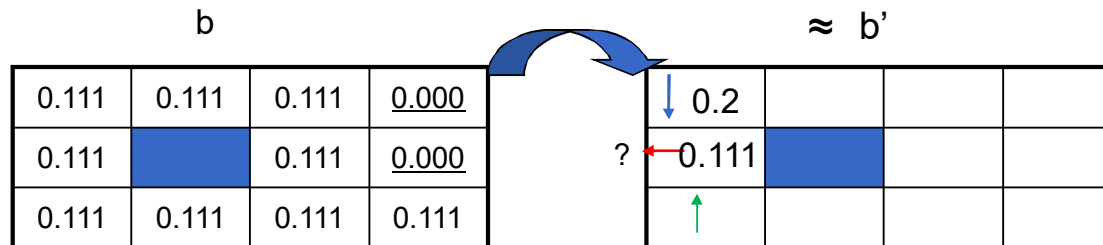
A = finite set of actions

$$\rho(b) = \sum_s b(s)R(s) \quad (\text{reward function})$$

$$P(b'|b, a) = \sum_e P(b'|e, a, b) \sum_{s'} P(e|s') \sum_s P(s'|s, a) b(s) \quad (\text{transition function})$$



Move left once, without observations

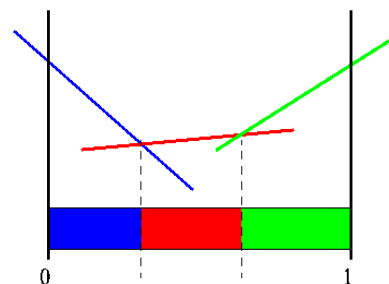


$$0.8 \cdot 0.111 + 0.1 \cdot 0.111 + 0.1 \cdot 0.111 = 0.111$$

Solutions for POMDP

- Methods based on *value* and *policy iteration*:

A policy $\pi(b)$ can be represented as a set of *regions* of belief state space, each of which is associated with a particular optimal action. The value function associates a distinct *linear* function of b with each region. Each value or policy iteration step refines the boundaries of the regions and may introduce new regions.



Value Iteration for POMDPS

- Consider an optimal policy π^* and its application in **belief state \mathbf{b}** .
- For this \mathbf{b} the policy is a “conditional plan”
 - ♦ Let the utility of executing a fixed conditional plan \mathbf{p} in \mathbf{s} be $u_p(\mathbf{s})$.
Expected utility $U_p(\mathbf{b}) = \sum_s \mathbf{b}(s) u_p(\mathbf{s})$
It varies linearly with \mathbf{b} , a hyperplane in a belief space
 - ♦ At any \mathbf{b} , the optimal policy will choose the conditional plan with the highest expected utility
 $U(\mathbf{b}) = U^{\pi^*}(\mathbf{b}) = \operatorname{argmax}_p \mathbf{b} \times u_p$ (summation of dot-prod.)
- $U(\mathbf{b})$ is the maximum of a collection of hyperplanes and will be piecewise linear and convex

19

Example: Conditional Plans

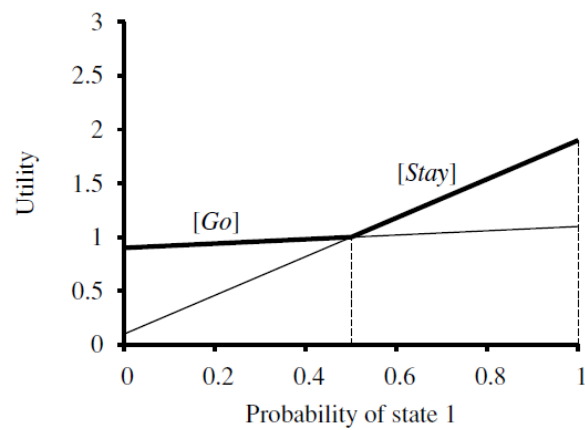
- Two state world 0,1
- Two actions: stay(s), go(s)
 - ♦ Actions achieve intended effect with some probability p
- One-step plan [go], [stay]
- Two-step plans are conditional
 - ♦ [a1, IF percept = 0 THEN a2 ELSE a3]
 - ♦ Shorthand notation: [a1, a2/a3]
- n-step plan are trees with nodes attached with actions and edges attached with percepts

Example

- Two state world 0,1. $R(0)=0$, $R(1)=1$
- Two actions: stay (0.9), go (0.9)
- The sensor reports the correct state with prob. 0.6
- Consider the one-step plans [stay] and [go]
 - ♦ $u_{[\text{stay}]}(0)=R(0) + 0.9R(0)+0.1R(1) = 0.1$
 - ♦ $u_{[\text{stay}]}(1)=R(1) + 0.9R(1)+0.1R(0) = 1.9$
 - ♦ $u_{[\text{go}]}(0)=R(0) + 0.9R(1)+0.1R(0) = 0.9$
 - ♦ $u_{[\text{go}]}(1)=R(1) + 0.9R(0)+0.1R(1) = 1.1$
- This is just the direct reward function (taken into account the probabilistic transitions)

21

Example



Utility of two one-step plans
as a function of $b(1)$

if $(b(1) > 0.5)$ stay else go

General formula

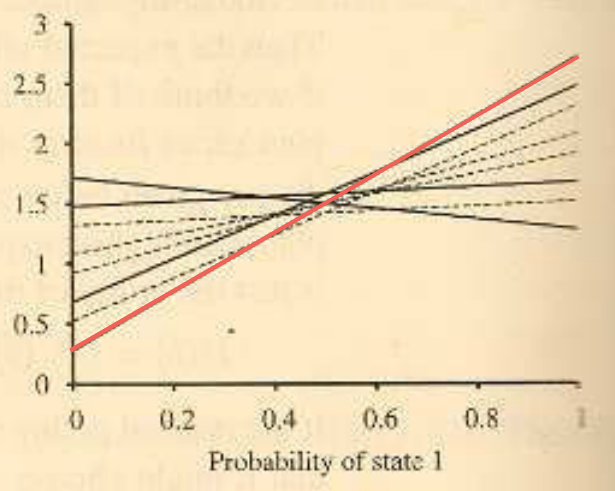
We can compute the utilities for conditional plans of depth-2 by considering each possible first action, each possible subsequent percept and then each way of choosing a depth-1 plan to execute for each percept

[Stay; **if** Percept =0 **then** Stay **else** Stay]
[Stay; **if** Percept =0 **then** Stay **else** Go] . . .

- Let p be a depth- d conditional plan whose initial action is a and whose depth- $(d-1)$ subplan for percept e is $p.e$, then

$$u_p(s) = R(s) + \sum_{s'} P(s' | s, a) \sum_e P(e | s') u_{p.e}(s')$$

Example



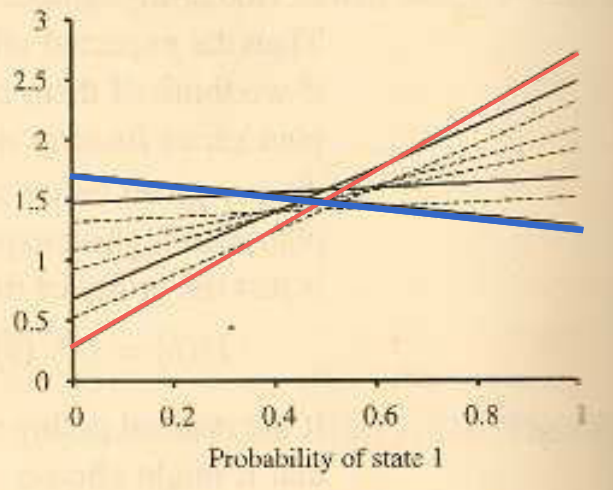
- $U_{[stay]}(0) = R(0) + 0.9R(0) + 0.1R(1) = 0.1$
- $U_{[stay]}(1) = R(1) + 0.9R(1) + 0.1R(0) = 1.9$
- $U_{[go]}(0) = R(0) + 0.9R(1) + 0.1R(0) = 0.9$
- $U_{[go]}(1) = R(1) + 0.9R(0) + 0.1R(1) = 1.1$

$$u_p(s) = R(s) + \sum_{s'} P(s'|s, a) \sum_e P(e|s') u_{p,e}(s')$$

$$u_{[stay, stay/stay]}(0) = R(0) + \underbrace{(0.9 * (0.6 * 0.1 + 0.4 * 0.1))}_{\text{use } u_{stay}(0)} + \underbrace{0.1 * (0.4 * 1.9 + 0.6 * 1.9)}_{\text{use } u_{stay}(1)} = 0.28$$

$$u_{[stay, stay/stay]}(1) = R(1) + \underbrace{(0.1 * (0.6 * 0.1 + 0.4 * 0.1))}_{u_{stay}(0)} + \underbrace{0.9 * (0.4 * 1.9 + 0.6 * 1.9)}_{u_{stay}(1)} = 2.72$$

Example



- ♦ $u_{[stay]}(0) = R(0) + 0.9R(0) + 0.1R(1) = 0.1$
- ♦ $u_{[stay]}(1) = R(1) + 0.9R(1) + 0.1R(0) = 1.9$
- ♦ $u_{[go]}(0) = R(0) + 0.9R(1) + 0.1R(0) = 0.9$
- ♦ $u_{[go]}(1) = R(1) + 0.9R(0) + 0.1R(1) = 1.1$

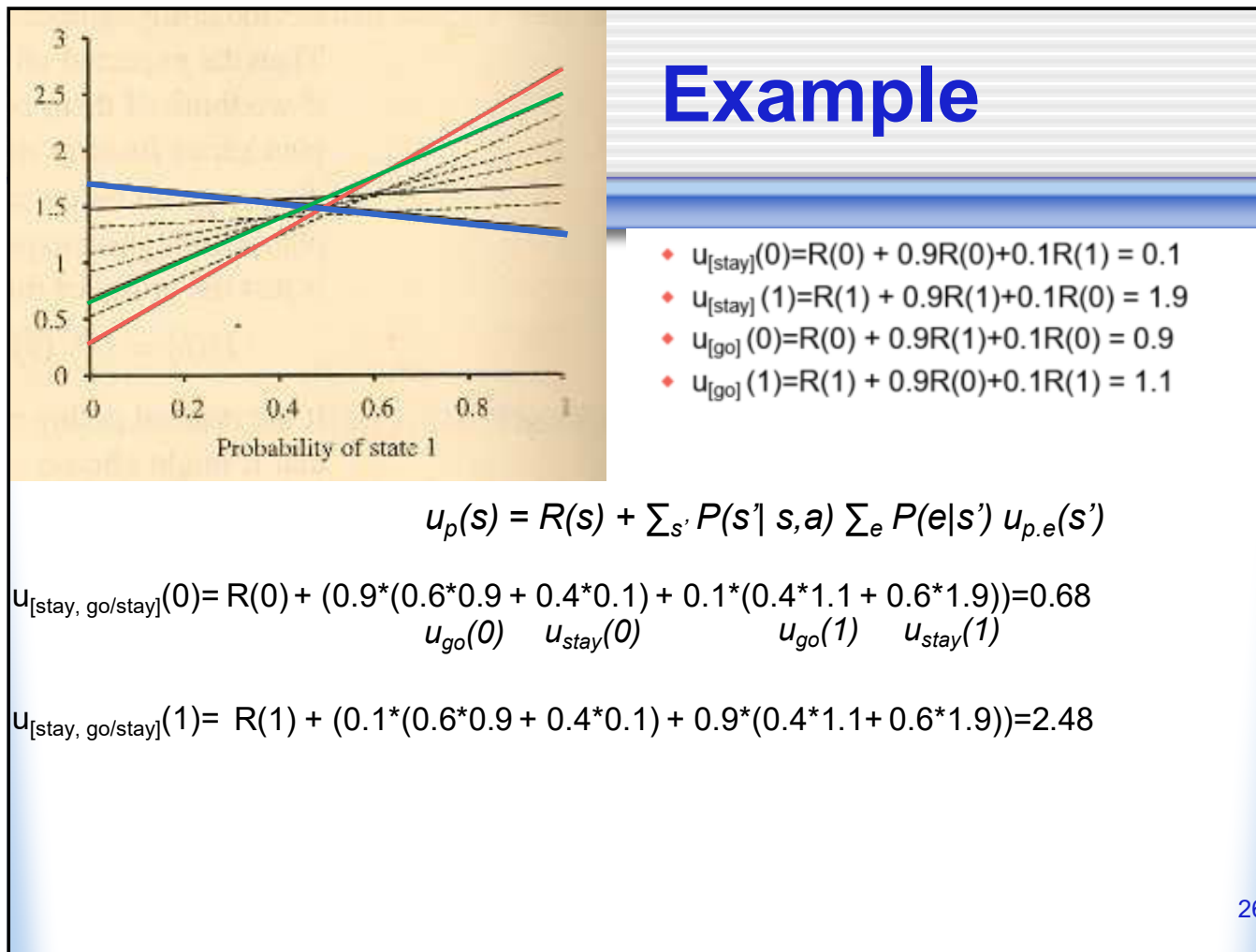
$$u_p(s) = R(s) + \sum_{s'} P(s' | s, a) \sum_e P(e | s') u_{p,e}(s')$$

$$u_{[go, stay/stay]}(0) = R(0) + (0.1 * (0.6 * 0.1 + 0.4 * 0.1) + 0.9 * (0.4 * 1.9 + 0.6 * 1.9)) = 1.72$$

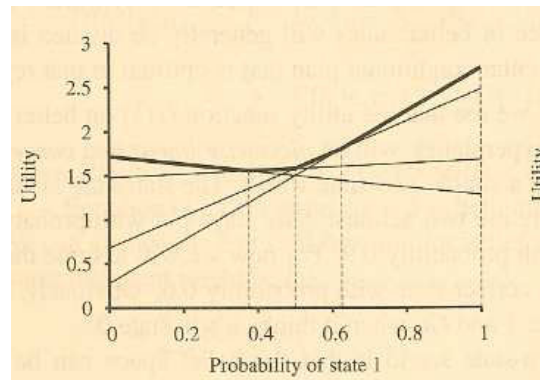
$$u_{[go, stay/stay]}(1) = R(1) + (0.9 * (0.6 * 0.1 + 0.4 * 0.1) + 0.1 * (0.4 * 1.9 + 0.6 * 1.9)) = 1.28$$

$$\underbrace{\hspace{10em}}_{u_{stay}(0)}$$

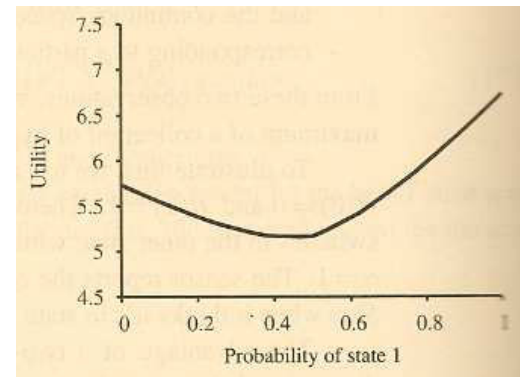
$$\underbrace{\hspace{10em}}_{u_{stay}(1)}$$



Example



Utility of four undominated two-step plans



Utility function for optimal eight step plans

Value Iteration

$$u_p(s) = R(s) + \sum_{s'} P(s'|s,a) \sum_e P(e|s') u_{p,e}(s')$$

- This give us a value iteration algorithm
- The elimination of dominated plans is essential for reducing doubly exponential growth:
the number of undominated plans with $d=8$ is just 144,
otherwise 2^{255} ($|A|^{O(|E|^{d-1})}$)
If you have n undominated plans you have to generate $|A| * n^{|E|}$ new plans.
- For large POMDPs this approach is highly inefficient

Model for POMDPs

- Dynamic Bayesian network
 - ♦ the transition and observation models
- Dynamic decision network (DDN)
 - ♦ decision and utility
- A filtering algorithm
 - ♦ incorporate each new percept and action and update the belief state representation.
- Decisions are made by projecting forward possible action sequences and choosing the best action sequence.

The Generic Structure of a Dynamic Decision Network

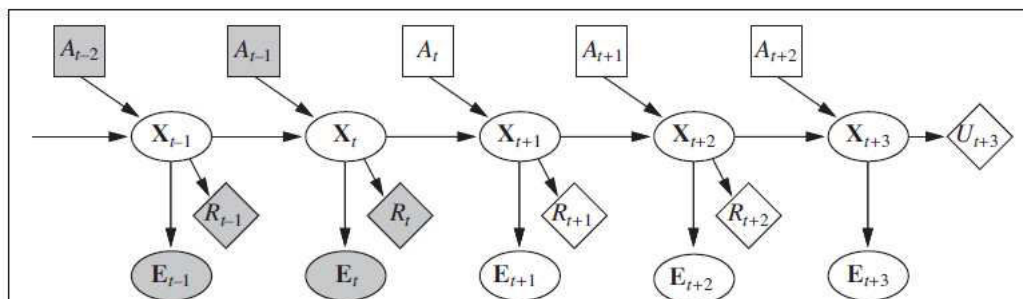
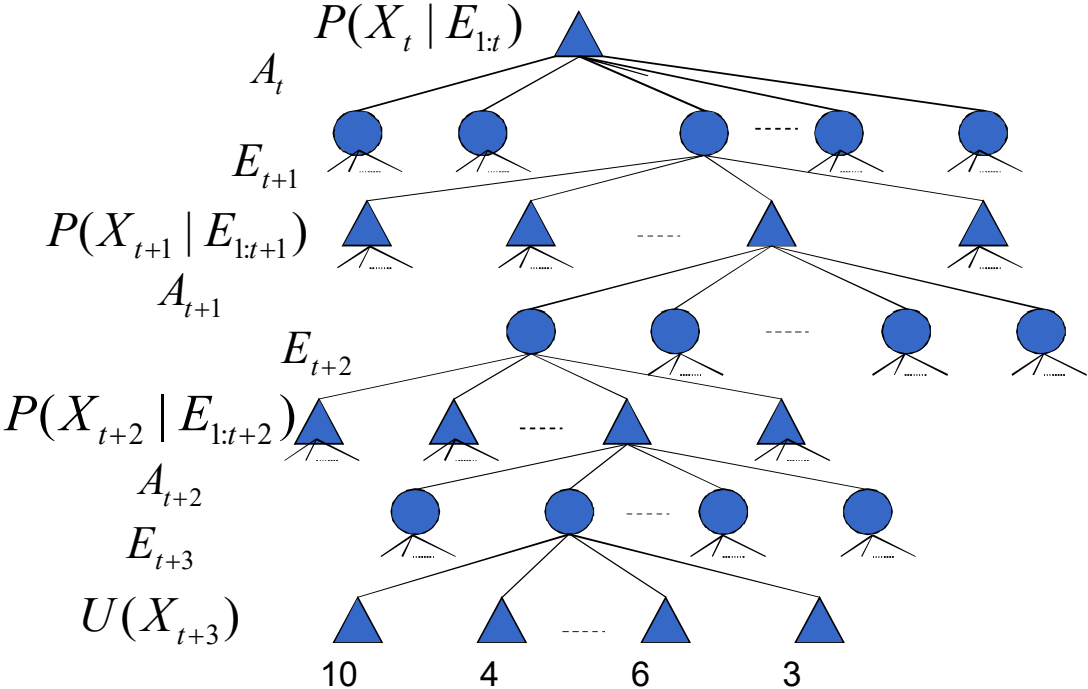


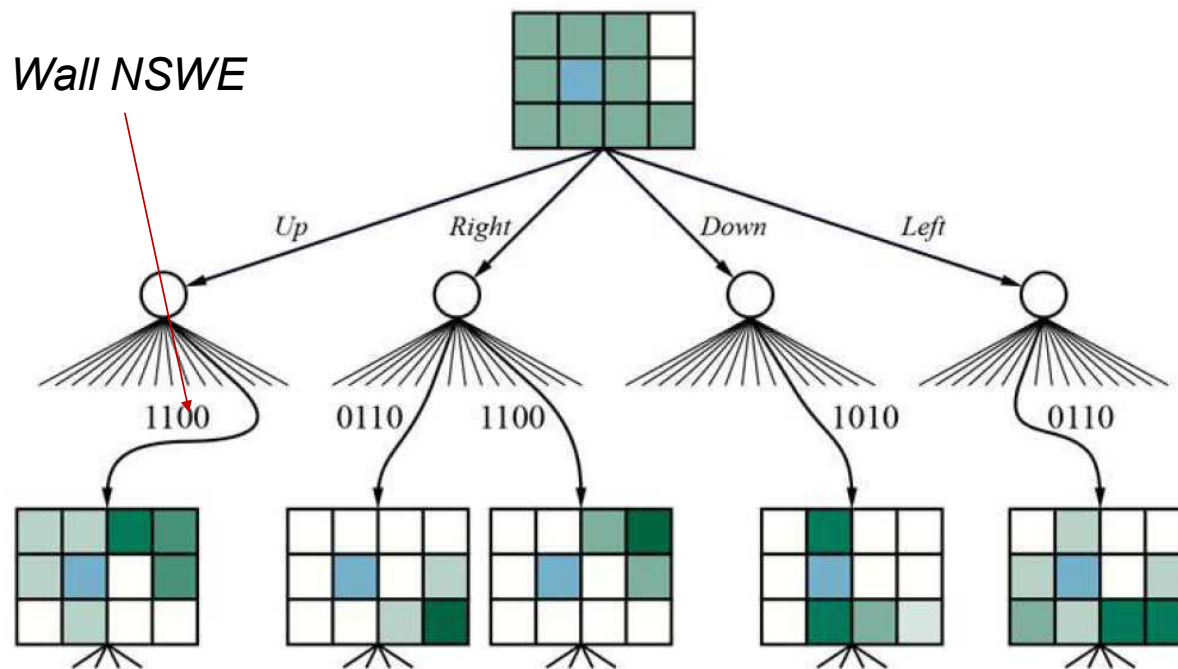
Figure 17.10 FILES: figures/generic-ddn.eps (Tue Nov 3 16:22:53 2009). The generic structure of a dynamic decision network. Variables with known values are shaded. The current time is t and the agent must decide what to do—that is, choose a value for A_t . The network has been unrolled into the future for three steps and represents future rewards, as well as the utility of the state at the look-ahead horizon.

- The decision problem involves calculating the value of A_t that maximizes the agent's expected utility over the remaining state sequence.

Search Tree of the Lookahead DDN



Search Tree of the Lookahead DDN



Search Tree: Exhaustive Enumeration

- The search tree of DDN is very similar to the EXPECTIMINIMAX algorithm for game trees with chance nodes, except that:
 - ♦ There can also be rewards at non-leaf states
 - ♦ The decision nodes correspond to belief states rather than actual states.
- The time complexity: $O(|A|^d \cdot |E|^d)$
 d is the depth, $|A|$ is the number of available actions, $|E|$ is the number of possible observations.
This is far less than value iteration.

33

Discussion of DDNs

- DDNs provide a **general, concise representation** for large POMDPs
- Agent systems moved from
 - ♦ **static, accessible, and simple** environments to
 - ♦ **dynamic, inaccessible, and complex** environments that are closer to the real world
- However, **exact algorithms** are **exponential**

Perspectives of DDNs to Reduce Complexity

- **Heuristic estimate** for the utility of the remaining steps
- Incremental **pruning** techniques
- Many **approximation** techniques as in our search lecture:
 - ♦ Using less detailed state variables for states in the distant future.
 - ♦ Using a greedy heuristic search through the space of decision sequences.

...

Intelligent Autonomous Agents and Cognitive Robotics

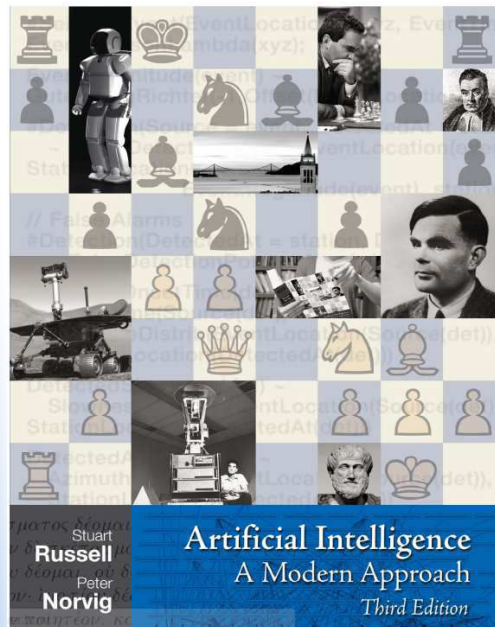
Topic 10: AgentS and Game Theory

Topic 11: Social Choice (Preference Aggregation)

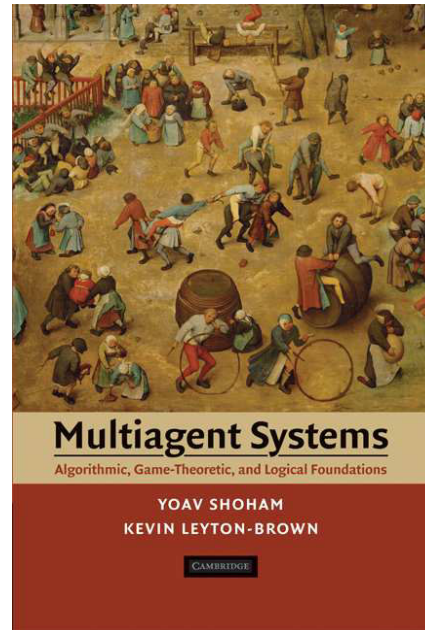
Ralf Möller, Rainer Marrone
Hamburg University of Technology

Literature

- Chapter 17



- Chapter 3



Game Theory

- So far we looked at uncertainty of *actions* and *sensors*
- Now, uncertainty due to the **behavior** of other *agents* !!!!!

→ Game theory

Game Theory: The Basics

- **A game:** Formal representation of a situation of *strategic interdependence* (extension of Adversarial search)
 - ♦ Set of agents, I ($|I|=n$)
 - AKA players
 - ♦ Each agent, j , has a set of actions, A_j
 - AKA moves
 - ♦ Actions define outcomes
 - For each possible action there is an outcome \rightarrow state.
 - ♦ Outcomes define payoffs
 - Agents' derive utility from different outcomes. Utilities can be the same or different.

4

Normal form game*

(matching pennies)

		Agent 2	
		H	T
Agent 1	Action H	-1, 1	1, -1
	T	1, -1	-1, 1

Outcome

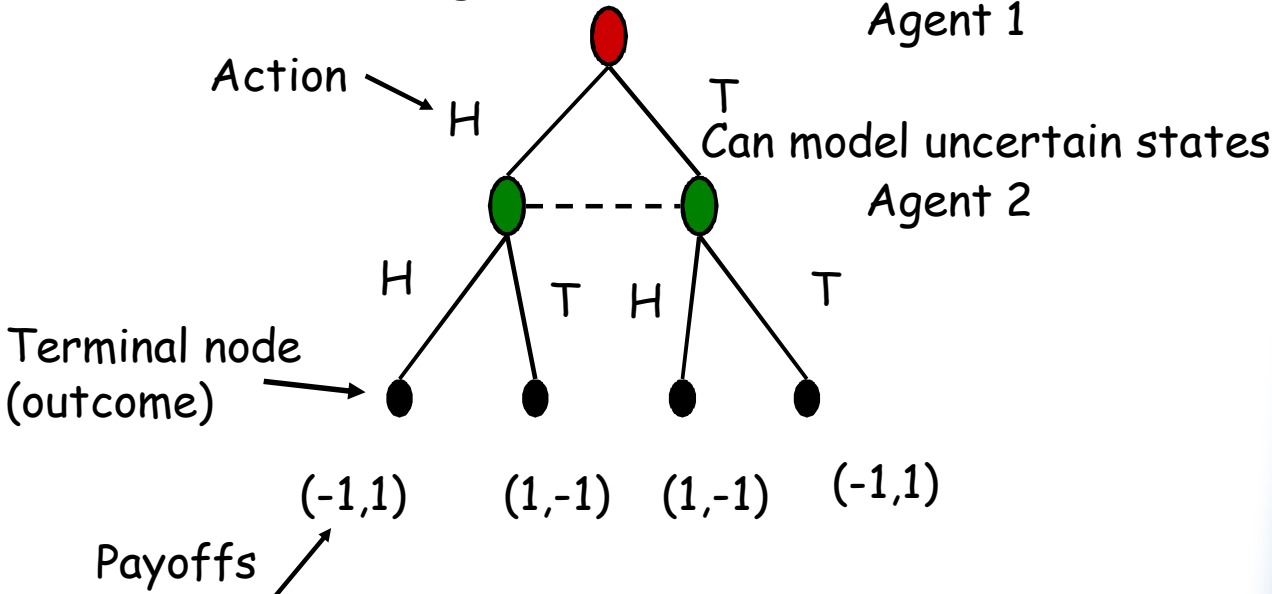
Payoffs

*aka strategic form, matrix form

Extensive form game

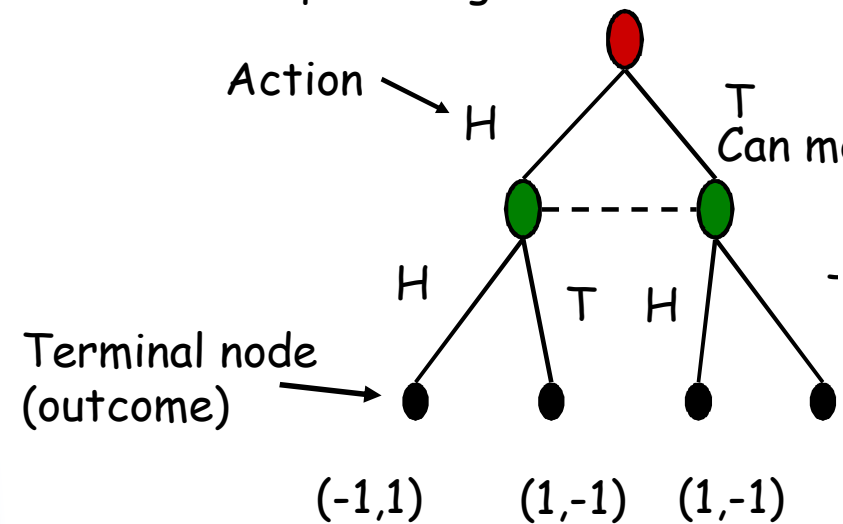
(matching pennies)

Can model sequential games



Extensive form game (matching pennies)

Can model sequential games



Agent 1

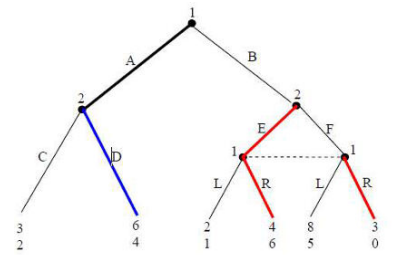
Can model uncertain states

Agent 2

Terminal node
(outcome)

Payoffs

Can also model partly sequential & partly parallel situations



Strategies (aka Policies)

- Strategy:
 - ♦ A strategy, s_j , is a **complete contingency plan** (policy); defines actions agent j should take for all possible states of the world
- Strategy profile:
 - ♦ $\mathbf{s} = (s_1, \dots, s_n)$ (all agents)
 - ♦ $\mathbf{s}_{-i} = (s_1, \dots, s_{i-1}, s_{i+1}, \dots, s_n)$ (all agents without i)
- Utility function: $u_i(\mathbf{s})$
 - ♦ Note that the utility of an agent depends on the strategy profile, not just its own strategy
 - ♦ We assume agents are **expected utility maximizers**

8

Normal form game*

(matching pennies)

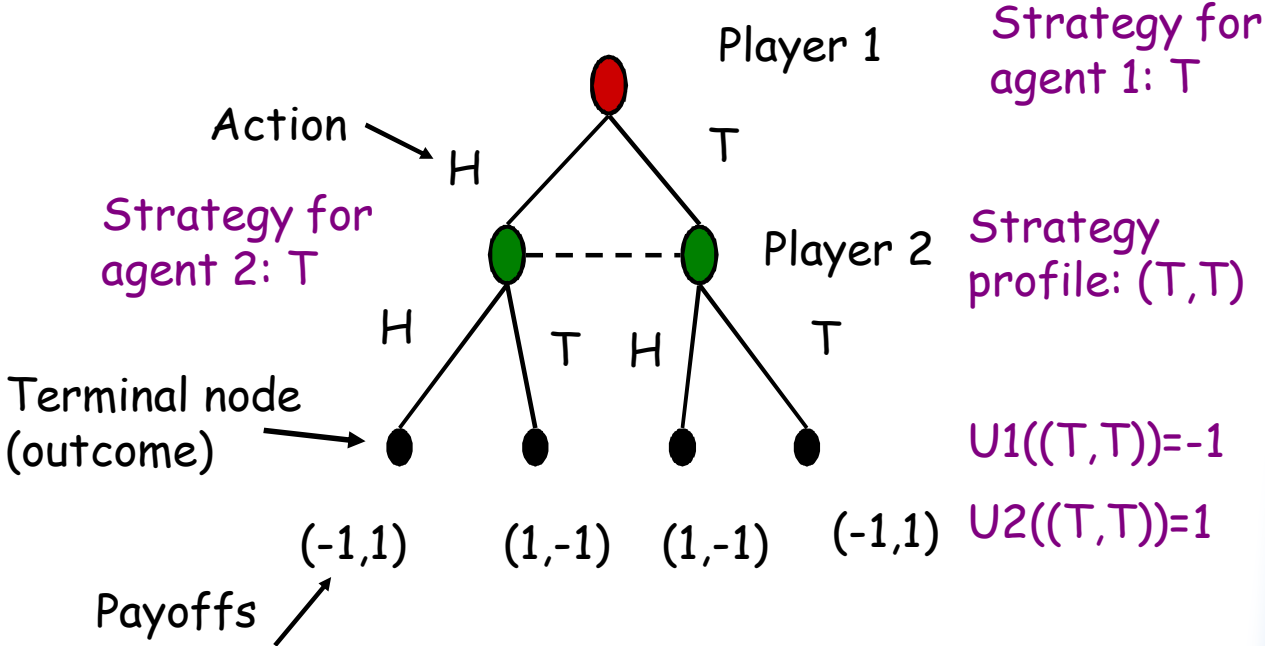
		Agent 2		Strategy for agent 2: T
		H	T	
Agent 1	H	-1, 1	1, -1	Strategy profile (H,T)
	T	1, -1	-1, 1	

$U_1((H, T)) = 1$
 $U_2((H, T)) = -1$

*aka strategic form, matrix form

Extensive form game

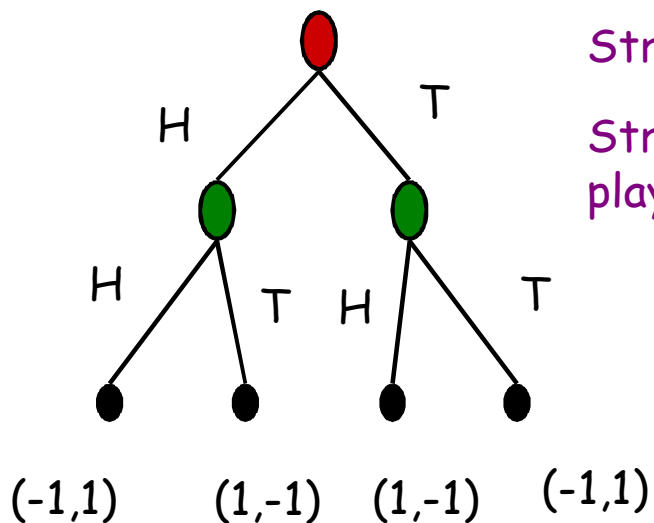
(matching pennies)



Extensive form game

(matching pennies, seq moves)

Recall: A strategy is a contingency plan for all states of the game



Strategy for agent 1: T

Strategy for agent 2: H if 1 plays H, T if 1 plays T (H,T)

Strategy profile: (T,(H,T))

$U_1((T,(H,T))) = -1$

$U_2((T,(H,T))) = 1$

Dominant Strategies

- Recall that
 - ♦ Agents' utilities depend on what strategies other agents are playing
 - ♦ Agents' are expected utility maximizers
- Agents' will play best-response strategies for s_{-i}
 s_i^* is a best response if $u_i(s_i^*, s_{-i}) \geq u_i(s_i', s_{-i})$ for all s_i'
- A dominant strategy is a best-response for all s_{-i}
 - ♦ They do not always exist
 - ♦ Inferior strategies are called dominated

Dominant Strategy Equilibrium

- A *dominant strategy equilibrium* is a strategy profile where the strategy for each player is dominant
 - ♦ $s^* = (s_1^*, \dots, s_n^*)$
 - ♦ $u_i(s_i^*, s_{-i}) \geq u_i(s_i', s_{-i})$ for all i , for all s_i' , for all s_{-i}
- **GOOD:**
Agents do not need to counterspeculate!

Example: Prisoner's Dilemma

- Two people are arrested for a crime. A prosecutor offers each a deal: if you **testify** against your partner as the leader of a burglary ring, you'll go free for being the cooperative one, while your partner will serve 10 years in prison. However, if both **testify** against each other, they both get 5 years. If both **refuse**, each get 1 year.

		A: testify	A: refuse
Dom. Str. Eq	B: testify	$B = -5$ $A = -5$	$B = 0$ $A = -10$
	B: refuse	$B = -10$ $A = 0$	$B = -1$ $A = -1$

Pareto Optimal Outcome

14

Example: Bach or Stravinsky

aka Battle of sexes

- A couple likes going to two concerts. One loves Bach but not Stravinsky. The other loves Stravinsky but not Bach. However, they prefer being together than being apart.

	B	S
B	2,1	0,0
S	0,0	1,2

No dom.
str. equil.

Nash Equilibrium

- Sometimes an agent's best-response depends on the strategies other agents are playing
 - ♦ No dominant strategy equilibria
- A strategy profile is a **Nash equilibrium** if no player has incentive to deviate from his strategy **given that others do not deviate**:
 - ♦ for every agent i , $u_i(s_i^*, s_{-i}) \geq u_i(s_i', s_{-i}) \rightarrow s_i^*$ is a best response to s_{-i}

	B	S
B	2,1	0,0
S	0,0	1,2

16

How to find (Nash) Equilibria

- Can agents rule out strategies?
 - ◆ Strategies an agent will not play
- Get rid of those strategies
 - ◆ Maybe there will exist a single solution

Example

	r	l	c
U	3,-3	7,-7	9,-15
D	9,-9	8,-8	10,-10

18

Iterated Elimination of Dominated Strategies

- Let $R_i \subseteq S_i$ be the set of removed strategies for agent i
- Initially $R_i = \emptyset$
- Choose agent i , and strategy s_i such that $s_i \in S_i \setminus R_i$ and there exists $s_i' \in S_i \setminus R_i$ such that

$$u_i(s_i', s_{-i}) > u_i(s_i, s_{-i}) \text{ for all } s_{-i} \in S_{-i} \setminus R_{-i}$$
- Add s_i to R_i , continue
- **Thm:** If a unique strategy profile, s^* , survives then it is a Nash Eq.
- **Thm:** If a profile, s^* , is a Nash Equilibrium then it must survive iterated elimination.

Nash Equilibrium

- Criticisms
 - ◆ They may not be unique (Bach or Stravinsky)
 - Ways of overcoming this
 - Refinements of equilibrium concept, Mediation, Learning
 - ◆ Do not exist in all games (in form defined)
 - ◆ They may be hard to find
 - ◆ People don't always behave based on what equilibria would predict (ultimatum games and notions of fairness,...)

Example: Matching Pennies

	H	T
H	-1, 1	1, -1
T	1, -1	-1, 1

There is NO (Nash) strategy in pure strategies

Example: Bach Stravinsky

	B	S
B	2,1	0,0
S	0,0	1,2

If I do not know, what the other agent is doing, and if communication is not possible, what should the agents do

So far we have talked only about **pure** strategy equilibria.

Not all games have pure strategy equilibria. Some equilibria are **mixed** strategy equilibria.

22

Example: Bach Stravinski

		<i>Husband</i>	
		<i>q</i> B	<i>1-q</i> S
<i>Wife</i>	<i>p</i> B	2, 1	0, 0
	<i>1-p</i> S	0, 0	1, 2

Mixed strategies can help if no communication is possible.
 Want to play a strategy so that the other is indifferent
 playing a pure strategy (B or S).

$$EU_{HB} = 1p + 0(1-p)$$

$$EU_{HS} = 0p + 2(1-p)$$

$$EU_{HB} = EU_{HS}$$

$$p = 2 - 2p$$

$$p = 2/3 \quad (\text{wife has mixed } \langle 2/3; 1/3 \rangle)$$

Example: Bach Stravinski

		<i>Husband</i>	
		q B	$1-q$ S
<i>Wife</i>	p B	2, 1	0, 0
	$1-p$ S	0, 0	1, 2

Mixed strategies can help if no communication is possible.
 Want to play a strategy so that the other is indifferent
 playing a pure strategy (B or S).

$$EU_{WB} = 2q + 0(1-q)$$

$$EU_{WS} = 0q + 1(1-q)$$

$$EU_{WB} = EU_{WS}$$

$$2q = 1 - 1q$$

$$q = 1/3 \quad (\text{husband has mixed } \langle 1/3; 2/3 \rangle)$$

Example: Bach Stravinski

- If Husband **strictly** plays B with $q=1/3$
 - ♦ Which distribution can his wife play
 - ♦ $EU_w(p,1-p) = 1/3 * p * 2 + 2/3 * p * 0 + 1/3 * (1-p) * 0 + 2/3 * (1-p) * 1 = 2/3 * p + 2/3 - 2/3 p = 2/3$
any distribution leads to 2/3 in average

		1/3 B	2/3 S
p	B	2, 1	0, 0
1-p	S	0, 0	1, 2

Example: Bach Stravinski

		<i>Husband</i>	
		q B	$1-q$ S
<i>Wife</i>	p B	2, 1	0, 0
	$1-p$ S	0, 0	1, 2

$$EU_{WB} = 2q + 0(1-q)$$

$$EU_{WS} = 0q + 1(1-q)$$

$$EU_{WB} = EU_{WS}$$

$$2q = 1 - 1q$$

$$q = 1/3$$

husband has mixed strategy $\langle 1/3; 2/3 \rangle$

wife has mixed strategy $\langle 2/3; 1/3 \rangle$

Example: Bach Stravinski

- If Husband **strictly** plays B with $q=1/3$
 - ♦ Which distribution should wife play
 - ♦ $Eu_w(p, 1-p) = 2/3$
- If Husband deviates $q < 1/3$
 - ♦ Wife deviates plays **S**
- If Husband $q > 1/3$
 - ♦ Wife plays **B**
- Equilibrium: $\{(2/3, 1/3); (1/3, 2/3)\}$

		Husband	
		q B	$1-q$ S
Wife	p B	2, 1	0, 0
	$1-p$ S	0, 0	1, 2

Mixed strategy equilibria

- Mixed strategy:
 $\sigma_i \in \Sigma_i$ defines a probability distribution over S_i
- Strategy profile: $\sigma = (\sigma_1, \dots, \sigma_n)$
- Expected utility: $u_i(\sigma) = \sum_{s \in S} (\prod_j \sigma(s_j)) u_i(s)$
- Nash Equilibrium:
 - ♦ σ^* is a mixed Nash equilibrium if
 $u_i(\sigma^*_i, \sigma^*_{-i}) \geq u_i(\sigma_i, \sigma^*_{-i})$ for all $\sigma_i \in \Sigma_i$, for all i

28

Mixed Nash Equilibrium

- Thm (Nash 50):
 - ◆ Every game in which the strategy sets S_1, \dots, S_n have a finite number of elements, has a **mixed strategy equilibrium**.
- Finding Nash Equilibria is another problem
 - ◆ “Together with prime factoring, the complexity of finding a Nash Eq is the most important concrete open question ...” (Papadimitriou)

Bayesian-Nash Equil

(Harsanyi 68)

- So far we have assumed that agents have complete information about each other (including payoffs)
 - ♦ **Very strong assumption!**
- Assume agent i has **type** $\theta_i \in \Theta_i$, defines the payoff $u_i(s, \theta_i)$
- Agents have common prior over distribution of types $p(\theta)$
 - ♦ Conditional probability $p(\theta_{-i} | \theta_i)$
(obtained by Bayes Rule when possible)

30

Battle of the sexes

- Shopping or Basketball?
- Sally knows Kevin's type
Kevin does not know Sally's type but possible types.

What should Sally play?

		Kevin	
		Basketball	Shopping
Sally	Basketball	3, 2	2, 1
	Shopping	0, 0	1, 3

Sally a basketball fan

Her **dominant** strategy!

		Kevin	
		Basketball	Shopping
Sally	Basketball	1, 2	0, 1
	Shopping	2, 0	3, 3

Sally a shopping fan

Battle of the sexes

- Sally should play her **dominant** strategy

$$\Theta_1 = \{\theta_{11}, \theta_{12}\}, \Theta_2 = \{\theta_2\},$$

$$P(\theta_{11}, \theta_2) = p \quad P(\theta_{12}, \theta_2) = (1-p)$$

$$2p + 0(1-p) > 1p + 3(1-p) \quad \text{basketball vs shopping}$$

$$2p > -2p + 3$$

$$p > 3/4$$

If $p > 3/4$ Basketball

If $p < 3/4$ Shopping

If $p = 3/4$??

		Kevin	
		Basketball	Shopping
Sally	Basketball	3, 2	2, 1
	Shopping	0, 0	1, 3

		Kevin	
		Basketball	Shopping
Sally	Basketball	1, 2	0, 1
	Shopping	2 , 0	3, 3

32

Battle of the sexes

- Sally's decision depends on her known type
- Kevin's decision depends on p

$$S_2^*(\theta_2) = \begin{cases} \text{Basketball} & \text{if } p > 3/4 \\ (q, 1-q), q \in [0, 1] & \text{if } p = 3/4 \\ \text{Shopping} & \text{if } p < 3/4 \end{cases}$$

Bayesian-Nash Equil

- **Strategy:** $\sigma_i(\theta_i)$ is the (mixed) strategy agent i plays if its type is θ_i .
- **Strategy profile:** $\sigma = (\sigma_1, \dots, \sigma_n)$
- **Expected utility:**
 - ♦ $U_i(\sigma_i(\theta_i), \sigma_{-i}(\cdot), \theta_i) = \sum_{\theta_{-i}} p(\theta_{-i} | \theta_i) u_i(\sigma_i(\theta_i), \sigma_{-i}(\theta_{-i}), \theta_i)$
- **Bayesian Nash Eq:** Strategy profile σ^* is a Bayesian-Nash Equilibrium if for all i , for all θ_i ,
 $U_i(\sigma_i^*(\theta_i), \sigma_{-i}^*(\cdot), \theta_i) \geq U_i(\sigma_i(\theta_i), \sigma_{-i}^*(\cdot), \theta_i)$

(best responding w.r.t. its beliefs about the types of the other agents, assuming they are also playing a best response)

Last time

- Definition of games
- Strategies & Strategy profiles

- ◆ Dominant strategy equilibrium

$$u_i(s_i^*, s_{-i}) \geq u_i(s_i', s_{-i}) \quad \forall s_i', \forall s_{-i}, \forall i,$$

- ◆ Nash equilibrium

$$u_i(s_i^*, s_{-i}) \geq u_i(s_i', s_{-i}) \quad \forall s_i', \forall i,$$

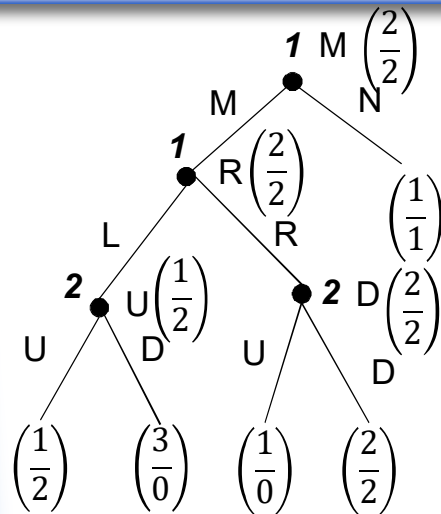
- ◆ Mixed Nash strategy equilibrium

$$u_i(\sigma_i^*, \sigma_{-i}) \geq u_i(\sigma_i', \sigma_{-i}) \quad \forall \sigma_i', \forall i,$$

- ◆ Bayesian Nash equilibrium

$$u_i(\sigma_i^*(\theta_i), \sigma_{-i}^*(\cdot), \theta_i) \geq u_i(\sigma_i(\theta_i), \sigma_{-i}^*(\cdot), \theta_i)$$

Extensive Form Games



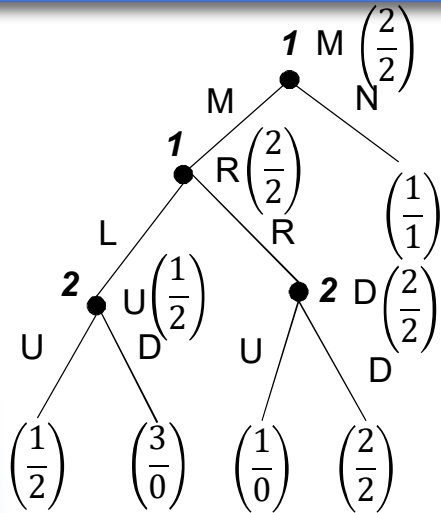
Any finite game of perfect information has a pure strategy Nash equilibrium.
It can be found by backward induction.

How to find a Nash Equilibrium?
By backward induction!
Have to define an action for every choice point.
(MR, UD)

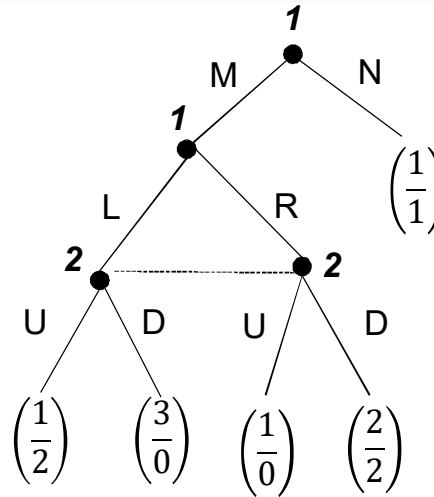
Subgame perfect equilibrium & credible threats

- Proper subgame = subtree (of the game tree) whose root is alone in its information set (agent knows his state)
- Subgame perfect equilibrium
 - ◆ Strategy profile that is in Nash equilibrium in every proper subgame (including the root), whether or not that subgame is reached along the equilibrium path of play

Subgame perfect equilibrium

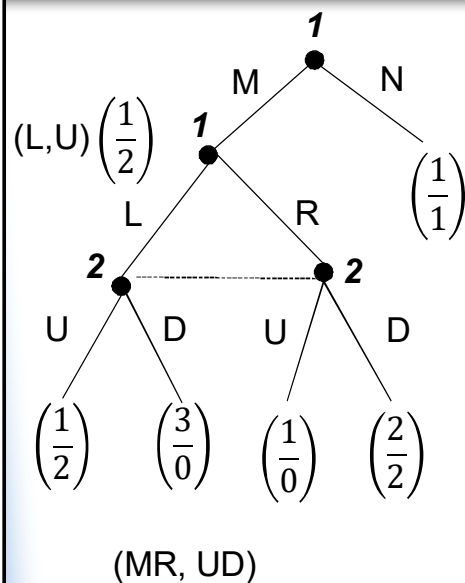


(MR, UD)



What is the strategy now?

Subgame perfect NE equilibrium



What is the strategy now?

	U	D
L	<u>1</u> , <u>2</u>	<u>3</u> , <u>0</u>
R	<u>1</u> , <u>0</u>	<u>2</u> , <u>2</u>

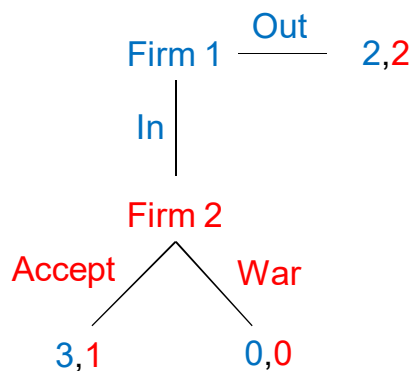
If U best is L or R
 If D best is L
 If L best is U
 If R best is D
 (L,U) is the NE

(ML, U) and (NL, U) are the SPNE of the game
 $\begin{pmatrix} 1 \\ 2 \end{pmatrix}$ $\begin{pmatrix} 1 \\ 1 \end{pmatrix}$

Non creditable threats

- A firm is deciding whether to enter the market, which another firm currently has a monopoly over.
- If the firm enters, the monopolist chooses whether to accept it or declare a price war.
 - ♦ The firm only wants to enter if the monopolist won't engage in a price war
 - ♦ A price war is unprofitable for the monopolist

Non credible threats



	Accept	War
In	3,1	0,0
Out	2,2	2,2

Firm 2 announce to make a price war if Firm 1 enters.

(out, war) is a Nash equilibria.

But, it is not subgame perfect →
This is a non credible threat

Social choice theory

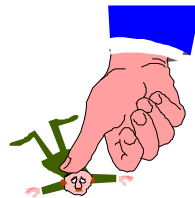
- Study of decision problems in which a group has to make the decision
- The decision affects all members of the group
 - ♦ Their opinions! should count
- Applications:
 - ♦ Political elections
 - ♦ Note that outcomes can be vectors
 - Allocation of money among agents, allocation of goods, tasks, resources...
- CS applications:
 - ♦ Multiagent planning [Ephrati&Rosenschein]
 - ♦ Accepting a joint project, rating Web articles [Avery,Resnick&Zeckhauser]
 - ♦ ...

Criteria for evaluating multiagent systems

- **Social welfare**: $\max_{\text{outcome}} \sum_i u_i(\text{outcome})$
- **Surplus**: social welfare of outcome – social welfare of status quo
 - ♦ Zero sum games have 0 surplus. Markets are not zero sum
- **Pareto efficiency**: An outcome o is Pareto efficient if there exists no other outcome o' s.t. some agent has higher utility in o' than in o and no agent has lower
 - ♦ Implied by social welfare maximization
- **Individual rationality**: Participating in the negotiation (or individual deal) is no worse than not participating
- **Stability**: No agents can increase their utility by changing their strategies
- **Symmetry**: No agent should be inherently preferred, e.g. dictator

Assumptions

1. Agents have preferences over alternatives
 - Agents can **rank order** the outcomes
 - $a > b > c = d$ is read as “a is preferred to b which is preferred to c which is equivalent to d”
2. Voters are **sincere**
 - They truthfully tell the center their preferences
3. Outcome is enforced on all agents



Voting

- Majority decision:
 - ◆ If more agents prefer a to b, then a should be chosen
- Two outcome setting is easy
 - ◆ Choose outcome with more votes!
- What happens if you have 3 or more possible outcomes?

Case 1: Agents specify their top preference

Ballot



47

Election System

- Plurality Voting
 - ◆ One name is ticked on a ballot
 - ◆ One round of voting
 - ◆ One candidate is chosen

Is this a "good"
system?

What do we mean by good?

Example: Plurality (Canada)

- 3 candidates
 - ♦ Lib, NDP, C
- 21 voters with the preferences
 - ♦ 10 Lib>NDP>C
 - ♦ 6 NDP>C>Lib
 - ♦ 5 C>NDP>Lib
- Result: **Lib 10**, NDP 6, C 5
 - ♦ But a majority of voters (11) prefer all other parties more than the Libs!

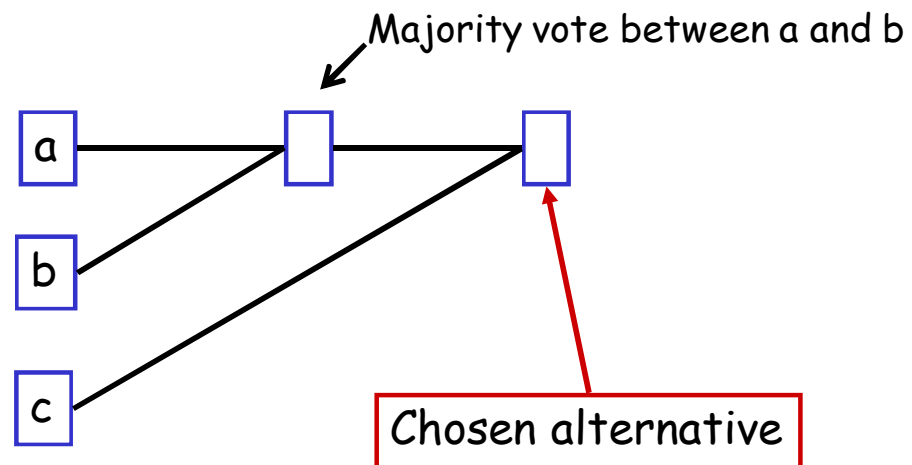
49

What can we do?

- Majority system
 - ♦ Works well when there are 2 alternatives
 - ♦ Not great when there are more than 2 choices
- Proposal:
 - ♦ Organize a series of votes between 2 alternatives at a time
 - ♦ How this is organized is called an **Agenda**
 - Or a cup (often in sports)

Agendas

- 3 alternatives {a,b,c}
- Agenda a,b,c



51

Example: Agenda

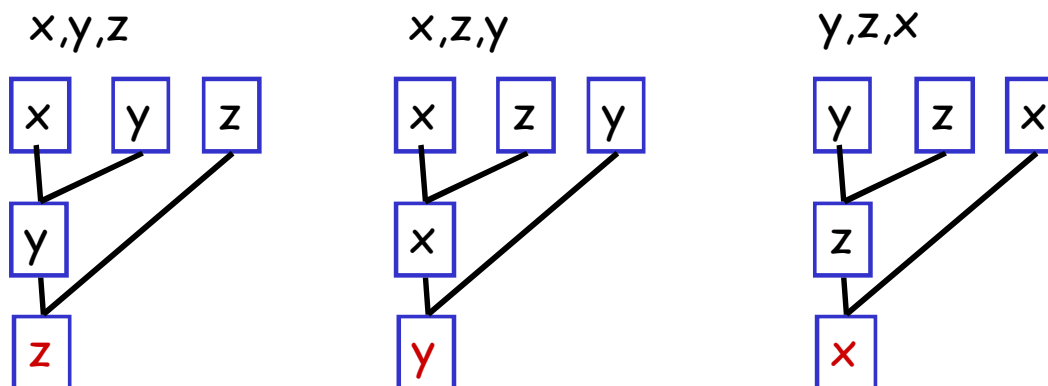
- *Binary protocol (majority rule)* Three types of agents:

1. $x \succ z \succ y$ (35%)

2. $y \succ x \succ z$ (33%)

3. $z \succ y \succ x$ (32%)

Chairman defines order

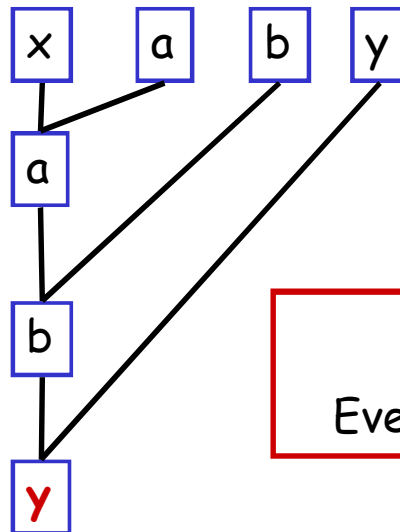


- Power of agenda setter (e.g. chairman)

Pareto dominated winner paradox

Agents:

1. $x > y > b > a$
2. $a > x > y > b$
3. $b > a > x > y$



BUT
Everyone prefers x to y!

53

Case 2: Agents specify their complete preferences

Maybe the
problem was with
the ballots!

Ballot

X>Y>Z



Now have
more
information

Condorcet

Marie Jean Antoine Nicolas Caritat, Marquis de Condorcet

- Proposed the following
 - ◆ Compare each pair of alternatives
 - ◆ Declare “a” is socially preferred to “b” if more voters strictly prefer a to b
- **Condorcet Principle:** If one alternative is preferred to all other candidates then it should be selected

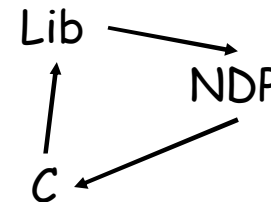
55

Example: Condorcet

- 3 candidates
 - ◆ Lib, NDP, C
- 21 voters with the preferences
 - ◆ 10 Lib>NDP>C
 - ◆ 6 NDP>C>Lib
 - ◆ 5 C>NDP>Lib
- Result:
 - ◆ **NDP win!** (11/21 prefer them to Lib, 16/21 prefer them to C)

A Problem

- 3 candidates
 - ♦ Lib, NDP, C
- 3 voters with the preferences
 - ♦ Lib>NDP>C
 - ♦ NDP>C>Lib
 - ♦ C>Lib>NDP
- Result:
 - ♦ No Condorcet Winner



57

Borda Count

- Each ballot is a list of ordered alternatives
- Over all ballots compute the rank of each alternative
- Rank order alternatives based on decreasing sum of their ranks

A>B>C

A>C>B

C>A>B



A: 4

B: 8

C: 6

Borda Count

- Simple. Only counting ranks
- Always a Borda Winner, but have to define a solution for ties.
- BUT does not always choose Condorcet winner!
Borda scores:
a:5, b:6, c:8, d:11
Therefore **a** wins
BUT **b** is the **Condorcet** winner
- 3 voters
 - ♦ 2: $b > a > c > d$
 - ♦ 1: $a > c > d > b$

59

Another example

- Borda rule with 4 alternatives
- Agents:
 1. $x \succ c \succ b \succ a$
 2. $a \succ x \succ c \succ b$
 3. $b \succ a \succ x \succ c$
 4. $x \succ c \succ b \succ a$
 5. $a \succ x \succ c \succ b$
 6. $b \succ a \succ x \succ c$
 7. $x \succ c \succ b \succ a$
- $x=13$
- $a=18$
- $b=19$
- $c=20$

The winner is dropped

- X went out → Remove x:

1. $x > c > b > a$
2. $a > x > c > b$
3. $b > a > x > c$
4. $x > c > b > a$
5. $a > x > c > b$
6. $b > a > x > c$
7. $x > c > b > a$



1. $c > b > a$
2. $a > c > b$
3. $b > a > c$
4. $c > b > a$
5. $a > c > b$
6. $b > a > c$
7. $c > b > a$

- $x=13$, $a=18$, $b=19$, $c=20$
- $c=13$
- $b=14$
- $a=15$

Inverted order paradox

Borda rule vulnerable to irrelevant alternatives

- Three types of agents:

x	y
35	70
66	33
64	32
165	135
second	first

1. $x > y$ (35%)
2. $y > x$ (33%)
3. $y > x$ (32%)

- Borda winner is y

Borda rule vulnerable to irrelevant alternatives

- Three types of agents:

x	y
35	70
66	33
64	32
165	135
second	first

1. $x > z > y$ (35%)
2. $y > x > z$ (33%)
3. $z > y > x$ (32%)

x	y	z
35	105	70
66	33	99
96	64	32
197	202	201
first	third	second

- Borda winner is y
- Add z Borda winner is x

The social preferences between alternatives x and y depend only on the individual preferences between x and y

64

Desirable properties for a voting protocol

- No dictators
- Universality (unrestricted domain)
 - ◆ It should work with any set of preferences
- Independence of irrelevant alternatives
 - ◆ The comparison of two alternatives should depend only on their standings among agents' preferences, not on the ranking of other alternatives
- Pareto efficient
 - ◆ If all agents prefer x to y then in the outcome x should be preferred to y

Arrow's Theorem (1951)

- **Thrm.** If there are 3 or more alternatives and a finite number of agents then there is **no** protocol which satisfies the 4 desired properties
- **Thrm.** Let $|O| \geq 3$, any social welfare function W that is Pareto efficient and independent of irrelevant alternatives is dictatorial.

Take-home Message

- Despair?
 - ◆ No ideal voting method
 - ◆ That would be boring!
- A group is more complex than an individual
- Weigh the pro's and con's of each system and understand the setting they will be used in
- Do not believe anyone who says they have the best voting system out there!

Intelligent Autonomous Agents and Cognitive Robotics

Topic 12: Mechanism Design

Ralf Möller, Rainer Marrone
Hamburg University of Technology

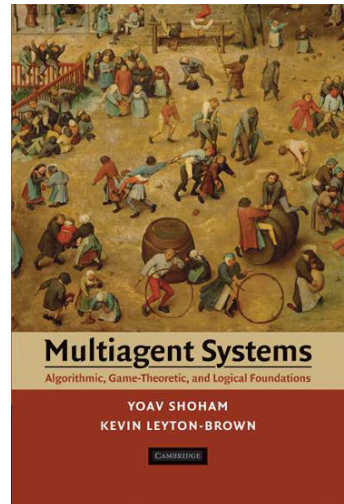
Acknowledgement

Material from CS 886

Advanced Topics in AI [LSEP] **Electronic Market Design**

Kate Larson

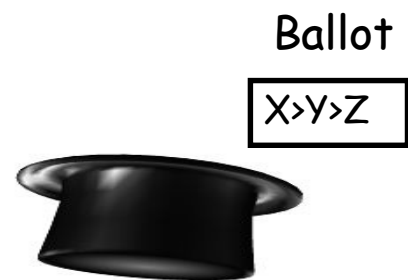
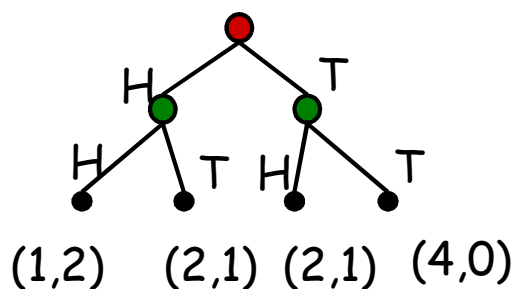
Waterloo Univ.



Introduction

So far we have looked at

- Game Theory
 - ♦ Given a game we are able to analyze the strategies agents will follow
- Social Choice Theory
 - ♦ Given a set of agents' preferences we can choose some outcome



Introduction

- Now: **Mechanism Design**
 - ♦ Game Theory + Social Choice
- Goal of Mechanism Design is to
 - ♦ Obtain a dedicated outcome (function of agents' preferences)
 - ♦ But agents are rational
They may lie about their preferences
- Goal:
Define the rules of a game so that in equilibrium the agents do what the social community in general wants

Fundamentals

- Set of possible outcomes, O
- Agents $i \in I$, $|I|=n$, each agent i has type $\theta_i \in \Theta_i$
 - ♦ Type captures all private information that is relevant to agent's decision making (its payoffs, which may be different)
- Utility $u_i(o, \theta_i)$, over outcome $o \in O$
- Recall: goal is to implement some **system-wide** solution
 - ♦ Captured by a social choice function (SCF)

$$f: \Theta_1 \times \dots \times \Theta_n \rightarrow O$$

$f(\theta_1, \dots, \theta_n) = o$ is a collective choice

Mechanisms

- Recall: We want to implement a social choice function
 - ♦ Need to know agents' preferences
 - ♦ They may not reveal them to us truthfully
- Example:
 - ♦ 1 item to allocate, and want to give it to the agent who values it the most
 - ♦ If we just ask agents to tell us their preferences, they may lie

I like the bear the most!



No, I do!

7

Mechanism Design Problem

- By having agents interact through an institution (M) we might be able to solve the problem
- Mechanism:

$$M = (S_1, \dots, S_n, g(\cdot))$$

↗
↑

Strategy spaces of agents
Outcome function

$$g: S_1 \times \dots \times S_n \rightarrow O$$

Implementation

- A mechanism $M=(S_1, \dots, S_n, g(\cdot))$ **implements** social choice function $f(\theta)$ if there is an equilibrium strategy profile

$$s^*(\cdot) = (s_1^*(\cdot), \dots, s_n^*(\cdot))$$

of the game induced by M such that

$$g(s_1^*(\theta_1), \dots, s_n^*(\theta_n)) = f(\theta_1, \dots, \theta_n)$$

for all

$$(\theta_1, \dots, \theta_n) \in \Theta_1 \times \dots \times \Theta_n$$

Implementation

- We did not specify the type of equilibrium in the definition
 - ◆ (Mixed) Nash
 - ◆ Bayes-Nash
 - ◆ Dominant

Direct Mechanisms

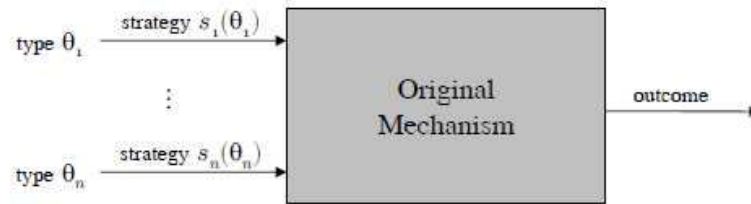
- Recall that a mechanism specifies the strategy sets of the agents
 - ♦ These sets can contain complex strategies
- **Direct mechanisms:**
 - ♦ Mechanism in which $S_i = \Theta_i$ for all i , and $g(\theta) = f(\theta)$ for all $\theta \in \Theta_1 \times \dots \times \Theta_n$
- **Incentive-compatible:**
 - ♦ A direct mechanism is incentive-compatible if it has an equilibrium s^* where $s_i^*(\theta_i) = \theta_i$ for all $\theta_i \in \Theta_i$ and all i
 - ♦ (truth telling by all agents is an equilibrium)
 - ♦ **Strategy-proof** if dominant-strategy equilibrium

11

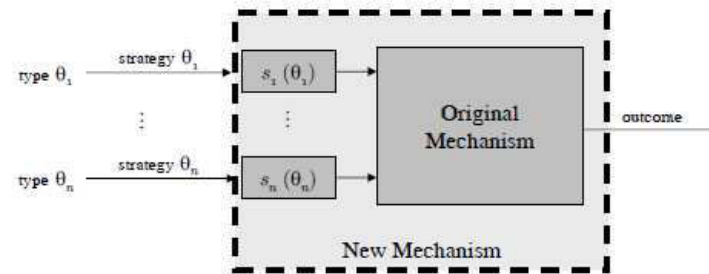
Dominant Strategy Implementation

- Is a certain social choice function implementable in dominant strategies? Did the mechanism enforce dominant strategies?
 - ♦ In principle we would need to consider all possible mechanisms
- **Revelation Principle** (for Dom Strategies)
 - ♦ Suppose there exists a **(in)direct** mechanism $M=(S_1, \dots, S_n, g(\cdot))$ that implements social choice function $f()$ in dominant strategies. Then there is a direct strategy-proof mechanism, M' , which also implements $f()$.

Revelation Principle: Intuition



(a) Revelation principle: original mechanism



(b) Revelation principle: new mechanism

Theoretical Implications

- Literal interpretation: Need only study direct mechanisms
 - This is a much smaller space of mechanisms
 - ♦ Negative results: If no direct mechanism can implement SCF $f()$ then no mechanism can do it => impossibility theorems, e.g. Arrow in voting.
 - ♦ Analysis tool:
 - Best direct mechanism gives us an upper bound on what we can achieve with an indirect mechanism
 - Analyze all direct mechanisms and choose the best one

Practical Implications

- Incentive-compatibility is “free” from an implementation perspective
- **BUT!!!**
 - ◆ A lot of mechanisms used in practice are not direct and incentive-compatible
 - ◆ Maybe there are some issues that are being ignored here

Quasi-Linear Preferences

- Outcome $o=(x,t_1,\dots,t_n)$
 - ♦ x is a “project choice” and $t_i \in \mathbf{R}$ are transfers (money)
- Utility function of agent i
 - ♦ $u_i(o,\theta_i)=u_i((x,t_1,\dots,t_n),\theta_i)=v_i(x,\theta_i)-t_i$
- Quasi-linear mechanism:
 $M=(S_1,\dots,S_n,g(\cdot))$ where $g(\cdot)=(x(\cdot),t_1(\cdot),\dots,t_n(\cdot))$

Social choice functions and quasi-linear settings

- SCF is **efficient** if for all types $\theta=(\theta_1,\dots,\theta_n)$
 - $\sum_{i=1}^n v_i(x(\theta),\theta_i) \geq \sum_{i=1}^n v_i(x'(\theta),\theta_i) \quad \forall x'(\theta)$
 - Aka social welfare maximizing, x is the selection function
- SCF is **budget-balanced** (BB) if
 - $\sum_{i=1}^n t_i(\theta)=0$
 - ♦ **Weakly budget-balanced** if $\sum_{i=1}^n t_i(\theta) \geq 0$

Groves Mechanisms

[Groves 1973]

- A **Groves mechanism**,
 $M=(S_1, \dots, S_n, (x, t_1, \dots, t_n))$ is defined by
 - ◆ Choice rule $x^*(\theta') = \operatorname{argmax}_x \sum_i v_i(x, \theta'_i)$
 - ◆ Transfer rules
 - $t_i(\theta') = h_i(\theta_{-i}') - \sum_{j \neq i} v_j(x^*(\theta'), \theta'_j)$

where $h_i(\cdot)$ is an (arbitrary) function that **does not depend** on the reported type θ'_i of agent i (quasi linear)

VCG Mechanism

(aka Clarke tax mechanism aka Pivotal mechanism)

- **Def:** Implement efficient outcome,

$$x^* = \operatorname{argmax}_x \sum_i v_i(x, \theta_i')$$

Compute transfers

$$t_i(\theta') = \sum_{j \neq i} v_j(x^{-i}, \theta_j') - \sum_{j \neq i} v_j(x^*, \theta_j')$$

Where $x^{-i} = \operatorname{argmax}_x \sum_{j \neq i} v_j(x, \theta_j')$

VCGs are efficient and strategy-proof

Agent's equilibrium utility is:

$$u_i(x^*, t_i, \theta_i) = v_i(x^*, \theta_i) - [\sum_{j \neq i} v_j(x^{-i}, \theta_j) - \sum_{j \neq i} v_j(x^*, \theta_j)]$$

$$= \sum_j v_j(x^*, \theta_j) - \sum_{j \neq i} v_j(x^{-i}, \theta_j)$$

= marginal contribution to the welfare of the system

Example: Building a pool

- The cost of building the pool is \$300
- If together all agents value the pool more than \$300 then it will be built
- VCG Mechanism:
 - ♦ Each agent announces their value, v_i
 - ♦ If $\sum v_i \geq 300$ then it is built
 - ♦ Payments $t_i(\theta_i') = \sum_{j \neq i} v_j(x^{-i}, \theta_j') - \sum_{j \neq i} v_j(x^*, \theta_j')$ if built, 0 otherwise

$$v_1=50, v_2=50, v_3=250$$

Pool should be built

$$t_1=(250+50)-(250+50)=0$$

$$t_2=(250+50)-(250+50)=0$$

$$t_3=(0)-(50+50)=-100$$

Not budget balanced

Example

- The government is deciding on number of street lights to be installed.
- Three beneficiaries - A, B, C.
- Four alternatives: $n = 0, 1, 2, 3$ where n is the number of street lights. The cost of a street light is 120.
- The government's objective to install the socially efficient number of street lights.

Net benefits with equal cost share

- If $n = 2$, the total cost is 240.
Hence, cost share for each is 80 (40 for each lamp).

Resident	No. of street lights			
	0	1	2	3
A	0	60	90	155
B	0	80	120	140
C	0	120	200	220
Cost	0	120	240	360

Net benefits with equal cost share

- The private net benefit for A is then $90 - 80 = 10$.
- Similarly for B and C and $n = 1, 3$. Figure show the benefits for each agent.

Resident	No. of street lights			
	0	1	2	3
A	0	60	90	155
B	0	80	120	140
C	0	120	200	220
Cost	0	120	240	360

Resident	No. of street lights			
	0	1	2	3
A	0	20	10	35
B	0	40	40	20
C	0	80	120	100
Social net benefit	0	140	170	155

28

Groves Clarke Taxes

- Is Person A pivotal? Does he has to pay a tax?

Resident	No. of street lights			
	0	1	2	3
A	0	20	10	35
B	0	40	40	20
C	0	80	120	100
Social net benefit	0	140	170	155

Resident	No. of street lights			
	0	1	2	3
B	0	40	40	20
C	0	80	120	100
Social net benefit	0	120	160	120

Person A is not pivotal. Without him, the net benefit is maximum at $n = 2$. With him the net benefit is maximum at $n = 2$. So his tax is zero.

Person B

Resident	No. of street lights					Resident	No. of street lights			
	0	1	2	3			0	1	2	3
A	0	20	10	35		A	0	20	10	35
B	0	40	40	20		C	0	80	120	100
C	0	80	120	100		Social net benefit	0	100	130	135
Social net benefit	0	140	170	155						

- ◆ Person B however is pivotal. With him the net benefit is maximum at $n = 2$. Without him the net benefit is maximum at $n = 3$.
- ◆ B's tax is the difference between the sum of net benefits of others at $n = 3$ and the sum of net benefits of others at $n = 2$, i.e. $135 - 130 = 5$.
- ◆ B is paying the tax because his report changes the decision from $n = 3$ to $n = 2$.

30

Person C

Resident	No. of street lights					Resident	No. of street lights			
	0	1	2	3			0	1	2	3
A	0	20	10	35		A	0	20	10	35
B	0	40	40	20		B	0	40	40	20
C	0	80	120	100		Social net benefit	0	60	50	55
Social net benefit	0	140	170	155						

- ◆ Person C is pivotal as well. With him the net benefit is maximum at $n = 2$. Without him the net benefit is maximum at $n = 1$
- ◆ C's tax is therefore the sum of others' benefits at $n = 1$ and the sum of others' benefits at $n = 2$, i.e. $60 - 50 = 10$.

Net benefits with taxes

Resident	No. of street lights				Tax
	0	1	2	3	
A	0	20	10	35	0
B	0	40	40	20	5
C	0	80	120	100	10
Social net benefit	0	140	170	155	

- Post tax net benefit from this scheme:
10 for A,
 $40 - 5 = 35$ for B,
 $120 - 10 = 110$ for C.

Incentives for truthful revelation

Resident	No. of street lights			
	0	1	2	3
A	0	20	10	35
B	0	40	40	20
C	0	80	120	100
Social net benefit	0	140	170	190

? 70

←

- ◆ Notice that A's net benefit is maximum at $n = 3$. Does he have an incentive to lie and change the decision to $n = 3$?
- ◆ Suppose A states his net benefit from $n = 3$ to be 70 instead of 35. Then, sum of stated net benefits is maximum at $n = 3$.

Incentives for truthful revelation

Resident	No. of street lights					Resident	No. of street lights			
	0	1	2	3			0	1	2	3
A	0	20	10	70		B	0	40	40	20
B	0	40	40	20		C	0	80	120	100
C	0	80	120	100		Social net benefit	0	120	160	120
Social net benefit	0	140	170	190						

- ♦ But then A becomes pivotal. Without him the sum of net benefits is maximum at $n = 2$.
His report changes the decision from $n = 2$ to $n = 3$.
- ♦ So he has to pay a tax and his tax will be equal to $160 - 120 = 40$.

Incentives for truthful revelation

- A's net benefit from lying will be
(Net benefit from $n = 3$) - Tax
 $= 35 - 40$
 $= -5$
- A's net benefit from truthfully reporting is 10.
- Hence A doesn't have incentive to lie.
- You can repeat the same exercise for B and C to verify that they do not have incentive to lie either.

Clarke tax mechanism...

- Pros
 - ◆ Social welfare maximizing outcome
 - ◆ Truth-telling is a dominant strategy
 - ◆ Feasible in that it does not need a benefactor ($\sum_i t_i \leq 0$) (not discussed here)

Participation Constraints

- Agents can not be forced to participate in a mechanism
 - ◆ It must be in their own best interest
- A mechanism is **individually rational** (IR) if an agent's (expected) utility from participating is (weakly) better than what it could get by not participating

Participation Constraints

- Can classify mechanisms based on participation constraints
 - ♦ Let $u_i^*(\theta_i)$ be an agent's utility if it does not participate and has type θ_i
 - ♦ **Ex ante IR:** An agent must decide to participate before it knows its own type and other types
 - $E_{\theta \in \Theta}[u_i(f(\theta), \theta_i)] \geq E_{\theta_i \in \Theta_i}[u_i^*(\theta_i)]$
 - ♦ **Interim IR:** An agent decides whether to participate once it knows its own type, but no other agent's type
 - $E_{\theta_{-i} \in \Theta_{-i}}[u_i(f(\theta_i, \theta_{-i}), \theta_i)] \geq u_i^*(\theta_i)$
 - ♦ **Ex post IR:** An agent decides whether to participate after it knows everyone's types (after the mechanism has completed)
 - $u_i(f(\theta), \theta_i) \geq u_i^*(\theta_i)$

Quick Review

- Gibbard-Satterthwaite
 - ♦ Impossible to get non-dictatorial mechanisms if using **dominant strategy implementation** and **general preferences**
- Groves
 - ♦ Possible to get dominant strategy implementation with quasi-linear utilities
 - Efficient
- Clarke (or VCG)
 - ♦ Possible to get dominant strategy implementation with quasi-linear utilities
 - Efficient, interim IR

The End

- Exam: 30.03, 9:00, Audimax I
- Remember comments in exercises
- There will be no questions about “Mechanism Design” in the exam!!!.