

Module 2: Linear Programs (Simplex – A First Attempt)

A Naive Strategy for Solving an LP

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Many details missing!

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- How do we find a feasible solution?

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- Will this ever terminate?

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The **SIMPLEX** algorithm works along these lines.

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- How do we find a “better” solution?
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The **SIMPLEX** algorithm works along these lines.

In this lecture: A first attempt at this algorithm.

A First Example

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Consider

$$\max (4, 3, 0, 0)x + 7$$

s.t.

$$\begin{pmatrix} 3 & 2 & 1 & 0 \\ 1 & 1 & 0 & 1 \end{pmatrix} x = \begin{pmatrix} 2 \\ 1 \end{pmatrix}$$

$$x_1, x_2, x_3, x_4 \geq 0$$

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- We have a feasible solution: $x_1 = 0, x_2 = 0, x_3 = 2$, and $x_4 = 1$.

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The feasible solution has objective value: $4 \times 0 + 3 \times 0 + 7 = 7$.

- Can we find a feasible solution with value larger than 7?

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Increase x_1 as much as possible, and keep x_2 unchanged, i.e.,

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Idea

Increase x_1 as much as possible, and keep x_2 unchanged, i.e.,

$$\begin{array}{ll} x_1 = t & \text{for some } t \geq 0 \text{ as large as possible} \\ x_2 = 0 & \end{array}$$

$$\max (4, 3, 0, 0)x + 7$$

s.t.

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$$x_1, x_2, x_3, x_4 \geq 0$$

$$x_1 = t$$

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Choose $t \geq 0$ as large as possible.

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It needs to satisfy

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Choose $t \geq 0$ as large as possible.

It needs to satisfy

1. the equality constraints, and
2. the non-negativity constraints.

Satisfying the Equality Constraints

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$$\begin{pmatrix} 2 \\ 1 \end{pmatrix} = \begin{pmatrix} 3 & 2 & 1 & 0 \\ 1 & 1 & 0 & 1 \end{pmatrix} x$$

Satisfying the Equality Constraints

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$$\begin{pmatrix} 2 \\ 1 \end{pmatrix} = \begin{pmatrix} 3 & 2 & 1 & 0 \\ 1 & 1 & 0 & 1 \end{pmatrix} x$$

$$= x_1 \begin{pmatrix} 3 \\ 1 \end{pmatrix} + x_2 \begin{pmatrix} 2 \\ 1 \end{pmatrix} + x_3 \begin{pmatrix} 1 \\ 0 \end{pmatrix} + x_4 \begin{pmatrix} 0 \\ 1 \end{pmatrix}$$

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$$= t \begin{pmatrix} 3 \\ 1 \end{pmatrix} + 0 \begin{pmatrix} 2 \\ 1 \end{pmatrix} + \begin{pmatrix} x_3 \\ 0 \end{pmatrix} + \begin{pmatrix} 0 \\ x_4 \end{pmatrix}$$

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$$\begin{pmatrix} x_3 \\ x_4 \end{pmatrix} = \begin{pmatrix} 2 \\ 1 \end{pmatrix} - t \begin{pmatrix} 3 \\ 1 \end{pmatrix}$$

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Remark

Equality constraints hold for any choice of t .

Satisfying the Non-Negativity Constraints

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Choose $t \geq 0$ as large as possible.

Satisfying the Non-Negativity Constraints

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$$x_1 = t \geq 0 \quad \checkmark$$

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Choose $t \geq 0$ as large as possible.

$$x_1 = t \geq 0 \quad \checkmark$$

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$$x_3 = 2 - 3t \geq 0 \quad \rightarrow \quad t \leq \frac{2}{3}$$

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Choose $t \geq 0$ as large as possible.

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$$x_3 = 2 - 3t \geq 0 \quad \longrightarrow \quad t \leq \frac{2}{3}$$

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Thus, the largest possible t is $\min \left\{ 1, \frac{2}{3} \right\} = \frac{2}{3}$.

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Thus, the largest possible t is $\min \{1, \frac{2}{3}\} = \frac{2}{3}$. The new solution is

$$x = (t, 0, 2 - 3t, 1 - t)^\top = \left(\frac{2}{3}, 0, 0, \frac{1}{3}\right)^\top$$

Repeating the Argument?

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s.t.

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Is the new solution optimal?

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Question

Is the new solution optimal? **NO!**

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Is the new solution optimal? **NO!**

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Can we use the same trick to get a better solution?

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What made it work the first time around?

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The LP needs to be in “canonical” form.

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$$x_4 = 1$$

Revised strategy:

Remark

The LP needs to be in “canonical” form.

$$\max \quad (4 \quad 3 \quad 0 \quad 0)x + 7$$

s.t.

$$\begin{pmatrix} 3 & 2 & 1 & 0 \\ 1 & 1 & 0 & 1 \end{pmatrix} x = \begin{pmatrix} 2 \\ 1 \end{pmatrix}$$

$$x_1, x_2, x_3, x_4 \geq 0$$

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Step 1. Find a feasible solution, x .

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Step 1. Find a feasible solution, x .

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Step 4. If LP is unbounded, STOP.

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Revised strategy:

- Step 1.** Find a feasible solution, x .
- Step 2.** Rewrite LP so that it is in “canonical” form.
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- Step 4.** If LP is unbounded, STOP.
- Step 5.** Find a “better” feasible solution.

Remark

The LP needs to be in “canonical” form.

$$\begin{array}{l} \max \quad (4 \quad 3 \quad 0 \quad 0)x + 7 \\ \text{s.t.} \\ \quad \begin{pmatrix} 3 & 2 & 1 & 0 \\ 1 & 1 & 0 & 1 \end{pmatrix} x = \begin{pmatrix} 2 \\ 1 \end{pmatrix} \\ \quad x_1, x_2, x_3, x_4 \geq 0 \end{array}$$

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- Step 1.** Find a feasible solution, x .
- Step 2.** Rewrite LP so that it is in “canonical” form.
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