

Module 1: Formulations (LP Models)

Constrained Optimization

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This class: all functions are **affine**.

Modeling: Linear Programs

Affine Functions

Definition

A function $f : \mathbb{R}^n \rightarrow \mathbb{R}$ is **affine** if $f(x) = a^T x + \beta$ for $a \in \mathbb{R}^n$, $\beta \in \mathbb{R}$.

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- (iii) $f(x) = 5x - 3 \cos(x) + \sqrt{x}$ (**not affine and not linear**)

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The optimization problem

$$\begin{aligned} \min \{ & f(x) : g_i(x) \leq b_i, \\ & \forall 1 \leq i \leq m, x \in \mathbb{R}^n \} \end{aligned} \quad (\text{P})$$

is called a **linear program** if f is **affine** and g_1, \dots, g_m is **finite** number of **linear** functions.

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Instead of **set notation**, we often write LPs more verbosely

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Sometimes replace **subject**

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We often write $x \geq 0$ as a short for **all variables are non-negative**.

$$\text{min} \quad -x_1 - 2x_2 - x_3$$

$$\text{s.t.} \quad 2x_1 + x_3 \geq 3$$

$$x_1 + 2x_2 = 2$$

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Second mathematical program is **not an LP**.

Three reasons:

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$$\text{subject to } 2x_1 + x_3 < 3$$

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Production revisited

$$\begin{aligned} \max \quad & 300x_1 + 260x_2 + 220x_3 + 180x_4 - 8y_s - 6y_u \\ & 11x_1 + 7x_2 + 6x_3 + 5x_4 \leq 700 \\ & 4x_1 + 6x_2 + 5x_3 + 4x_4 \leq 500 \\ & 8x_1 + 5x_2 + 5x_3 + 6x_4 \leq y_s \\ \text{s.t.} \quad & 7x_1 + 8x_2 + 7x_3 + 4x_4 \leq y_u \\ & y_s \leq 600 \\ & y_u \leq 650 \\ & x_1, x_2, x_3, x_4, y_u, y_s \geq 0. \end{aligned}$$

The mathematical program for **WaterTech** example from last class is in fact an LP!

Example: Multiperiod Models

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One such example: **Multiperiod Models**:

Time is split into **periods**,

We have to make a **decision in each period**, and

All decisions influence the final outcome.

KW Oil

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Needs to decide on **how much oil to purchase** in order to **satisfy demand** of its customers

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Demand (ℓ)	5000	8000	9000	6000

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Question: **When** should we purchase **how much** oil?

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Additional complication: Company has storage tank that
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Additional complication: Company has storage tank that
Currently (beginning of month 1) contains 2000 litres of oil, and
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Assumption: Oil is delivered at beginning of month, and
consumption occurs mid month

KW Oil Model – Variables

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Objective function

Minimize **cost of oil** procurement.

Variables:

p_i : oil purchase in month i

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Constraints: when do

$$t_1, \dots, t_4, p_1, \dots, p_4$$

correspond to a **feasible purchasing scheme?**

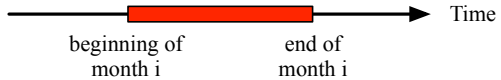
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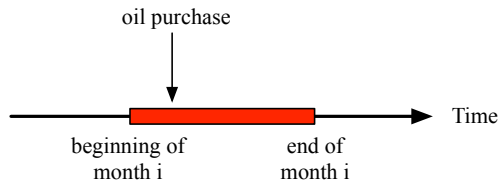
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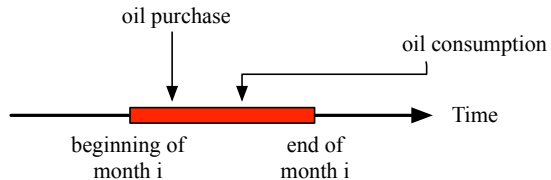
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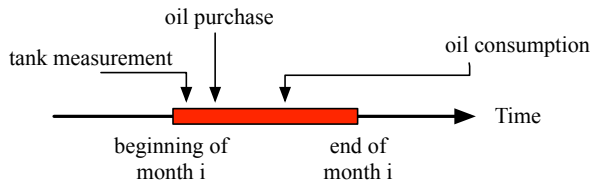
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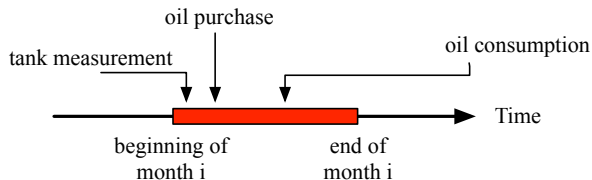
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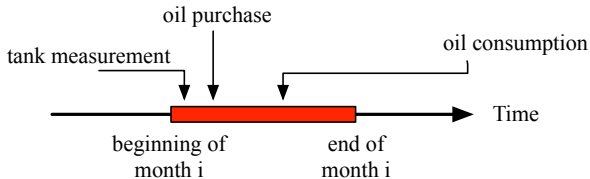
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We need: $p_i + t_i \geq \{\text{demand in month } i\}$

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Variables:

p_i : oil purchase in month i

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[Balance Equation]
$$p_i + t_i = \{\text{demand in month } i\} + t_{i+1}$$

Constraints

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Tank content in month 1: 2000 litres

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Month 1:

$$p_1 + 2000 = 5000 + t_2$$

Month	1	2	3	4
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Month 4:

$$p_4 + t_4 \geq 6000$$

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Tank content in month 1: 2000 litres

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subject to

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Solution: $p = (3000, 12000, 5000, 6000)^T$, and
 $t = (2000, 0, 4000, 0)^T$

KW Oil: Add-Ons

Can easily capture **additional features**. E.g. ...

Storage comes at a cost:
storage cost is \$.15 per
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- (ii) Will have to add constraints

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$$p_1 + t_1 = 5000 + t_2$$

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$$p_3 + t_3 = 9000 + t_4$$

$$p_4 + t_4 \geq 6000$$

$$t_1 = 2000$$

$$t_i \leq 4000 \quad (i = 2, 3, 4)$$

$$t_i, p_i \geq 0 \quad (i = 1, 2, 3, 4)$$

KW Oil: Add-Ons

- (i) Add variable M for maximum #l purchased over all months.

$$\min \quad 0.75p_1 + 0.72p_2 + 0.92p_3 + 0.90p_4$$

subject to

$$p_1 + t_1 = 5000 + t_2$$

$$p_2 + t_2 = 8000 + t_3$$

$$p_3 + t_3 = 9000 + t_4$$

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Goal: Minimize maximum #l of oil purchased over all months.

KW Oil: Add-Ons

- (i) Add variable M for maximum #l purchased over all months.
- (ii) Add constraints

$$p_i \leq M$$

for all $i \in [4]$.

$$\min \quad 0.75p_1 + 0.72p_2 + 0.92p_3 + 0.90p_4$$

subject to

$$p_1 + t_1 = 5000 + t_2$$

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Goal: Minimize maximum #l of oil purchased over all months.

KW Oil: Add-Ons

- (i) Add variable M for maximum #l purchased over all months.
- (ii) Add constraints

$$p_i \leq M$$

for all $i \in [4]$.

- (iii) Done?

$$\min \quad 0.75p_1 + 0.72p_2 + 0.92p_3 + 0.90p_4$$

subject to

$$p_1 + t_1 = 5000 + t_2$$

$$p_2 + t_2 = 8000 + t_3$$

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Goal: Minimize maximum #l of oil purchased over all months.

KW Oil: Add-Ons

- (i) Add variable M for maximum #l purchased over all months.
- (ii) Add constraints

$$p_i \leq M$$

for all $i \in [4]$.

- (iii) **Done?** No! Need to replace objective function by

$$\min M$$

$$\min 0.75p_1 + 0.72p_2 + 0.92p_3 + 0.90p_4$$

subject to

$$p_1 + t_1 = 5000 + t_2$$

$$p_2 + t_2 = 8000 + t_3$$

$$p_3 + t_3 = 9000 + t_4$$

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Goal: Minimize maximum #l of oil purchased over all months.

Minimizing the Maximum Purchase: LP

min M

s.t.

$$p_1 + t_1 = 5000 + t_2$$

$$p_2 + t_2 = 8000 + t_3$$

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$$t_i \leq 4000 \quad (i = 2, 3, 4)$$

$$p_i \leq M \quad (i = 1, 2, 3, 4)$$

$$t_i, p_i \geq 0 \quad (i = 1, 2, 3, 4)$$

KW Oil: Correctness

Why is this a **correct** model?

min M

s.t.

$$p_1 + t_1 = 5000 + t_2$$

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KW Oil: Correctness

Why is this a **correct** model?

Suppose that

$M, p_1, \dots, p_4, t_1, \dots, t_4$

is an **optimal**

solution to the LP

min M

s.t.

$$p_1 + t_1 = 5000 + t_2$$

$$p_2 + t_2 = 8000 + t_3$$

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KW Oil: Correctness

Why is this a **correct** model?

Suppose that

$M, p_1, \dots, p_4, t_1, \dots, t_4$

is an **optimal**

solution to the LP

Clearly:

$M \geq \max_i p_i$

min M

s.t.

$$p_1 + t_1 = 5000 + t_2$$

$$p_2 + t_2 = 8000 + t_3$$

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KW Oil: Correctness

Why is this a **correct** model?

Suppose that

$M, p_1, \dots, p_4, t_1, \dots, t_4$

is an **optimal**

solution to the LP

Clearly:

$$M \geq \max_i p_i$$

Since M, p, t is

optimal we must

have $M = \max_i p_i$.

Why?

$$\min M$$

s.t.

$$p_1 + t_1 = 5000 + t_2$$

$$p_2 + t_2 = 8000 + t_3$$

$$p_3 + t_3 = 9000 + t_4$$

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