Module 1: Formulations (LP Models)

 $\min\{f(x) : g_i(x) \le b_i, (1 \le i \le m), x \in \mathbb{R}^n\},\$ 

where

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 $n, m \in \mathbb{N}$ ,  $b_1, \ldots, b_m \in \mathbb{R}$ , and f,  $g_1, \ldots, g_m$  are functions with from  $\mathbb{R}^n$  to  $\mathbb{R}$ .

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## Modeling: Linear Programs

## **Definition**

 $\mathsf A$  function  $f:\mathbb R^n\to\mathbb R$  is affine if  $f(x)=a^Tx+\beta$  for  $a\in\mathbb R^n,$   $\beta\in\mathbb R.$ 

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\n

The optimization problem

$$
\min\{f(x) : g_i(x) \le b_i, \forall 1 \le i \le m, x \in \mathbb{R}^n\}
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 (P)

is called a linear program if f is affine and  $g_1, \ldots, g_m$  is finite number of linear functions.

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#### Comments:

Instead of set notation, we often write LPs more verbosely

 $max \quad -2x_1 + 7x_3$ subject to  $x_1 + 7x_2 \leq 3$  $3x_2 + 4x_3 < 2$  $x_1, x_3 \geq 0$ 

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Sometimes replace subject

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 $\min$  –  $x_1 - 2x_2 - x_3$ s.t.  $2x_1 + x_3 \ge 3$  $x_1 + 2x_2 = 2$  $x \geq 0$ 

Second mathematical program is not an LP. Three reasons:

min 
$$
-x_1 - 2x_2 - x_3
$$
  
s.t.  $2x_1 + x_3 \ge 3$   
 $x_1 + 2x_2 = 2$   
 $x \ge 0$ 

max  $-1/x_1 - x_3$ subject to  $2x_1 + x_3 < 3$  $x_1 + \alpha x_2 = 2 \quad \forall \alpha \in \mathbb{R}$ 

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> Dividing by variables is not allowed

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## Production revisited

max  $300x_1 + 260x_2 + 220x_3 + 180x_4 - 8y_5 - 6y_4$ s.t.  $7x_1 + 8x_2 + 7x_3 + 4x_4 \leq y_u$  $11x_1 + 7x_2 + 6x_3 + 5x_4 \leq 700$  $4x_1 + 6x_2 + 5x_3 + 4x_4 \leq 500$  $8x_1 + 5x_2 + 5x_3 + 6x_4 \leq y_s$  $y_s \leq 600$  $y_u \leq 650$  $x_1, x_2, x_3, x_4, y_u, y_s > 0.$ 

The mathematical program for WaterTech example from last class is in fact an LP!

## Example: Multiperiod Models

Main feature of WaterTech production model: Decisions about production levels have to be made once and for all. Main feature of WaterTech production model: Decisions about production levels have to be made once and for all. In practice, we often have to make a sequence of decision that influence each other.

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One such example: Multiperiod Models:

Time is split into periods, We have to make a decision in each period, and All decisions influence the final outcome.

KW Oil is local supplier of heating oil Needs to decide on how much oil to purchase in order to satisfy demand of its customers

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Years of experience give the following demand forecast for the next 4 months:



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The projected price of oil fluctuates from month to month:



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Years of experience give the following demand forecast for the next 4 months:



The projected price of oil fluctuates from month to month:



Question: When should we purchase how much oil?

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Additional complication: Company has storage tank that

Currently (beginning of month 1) contains 2000 litres of oil, and Has a capacity of 4000 litres of oil.
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Additional complication: Company has storage tank that

Currently (beginning of month 1) contains 2000 litres of oil, and Has a capacity of 4000 litres of oil.

Assumption: Oil is delivered at beginning of month, and consumption occurs mid month









(i) Need to decide how many litres of oil to purchase in each month  $i$ .





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- (ii) How much oil is stored in the tank at beginning of month  $i$ ?





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- (ii) How much oil is stored in the tank at beginning of month  $i$ ?  $\longrightarrow$  variable  $t_i$  for  $i \in [4]$

# Objective function

Minimize cost of oil procurement. Variables:

 $p_i$  : oil purchase in month  $i$  $t_i$  : tank level in month  $i$ 

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Minimize cost of oil procurement.

min  $0.75p_1 + 0.72p_2 + 0.92p_3 + 0.90p_4$ 

Variables:

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# Objective function

Minimize cost of oil procurement.

min  $0.75p_1 + 0.72p_2 + 0.92p_3 + 0.90p_4$ 

Constraints: when do

 $t_1, \ldots, t_4, p_1, \ldots, p_4$ 

correspond to a feasible purchasing scheme?

Variables:

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#### Variables:

- $p_i$  : oil purchase in month  $i$
- $t_i$  : tank level in month  $i$





[Balance Equation]  $p_i + t_i = \{$ demand in month  $i\} + t_{i+1}$ 

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Tank content in month 1: 2000 litres

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Tank content in month 1: 2000 litres

[Balance Equation]  $p_i + t_i = \{$ demand in month  $i\} + t_{i+1}$ 



Tank content in month 1: 2000 litres

 $p_1 + 2000 = 5000 + t_2$ Month 2:  $p_2 + t_2 = 8000 + t_3$ 

[Balance Equation]  $p_i + t_i = \{$ demand in month  $i\} + t_{i+1}$ 





Tank content in month 1: 2000 litres

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Tank content in month 1: 2000 litres

## KW Oil: Entire LP

min  $0.75p_1 + 0.72p_2 + 0.92p_3 + 0.90p_4$ subject to

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p_1+t_1 = 5000+t_2
$$
  
\n
$$
p_2+t_2 = 8000+t_3
$$
  
\n
$$
p_3+t_3 = 9000+t_4
$$
  
\n
$$
p_4+t_4 \ge 6000
$$
  
\n
$$
t_1 = 2000
$$
  
\n
$$
t_i \le 4000
$$
 (*i*=2,3,4)  
\n
$$
t_i, p_i \ge 0
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p_4 + t_4 \ge 6000
$$
  
\n
$$
t_1 = 2000
$$
  
\n
$$
t_i \le 4000
$$
 (*i* = 2, 3, 4)  
\n
$$
t_i, p_i \ge 0
$$
 (*i* = 1, 2, 3, 4)

Solution:  $p = (3000, 12000, 5000, 6000)^T$ , and  $t = (2000, 0, 4000, 0)^T$ 

Can easily capture additional features. E.g. ...

> Storage comes at a cost: storage cost is \$.15 per litre/month.

min  $0.75p_1 + 0.72p_2 + 0.92p_3 + 0.90p_4$ subject to

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(i = 1, 2, 3, 4)
$$

 $H = 0.00$ 

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Add  $\sum_{i=1}^4.15t_i$  to objective.

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Minimize maximum  $#I$  of oil purchased over all months

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p_1 + t_1 = 5000 + t_2
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$$
  
\n
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t_i = 4000
$$
  
\n
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t_i, p_i \ge 0
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(i = 2, 3, 4)
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$$

 $5000$ 

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(i) We will need a new variable  $M$  for maximum #l purchased

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p_4+t_4 \ge 6000
$$
  
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$$
t_1 \le 4000 \qquad (i=2,3,4)
$$
  
\n
$$
t_i, p_i \ge 0 \qquad (i=1,2,3,4)
$$

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Add  $\sum_{i=1}^4.15t_i$  to objective.

Minimize maximum  $#I$  of oil purchased over all months

- (i) We will need a new variable  $M$  for maximum #l purchased
- (ii) Will have to add constraints

min  $0.75p_1 + 0.72p_2 + 0.92p_3 + 0.90p_4$ subject to

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p_1 + t_1 = 5000 + t_2
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min

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$$

Goal: Minimize maximum #l of oil purchased over all months.

- (i) Add variable  $M$  for maximum  $#$ l purchased over all months.
- (ii) Add constraints

$$
p_i \leq M
$$
 for all  $i \in [4].$ 

min  $0.75p_1 + 0.72p_2 + 0.92p_3 + 0.90p_4$ subject to

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p_1+t_1 =
$$

$$
p_2 + t_2 = 8000 + t_3
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t_1 = 2000
$$
  
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$$
t_i \le 4000 \qquad (i = 2, 3, 4)
$$
  
\n
$$
t_i, p_i \ge 0 \qquad (i = 1, 2, 3, 4)
$$

 $5000 + t_2$ 

Goal: Minimize maximum #l of oil purchased over all months.

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#### Goal: Minimize maximum #l of oil purchased over all months.

(iii) Done?

- (i) Add variable  $M$  for maximum #l purchased over all months.
- (ii) Add constraints

$$
p_i \leq M
$$

for all  $i \in [4]$ .

(iii) Done? No! Need to replace objective function by

min  $0.75p_1 + 0.72p_2 + 0.92p_3 + 0.90p_4$ 

subject to

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p_1+t_1 = 5000+t_2
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t_i \le 4000 \qquad (i=2,3,4)
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t_i,p_i \ge 0 \qquad (i=1,2,3,4)
$$

 $\min$  M Goal: Minimize maximum #l of oil purchased over all months.

## Minimizing the Maximum Purchase: LP

min M

s.t.

$$
p_1 + t_1 = 5000 + t_2
$$
  
\n
$$
p_2 + t_2 = 8000 + t_3
$$
  
\n
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$$
p_i \le M \qquad (i = 1, 2, 3, 4)
$$
  
\n
$$
t_i, p_i \ge 0 \qquad (i = 1, 2, 3, 4)
$$

Why is this a correct model?

 $min$   $M$ s.t.



Why is this a correct  $min$   $M$ model? s.t. Suppose that  $M, p_1, \ldots, p_4, t_1, \ldots, t_4$ is an optimal solution to the LP





