Module 1: Formulations (LP Models)

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This class: all functions are affine.

Modeling: Linear Programs

Definition

A function $f : \mathbb{R}^n \to \mathbb{R}$ is affine if $f(x) = a^T x + \beta$ for $a \in \mathbb{R}^n$, $\beta \in \mathbb{R}$.

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(iii) $f(x) = 5x - 3\cos(x) + \sqrt{x}$ (not affine and not linear)

The optimization problem

$$\min\{f(x) : g_i(x) \le b_i, \\ \forall 1 \le i \le m, \, x \in \mathbb{R}^n\}$$
(P)

is called a linear program if f is affine and g_1, \ldots, g_m is finite number of linear functions.

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Comments:

Instead of set notation, we often write LPs more verbosely

 $\begin{array}{ll} \max & -2x_1+7x_3\\ \text{subject to} & x_1+7x_2\leq 3\\ & 3x_2+4x_3\leq 2\\ & x_1,x_3\geq 0 \end{array}$

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Sometimes replace subject

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 $\begin{array}{ll} \min & -x_1 - 2x_2 - x_3 \\ \text{s.t.} & 2x_1 + x_3 \geq 3 \\ & x_1 + 2x_2 = 2 \\ & x \geq 0 \end{array}$

Second mathematical program is not an LP. Three reasons:

min
$$-x_1 - 2x_2 - x_3$$

s.t. $2x_1 + x_3 \ge 3$
 $x_1 + 2x_2 = 2$
 $x \ge 0$

$$\begin{array}{ll} \max & -1/x_1 - x_3\\ \text{subject to} & 2x_1 + x_3 < 3\\ & x_1 + \alpha \, x_2 = 2 \quad \forall \alpha \in \mathbb{R} \end{array}$$

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Production revisited

The mathematical program for WaterTech example from last class is in fact an LP!

Example: Multiperiod Models

Main feature of WaterTech production model: Decisions about production levels have to be made once and for all. Main feature of WaterTech production model: Decisions about production levels have to be made once and for all. In practice, we often have to make a sequence of decision that influence each other. Main feature of WaterTech production model: Decisions about production levels have to be made once and for all.

In practice, we often have to make a sequence of decision that influence each other.

One such example: Multiperiod Models:

Time is split into periods, We have to make a decision in each period, and All decisions influence the final outcome.

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Month	1	2	3	4
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Assumption: Oil is delivered at beginning of month, and consumption occurs mid month

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- (ii) How much oil is stored in the tank at beginning of month i?

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- (i) Need to decide how many litres of oil to purchase in each month i. \longrightarrow variable p_i for $i \in [4]$
- (ii) How much oil is stored in the tank at beginning of month i? \longrightarrow variable t_i for $i \in [4]$

Objective function

Minimize cost of oil procurement.

Variables:

 p_i : oil purchase in month i t_i : tank level in month i

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min $0.75p_1 + 0.72p_2 + 0.92p_3 + 0.90p_4$

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Objective function

Minimize cost of oil procurement.

 $\min \quad 0.75p_1 + 0.72p_2 + 0.92p_3 + 0.90p_4$

Constraints: when do

 $t_1,\ldots,t_4,p_1,\ldots,p_4$

correspond to a feasible purchasing scheme?

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- p_i : oil purchase in month i
- t_i : tank level in month i





[Balance Equation] $p_i + t_i = \{ \text{demand in month } i \} + t_{i+1}$

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Month 1	Month	1	2	3	4
m + 2000 - 5000 + t	Demand (ℓ)	5000	8000	9000	6000
$p_1 + 2000 - 3000 + \iota_2$					

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 $p_1 + 2000 = 5000 +$ Month 2: $p_2 + t_2 = 8000 + t_3$

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Month 1: $p_1 + 2000 = 5000 + t_2$ Month 2: $p_2 + t_2 = 8000 + t_3$ Month 3: $p_3 + t_3 = 9000 + t_4$

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Demand (ℓ)	5000	8000	9000	6000

Tank content in month 1: 2000 litres

KW Oil: Entire LP

$$p_{1}+t_{1} = 5000+t_{2}$$

$$p_{2}+t_{2} = 8000+t_{3}$$

$$p_{3}+t_{3} = 9000+t_{4}$$

$$p_{4}+t_{4} \ge 6000$$

$$t_{1} = 2000$$

$$t_{i} \le 4000 \quad (i=2,3,4)$$

$$t_{i},p_{i} \ge 0 \quad (i=1,2,3,4)$$

KW Oil: Entire LP

min $0.75p_1 + 0.72p_2 + 0.92p_3 + 0.90p_4$ subject to

$$p_{1}+t_{1} = 5000+t_{2}$$

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$$t_{1} = 2000$$

$$t_{i} \le 4000 \quad (i=2,3,4)$$

$$t_{i},p_{i} \ge 0 \quad (i=1,2,3,4)$$

Solution: $p = (3000, 12000, 5000, 6000)^T$, and $t = (2000, 0, 4000, 0)^T$

Can easily capture additional features. E.g. ...

Storage comes at a cost: storage cost is \$.15 per litre/month.

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Minimize maximum #I of oil purchased over all months

- (i) We will need a new variable M for maximum #I purchased
- (ii) Will have to add constraints

min $0.75p_1 + 0.72p_2 + 0.92p_3 + 0.90p_4$ subject to

(i) Add variable *M* for maximum #I purchased over all months.

min $0.75p_1 + 0.72p_2 + 0.92p_3 + 0.90p_4$

subject to

- (i) Add variable M for maximum #I purchased over all months.
- (ii) Add constraints

$$p_i \leq M$$

for all $i \in [4]$.

min $0.75p_1 + 0.72p_2 + 0.92p_3 + 0.90p_4$ subject to

- (i) Add variable M for maximum #I purchased over all months.
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(iii) Done?

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for all $i \in [4]$.

(iii) Done? No! Need to replace objective function by

min M

min $0.75p_1 + 0.72p_2 + 0.92p_3 + 0.90p_4$ subject to

Minimizing the Maximum Purchase: LP

min M

s.t.

$$p_{1} + t_{1} = 5000 + t_{2}$$

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$$p_{4} + t_{4} \ge 6000$$

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Why is this a correct model?

min Ms.t.

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model?minMSuppose thats.t. $M, p_1, \dots, p_4, t_1, \dots, t_4$
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Suppose that	s.t.				
$M, p_1, \ldots, p_4, t_1, \ldots, t_4$		p_1+t_1	=	$5000 + t_2$	
is an optimal		$p_2 + t_2$	=	$8000 + t_3$	
solution to the LP		$p_3 + t_3$	=	$9000 + t_4$	
		$p_4 + t_4$	\geq	6000	
Clearly:		t_1	=	2000	
$M \ge \max_i p_i$		t_i	\leq	4000	(i = 2, 3, 4)
		p_i	\leq	M	(i = 1, 2, 3, 4)
		t_i, p_i	\geq	0	$\left(i=1,2,3,4\right)$

Why is this a correct model?	min	M			
Suppose that $M, p_1, \ldots, p_4, t_1, \ldots, t_4$ is an optimal	s.t.	$p_1+t_1 \ p_2+t_2$	=	$5000 + t_2$ $8000 + t_3$	
solution to the LP Clearly: $M \ge \max_i p_i$		$p_3+t_3\ p_4+t_4\ t_1\ t_i$	> <	$9000 + t_4$ 6000 2000 4000	(i = 2, 3, 4)
Since M, p, t is optimal we must have $M = \max_i p_i$. Why?		$rac{p_i}{t_i,p_i}$	<mark> </mark> <u> </u>	<i>M</i> 0	(i = 1, 2, 3, 4) (i = 1, 2, 3, 4)