Module 2: Linear Programs (Formalizing the Simplex)







Consider $B = \{1, 4\}.$

• A_B is square and non-singular



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• A_B is square and non-singular $\implies B$ is a basis



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$$A_B = I$$
 and $c_B = \mathbf{0}$



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- $A_B = I$ and $c_B = 0$ \implies LP is in canonical form for B



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- $\bar{x} = (2, 0, 0, 5)^{\top}$ is a the basic solution for B.



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- $\bar{x} = (2, 0, 0, 5)^{\top}$ is a the basic solution for B.
- $\bar{x} \ge \mathbf{0}$



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• $A_B = I$ and $c_B = 0$ \implies LP is in canonical form for B

- $\bar{x} = (2, 0, 0, 5)^{\top}$ is a the basic solution for B.
- $\bar{x} \ge \mathbf{0} \implies \bar{x}$ is feasible, i.e., B is feasible







Question

How do we find a better feasible solution?



Idea

Pick $k \notin B$ such that $c_k > 0$.



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Set $x_k = t \ge 0$ as large as possible.



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Keep all other non-basic variables at 0.

Pick k = 2. Set $x_2 = t \ge 0$.



Idea

Pick
$$k \notin B$$
 such that $c_k > 0$.

Set $x_k = t \ge 0$ as large as possible.

Keep all other non-basic variables at 0.

Pick
$$k = 2$$
. Set $x_2 = t \ge 0$.

Keep $x_3 = 0$.





Idea



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$$\begin{pmatrix} 2\\5 \end{pmatrix} = \begin{pmatrix} 1 & 1 & 2 & 0\\ 0 & 1 & 1 & 1 \end{pmatrix} x$$



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$$\begin{pmatrix} 2\\5 \end{pmatrix} = \begin{pmatrix} 1 & 1 & 2 & 0\\ 0 & 1 & 1 & 1 \end{pmatrix} x = x_1 \begin{pmatrix} 1\\0 \end{pmatrix} + x_2 \begin{pmatrix} 1\\1 \end{pmatrix} + x_3 \begin{pmatrix} 2\\1 \end{pmatrix} + x_4 \begin{pmatrix} 0\\1 \end{pmatrix}$$



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$$B = \{1, 4\}$$
 is a basis
 $x_2 = t \ge 0, x_3 = 0$

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 $B = \{1, 4\} \text{ is a basis}$ $x_2 = t \ge 0, x_3 = 0$ $\underbrace{\begin{pmatrix} x_1 \\ x_4 \end{pmatrix}}_{x_B} = \underbrace{\begin{pmatrix} 2 \\ 5 \end{pmatrix}}_{b} - t \underbrace{\begin{pmatrix} 1 \\ 1 \end{pmatrix}}_{A_2}$





Choose $t \ge 0$ as large as possible.



Basic variables must remain non-negative.



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Basic variables must remain non-negative.

 $x_1 = 2 - t \ge 0 \quad \Longrightarrow \quad t \le \frac{2}{1}$



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Thus, the largest possible $t = \min \left\{ \frac{2}{1}, \frac{5}{1} \right\}$.





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Thus, the largest possible $t = \min\left\{\frac{2}{1}, \frac{5}{1}\right\}$.

The new feasible solution is $x = (0, 2, 0, 3)^{\top}$. It has value 2 > 0.

$$\begin{array}{ll} \max & (0 & 1 & 3 & 0)x \\ \text{s.t.} & \\ & \begin{pmatrix} 1 & 1 & 2 & 0 \\ 0 & 1 & 1 & 1 \end{pmatrix} x = \begin{pmatrix} 2 \\ 5 \end{pmatrix} \\ & x_1, x_2, x_3, x_4 \ge 0 \end{array}$$

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Remark

The new feasible solution $x = (0, 2, 0, 3)^{\top}$ is a basic solution.

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Question

For what basis B is $x = (0, 2, 0, 3)^{\top}$ a basic solution?

$$\max \quad (0 \quad 1 \quad 3 \quad 0)x$$

s.t.
$$\begin{pmatrix} 1 \quad 1 \quad 2 \quad 0\\ 0 \quad 1 \quad 1 \quad 1 \end{pmatrix} x = \begin{pmatrix} 2\\ 5 \end{pmatrix}$$
$$x_1, x_2, x_3, x_4 \ge 0$$

Remark

The new feasible solution $x = (0, 2, 0, 3)^{\top}$ is a basic solution.

Question

For what basis B is $x = (0, 2, 0, 3)^{\top}$ a basic solution?

$$x_2 \neq 0 \implies 2 \in B$$
max
$$(0 \ 1 \ 3 \ 0)x$$

s.t.
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For what basis B is $x = (0, 2, 0, 3)^{\top}$ a basic solution?

$$\begin{array}{ccc} x_2 \neq 0 & \longrightarrow & 2 \in B \\ x_4 \neq 0 & \longrightarrow & 4 \in B \end{array}$$

max
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Question

For what basis B is $x = (0, 2, 0, 3)^{\top}$ a basic solution?

$$\begin{array}{ccc} x_2 \neq 0 & \longrightarrow & 2 \in B \\ x_4 \neq 0 & \longrightarrow & 4 \in B \end{array}$$

As |B| = 2, $B = \{2, 4\}$.

| \max | (<mark>0</mark> | 1 | 3 | 0)x | | | |
|----------------------------|------------------|---|---|-----------------------------|--|--|--|
| s.t. | | | | | | | |
| | (1 | 1 | 2 | $\binom{0}{r-\binom{2}{2}}$ | | | |
| | (0 | 1 | 1 | $1)^{x} = (5)$ | | | |
| $x_1, x_2, x_3, x_4 \ge 0$ | | | | | | | |

$$\begin{array}{ll} \max & (0 & 1 & 3 & 0)x \\ \text{s.t.} & \\ & \begin{pmatrix} 1 & 1 & 2 & 0 \\ 0 & 1 & 1 & 1 \end{pmatrix} x = \begin{pmatrix} 2 \\ 5 \end{pmatrix} \\ & x_1, x_2, x_3, x_4 \ge 0 \end{array}$$

 $\{1,4\}$ is a feasible basis

Canonical form for $\{1,4\}$

$$\max \quad (0 \quad 1 \quad 3 \quad 0)x$$
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NEW

 $\{2,4\}$ is a feasible basis

Canonical form for $\{2,4\}$

$$\max \quad \begin{pmatrix} 0 & 1 & 3 & 0 \end{pmatrix} x$$

s.t.
$$\begin{pmatrix} 1 & 1 & 2 & 0 \\ 0 & 1 & 1 & 1 \end{pmatrix} x = \begin{pmatrix} 2 \\ 5 \end{pmatrix}$$
$$x_1, x_2, x_3, x_4 \ge 0$$

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Canonical form for $\{1,4\}$



| max | (-1) | 0 | 1 | (0)x + 2 | |
|------|--|---------------------------------------|-----------|--|---------------------------------------|
| s.t. | | | | | |
| | $\begin{pmatrix} 1\\ -1 \end{pmatrix}$ | $\begin{array}{c} 1 \\ 0 \end{array}$ | $2 \\ -1$ | $\begin{pmatrix} 0\\1 \end{pmatrix} x =$ | $\begin{pmatrix} 2\\ 3 \end{pmatrix}$ |
| | $x_1, x_2,$ | x_3 , | $x_4 \ge$ | 0 | |

/ 1

NEW

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$$\max \quad \begin{pmatrix} 0 & 1 & 3 & 0 \end{pmatrix} x$$

s.t.
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$$x_1, x_2, x_3, x_4 \ge 0$$

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max $(-1 \ 0 \ 1 \ 0)x + 2$ s.t. $\begin{pmatrix} 1 \ 1 \ 2 \ 0 \\ -1 \ 0 \ -1 \ 1 \end{pmatrix}x = \begin{pmatrix} 2 \\ 3 \end{pmatrix}$ $x_1, x_2, x_3, x_4 \ge 0$

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Remark

We only need to know how to go from the OLD basis to a NEW basis!

$$\max \quad (0 \quad 1 \quad 3 \quad 0)x$$

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Canonical form for $\{1,4\}$



| \max | (-1) | 0 | 1 | (0)x + 2 | |
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• 2 <u>entered</u> the basis.

$$\begin{array}{ll} \max & (0 & 1 & 3 & 0)x \\ \text{s.t.} & \\ & \begin{pmatrix} 1 & 1 & 2 & 0 \\ 0 & 1 & 1 & 1 \end{pmatrix} x = \begin{pmatrix} 2 \\ 5 \end{pmatrix} \\ & x_1, x_2, x_3, x_4 \ge 0 \end{array}$$

 $\{1,4\}$ is a feasible basis

Canonical form for $\{1,4\}$

 $\mathbf{1}$

 $\begin{array}{ll} \max & (-1 & 0 & 1 & 0)x + 2 \\ \text{s.t.} & \\ & \begin{pmatrix} 1 & 1 & 2 & 0 \\ -1 & 0 & -1 & 1 \end{pmatrix} x = \begin{pmatrix} 2 \\ 3 \end{pmatrix} \\ & x_1, x_2, x_3, x_4 \ge 0 \end{array}$

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 $\{2,4\}$ is a feasible basis

Canonical form for $\{2,4\}$

Remark

We only need to know how to go from the OLD basis to a NEW basis!

- 2 <u>entered</u> the basis.
- 1 left the basis.

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Canonical form for $\{1,4\}$

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NEW

 $\{2,4\}$ is a feasible basis

Canonical form for $\{2,4\}$

Remark

We only need to know how to go from the OLD basis to a NEW basis!

- 2 <u>entered</u> the basis. WHY?
- 1 left the basis.

$$\max (0 \ 1 \ 3 \ 0)x$$
s.t.
$$\begin{pmatrix} 1 \ 1 \ 2 \ 0\\ 0 \ 1 \ 1 \ 1 \end{pmatrix}x = \begin{pmatrix} 2\\ 5 \end{pmatrix}$$

$$x_1, x_2, x_3, x_4 \ge 0$$

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Canonical form for $\{1, 4\}$

Pick $2 \notin B$ and set $x_2 = t \ge 0$.

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2 enters the basis

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Canonical form for $\{1, 4\}$

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2 enters the basis

Set
$$\begin{pmatrix} x_1 \\ x_4 \end{pmatrix} = \begin{pmatrix} 2 \\ 5 \end{pmatrix} - t \begin{pmatrix} 1 \\ 1 \end{pmatrix}$$
 and

$$\begin{array}{ll} \max & (0 & 1 & 3 & 0)x \\ \text{s.t.} & \\ & \begin{pmatrix} 1 & 1 & 2 & 0 \\ 0 & 1 & 1 & 1 \end{pmatrix} x = \begin{pmatrix} 2 \\ 5 \end{pmatrix} \\ & x_1, x_2, x_3, x_4 \ge 0 \end{array}$$

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 and $t = \min\left\{\frac{2}{1}, \frac{5}{1}\right\} = 2$.

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Canonical form for $\{1, 4\}$

Pick $2 \notin B$ and set $x_2 = t \ge 0$.

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 and $t = \min\left\{\frac{2}{1}, \frac{5}{1}\right\} = 2$.

 \implies $x_1 = 0$ and 1 leaves the basis

$$\begin{array}{ll} \max & (0 & 1 & 3 & 0)x \\ \text{s.t.} & \\ & \begin{pmatrix} 1 & 1 & 2 & 0 \\ 0 & 1 & 1 & 1 \end{pmatrix} x = \begin{pmatrix} 2 \\ 5 \end{pmatrix} \\ & x_1, x_2, x_3, x_4 \ge 0 \end{array}$$

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Canonical form for $\{1, 4\}$

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Set
$$\begin{pmatrix} x_1\\ x_4 \end{pmatrix} = \begin{pmatrix} 2\\ 5 \end{pmatrix} - t \begin{pmatrix} 1\\ 1 \end{pmatrix}$$
 and $t = \min\left\{\frac{2}{1}, \frac{5}{1}\right\} = 2.$

 \implies $x_1 = 0$ and 1 leaves the basis

The NEW basis is $\{2, 4\}$.



 $B = \{2, 4\}$ is a feasible basis Canonical form for B



 $B = \{2, 4\}$ is a feasible basis Canonical form for B

Pick $k \notin B$ such that $c_k > 0$ and set $x_k = t$:



 $B = \{2, 4\} \text{ is a feasible basis}$ Canonical form for B

Pick $k \notin B$ such that $c_k > 0$ and set $x_k = t$:

 $x_3 = t$



 $B = \{2, 4\} \text{ is a feasible basis}$ Canonical form for B

Pick $k \notin B$ such that $c_k > 0$ and set $x_k = t$:

 $x_3 = t \implies 3 \text{ <u>enters</u>}$ the basis



 $B = \{2, 4\} \text{ is a feasible basis}$ Canonical form for B

Pick $k \notin B$ such that $c_k > 0$ and set $x_k = t$:

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Pick $x_B = b - tA_k$:



 $B = \{2, 4\}$ is a feasible basis Canonical form for B

Pick $k \notin B$ such that $c_k > 0$ and set $x_k = t$:

 $x_3 = t$ \longrightarrow 3 <u>enters</u> the basis

Pick $x_B = b - tA_k$: $\begin{pmatrix} x_2 \\ x_4 \end{pmatrix} = \begin{pmatrix} 2 \\ 3 \end{pmatrix} - t \begin{pmatrix} 2 \\ -1 \end{pmatrix}$



 $B = \{2, 4\}$ is a feasible basis Canonical form for B

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Pick $k \notin B$ such that $c_k > 0$ and set $x_k = t$:

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Pick $x_B = b - tA_k$: $\begin{pmatrix} x_2 \\ x_4 \end{pmatrix} = \begin{pmatrix} 2 \\ 3 \end{pmatrix} - t \begin{pmatrix} 2 \\ -1 \end{pmatrix}$ $t = \min \left\{ \frac{2}{2}, - \right\} = 2$ thus $x_2 = 0$ \longrightarrow 2 <u>leaves</u> the basis



 $B = \{2, 4\} \text{ is a feasible basis}$ Canonical form for B

Pick $k \notin B$ such that $c_k > 0$ and set $x_k = t$:

 $x_3 = t \implies 3 \text{ <u>enters</u>}$ the basis

Pick $x_B = b - tA_k$: $\begin{pmatrix} x_2 \\ x_4 \end{pmatrix} = \begin{pmatrix} 2 \\ 3 \end{pmatrix} - t \begin{pmatrix} 2 \\ -1 \end{pmatrix}$ $t = \min \left\{ \frac{2}{2}, - \right\} = 2$ thus $x_2 = 0 \implies 2$ <u>leaves</u> the basis The NEW basis is $B = \{3, 4\}$.





Pick $k \notin B$ such that $c_k > 0$ and set $x_k = t$: ???



Pick $k \notin B$ such that $c_k > 0$ and set $x_k = t$: ???

Claim $(0,0,1,4)^\top \text{ has value } 3. \text{ It is optimal because } 3 \text{ is an upper bound.}$



Pick $k \notin B$ such that $c_k > 0$ and set $x_k = t$: ???

Claim

 $(0,0,1,4)^{\top}$ has value 3. It is optimal because 3 is an upper bound.

Proof



Pick $k \notin B$ such that $c_k > 0$ and set $x_k = t$: ???

Claim

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Proof

Let x be a feasible solution.



Pick $k \notin B$ such that $c_k > 0$ and set $x_k = t$: ???

Claim

 $(0,0,1,4)^{\top}$ has value 3. It is optimal because 3 is an upper bound.

Proof

Let x be a feasible solution. Then

$$\underbrace{(-1.5, \ 0.5, \ 0, \ 0)}_{\leq \mathbf{0}} \underbrace{x}_{\geq \mathbf{0}} + 3$$



Pick $k \notin B$ such that $c_k > 0$ and set $x_k = t$: ???

Claim

 $(0,0,1,4)^{\top}$ has value 3. It is optimal because 3 is an upper bound.

Proof

Let x be a feasible solution. Then

$$\underbrace{(-1.5, \ 0.5, \ 0, \ 0)}_{\leq \mathbf{0}} \underbrace{x}_{\geq \mathbf{0}} + 3 \leq 3.$$
$\begin{array}{ll} \max & (0 & -4 & 3 & 0 & 0)x \\ \text{s.t.} & \\ & \begin{pmatrix} 1 & -2 & 1 & 0 & 0 \\ 0 & 5 & -3 & 1 & 0 \\ 0 & 4 & -2 & 0 & 1 \end{pmatrix} x = \begin{pmatrix} 1 \\ 1 \\ 2 \end{pmatrix} \\ & x_1, x_2, x_3, x_4 \ge 0 \end{array}$

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 $x_3 = t \implies 3$ enters the basis

 $\begin{array}{ll} \max & (0 & -4 & 3 & 0 & 0)x \\ \text{s.t.} & \\ & \begin{pmatrix} 1 & -2 & 1 & 0 & 0 \\ 0 & 5 & -3 & 1 & 0 \\ 0 & 4 & -2 & 0 & 1 \end{pmatrix} x = \begin{pmatrix} 1 \\ 1 \\ 2 \end{pmatrix} \\ & x_1, x_2, x_3, x_4 \ge 0 \end{array}$

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$$\begin{pmatrix} x_1 \\ x_4 \\ x_5 \end{pmatrix} = \begin{pmatrix} 1 \\ 1 \\ 2 \end{pmatrix} - t \begin{pmatrix} 1 \\ -3 \\ -2 \end{pmatrix}$$

 $\begin{array}{ll} \max & (0 & -4 & 3 & 0 & 0)x \\ \text{s.t.} & \\ & \begin{pmatrix} 1 & -2 & 1 & 0 & 0 \\ 0 & 5 & -3 & 1 & 0 \\ 0 & 4 & -2 & 0 & 1 \end{pmatrix} x = \begin{pmatrix} 1 \\ 1 \\ 2 \end{pmatrix} \\ & x_1, x_2, x_3, x_4 \ge 0 \end{array}$

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Pick $k \notin B$ such that $c_k > 0$ and set $x_k = t$:

 $x_3 = t$ \implies 3 <u>enters</u> the basis

Pick $x_B = b - tA_k$: $\begin{pmatrix} x_1 \\ x_4 \\ x_5 \end{pmatrix} = \begin{pmatrix} 1 \\ 1 \\ 2 \end{pmatrix} - t \begin{pmatrix} 1 \\ -3 \\ -2 \end{pmatrix}$ $t = \min \left\{ \frac{1}{1}, -, - \right\} = 1 \text{ thus } x_1 = 0 \implies 1 \text{ <u>leaves the basis}</u>$

 $\begin{array}{ll} \max & (0 & -4 & 3 & 0 & 0)x \\ \text{s.t.} & \\ & \begin{pmatrix} 1 & -2 & 1 & 0 & 0 \\ 0 & 5 & -3 & 1 & 0 \\ 0 & 4 & -2 & 0 & 1 \end{pmatrix} x = \begin{pmatrix} 1 \\ 1 \\ 2 \end{pmatrix} \\ & x_1, x_2, x_3, x_4 \ge 0 \end{array}$

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$$x_B = b - tA_k$$
:

$$\begin{pmatrix} x_1 \\ x_4 \\ x_5 \end{pmatrix} = \begin{pmatrix} 1 \\ 1 \\ 2 \end{pmatrix} - t \begin{pmatrix} 1 \\ -3 \\ -2 \end{pmatrix}$$

$$t = \min\left\{\frac{1}{1}, -, -\right\} = 1 \text{ thus } x_1 = 0 \implies 1 \text{ leaves the basis}$$
The NEW basis is $B = \{3, 4, 5\}$

Pick $k \notin B$ such that $c_k > 0$ and set $x_k = t$:

$$\begin{array}{rll} \max & (-3 & 2 & 0 & 0 & 0)x + 3\\ \text{s.t.} & \\ & \begin{pmatrix} 1 & -2 & 1 & 0 & 0\\ 3 & -1 & 0 & 1 & 0\\ 2 & 0 & 0 & 0 & 1 \end{pmatrix} x = \begin{pmatrix} 1\\ 4\\ 4 \end{pmatrix} \\ & x_1, x_2, x_3, x_4 \ge 0 \end{array}$$

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$$x_2 = t$$

Pick $k \notin B$ such that $c_k > 0$ and set $x_k = t$:

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max
$$(-3 \ 2 \ 0 \ 0 \ 0)x + 3$$

s.t.
$$\begin{pmatrix} 1 \ -2 \ 1 \ 0 \ 0 \\ 3 \ -1 \ 0 \ 1 \ 0 \\ 2 \ 0 \ 0 \ 0 \ 1 \end{pmatrix} x = \begin{pmatrix} 1 \\ 4 \\ 4 \end{pmatrix}$$
 $x_1, x_2, x_3, x_4 \ge 0$

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max
$$(-3 \ 2 \ 0 \ 0 \ 0)x + 3$$

s.t.
$$\begin{pmatrix} 1 \ -2 \ 1 \ 0 \ 0 \\ 3 \ -1 \ 0 \ 1 \ 0 \\ 2 \ 0 \ 0 \ 0 \ 1 \end{pmatrix} x = \begin{pmatrix} 1 \\ 4 \\ 4 \end{pmatrix}$$
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:

$$\begin{pmatrix} x_3 \\ x_4 \\ x_5 \end{pmatrix} = \begin{pmatrix} 1 \\ 4 \\ 4 \end{pmatrix} - t \begin{pmatrix} -2 \\ -1 \\ 0 \end{pmatrix} \qquad \text{Choose } t = ???$$

Claim

The linear program is unbounded.

$$\begin{aligned} x_1 &= 0\\ \begin{pmatrix} x_3\\ x_4\\ x_5 \end{pmatrix} = \begin{pmatrix} 1\\ 4\\ 4 \end{pmatrix} - t \begin{pmatrix} -2\\ -1\\ 0 \end{pmatrix} \end{aligned}$$

The linear program is unbounded.

The linear program is unbounded.

Proof

$$\begin{bmatrix} \max & z = (-3 & 2 & 0 & 0 & 0)x + 3 \\ \text{s.t.} & & & \\ & \begin{pmatrix} 1 & -2 & 1 & 0 & 0 \\ 3 & -1 & 0 & 1 & 0 \\ 2 & 0 & 0 & 0 & 1 \end{pmatrix} x = \begin{pmatrix} 1 \\ 4 \\ 4 \end{pmatrix} \\ & \begin{pmatrix} x_3 \\ x_4 \\ x_5 \end{pmatrix} = \begin{pmatrix} 1 \\ 4 \\ 4 \end{pmatrix} - t \begin{pmatrix} -2 \\ -1 \\ 0 \end{pmatrix} x_5 = \begin{pmatrix} 1 \\ 4 \\ 4 \end{pmatrix} = \frac{1}{2} \begin{bmatrix} 1 \\ 4 \\ 4 \end{bmatrix} =$$

The linear program is unbounded.

Proof

$$x(t) = \begin{pmatrix} 0\\t\\1+2t\\4+t\\4 \end{pmatrix} =$$

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Proof

$$x(t) = \begin{pmatrix} 0 \\ t \\ 1+2t \\ 4+t \\ 4 \end{pmatrix} = \underbrace{\begin{pmatrix} 0 \\ 0 \\ 1 \\ 4 \\ 4 \end{pmatrix}}_{=\bar{x}} + t \underbrace{\begin{pmatrix} 0 \\ 1 \\ 2 \\ 1 \\ 0 \\ 0 \end{bmatrix}}_{=r}$$

Proof

The linear program is unbounded.

$$x(t) = \begin{pmatrix} 0 \\ t \\ 1+2t \\ 4+t \\ 4 \end{pmatrix} = \underbrace{\begin{pmatrix} 0 \\ 0 \\ 1 \\ 4 \\ 4 \end{pmatrix}}_{=\bar{x}} + t \underbrace{\begin{pmatrix} 0 \\ 1 \\ 2 \\ 1 \\ 0 \\ =r \end{pmatrix}}_{=r}$$

• x(t) is feasible for all $t \ge 0$.

The linear program is unbounded.

Proof $x(t) = \begin{pmatrix} 0 \\ t \\ 1+2 \\ 4+2$

$$\begin{pmatrix} 0\\t\\1+2t\\4+t\\4 \end{pmatrix} = \underbrace{\begin{pmatrix} 0\\0\\1\\4\\4 \end{pmatrix}}_{=\bar{x}} + t \underbrace{\begin{pmatrix} 0\\1\\2\\1\\0 \end{pmatrix}}_{=r}$$

- x(t) is feasible for all $t \ge 0$.
- $z \to \infty$ when $t \to \infty$.

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Proof

$$x(t) = \begin{pmatrix} 0\\t\\1+2t\\4+t\\4 \end{pmatrix} = \underbrace{\begin{pmatrix} 0\\0\\1\\4\\4 \end{pmatrix}}_{=\bar{x}} + t \underbrace{\begin{pmatrix} 0\\1\\2\\1\\0 \end{pmatrix}}_{=r}$$

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- $z \to \infty$ when $t \to \infty$.

 (\bar{x}, r) : certificate of unboundedness.)

$$\begin{array}{ll} \max & c^{\top}x\\ \text{s.t.} & \\ & Ax = b\\ & x \geq \mathbf{0} \end{array}$$

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INPUT:

$$\begin{array}{ll} \max & c^{\top}x \\ \text{s.t.} & \\ & Ax = b \\ & x \ge \mathbf{0} \end{array}$$

<u>INPUT:</u> a feasible basis B.

$$\begin{array}{ll} \max & c^{\top}x\\ \text{s.t.} & \\ & Ax = b\\ & x \geq \mathbf{0} \end{array}$$

<u>INPUT:</u> a feasible basis B.

<u>Output:</u>

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- <u>INPUT:</u> a feasible basis B.
- $\underline{OUTPUT:} \quad \text{an optimal solution OR}$

$$\begin{array}{ll} \max & c^{\top}x\\ \mathsf{s.t.} & \\ & Ax = b\\ & x \geq \mathbf{0} \end{array}$$

<u>INPUT:</u> a feasible basis B.

<u>OUTPUT:</u> an optimal solution OR it detects that the LP is unbounded.

$$\begin{array}{ll} \max & c^{\top}x\\ \text{s.t.} & \\ & Ax = b\\ & x \geq \mathbf{0} \end{array}$$

<u>INPUT:</u> a feasible basis B.

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Step 1. Rewrite in canonical form for the basis *B*.

$$\begin{array}{ll} \max & c^{\top}x\\ \text{s.t.} & \\ & Ax = b\\ & x \geq \mathbf{0} \end{array}$$

<u>INPUT:</u> a feasible basis B.

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Step 1. Rewrite in canonical form for the basis *B*.

Step 2. Find a better basis *B* or get required outcome.
The Simplex Algorithm

$$\begin{array}{ll} \max & c^{\top}x\\ \text{s.t.} & \\ & Ax = b\\ & x \geq \mathbf{0} \end{array}$$

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Step 1. Rewrite in canonical form for the basis *B*.



Step 2. Find a better basis B or get required outcome.

$$\begin{array}{ll} \max & z = c_N^\top x_N + \bar{z} \\ \text{s.t.} & \\ & x_B + A_N x_N = b \\ & x \geq \mathbf{0} \end{array}$$

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B is a feasible basis, $N = \{j \notin B\}$

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 $\bar{\boldsymbol{x}}$ is a basic solution

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If $c_N \leq \mathbf{0}$, then STOP. The basic solution \bar{x} is optimal.

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If $A_k \leq \mathbf{0}$, then STOP. The LP is unbounded.

$$\begin{array}{ll} \max & z = c_N^\top x_N + \bar{z} \\ \text{s.t.} & \\ & x_B + A_N x_N = b \\ & x \geq \mathbf{0} \end{array}$$

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Choose
$$t = \min\left\{\frac{b_i}{A_{ik}} : \text{for all } i \text{ such that } A_{ik} > 0\right\}$$
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Choose
$$t = \min\left\{\frac{b_i}{A_{ik}} : \text{for all } i \text{ such that } A_{ik} > 0\right\}.$$

Let x_r be a basic variable forced to 0.

$$\begin{array}{ll} \max & z = c_N^\top x_N + \bar{z} \\ \text{s.t.} & \\ & x_B + A_N x_N = b \\ & x \geq \mathbf{0} \end{array}$$

B is a feasible basis, $N=\{j\notin B\}$

Canonical form for B

 $\bar{\boldsymbol{x}}$ is a basic solution

If $c_N \leq \mathbf{0}$, then STOP. The basic solution \bar{x} is optimal.

Pick $k \notin B$ such that $c_k > 0$ and set $x_k = t$. Pick $x_B = b - tA_k$.

If $A_k \leq 0$, then STOP. The LP is unbounded.

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$$t = \min \left\{ \frac{b_i}{A_{ik}} : \text{for all } i \text{ such that } A_{ik} > 0 \right\}.$$

Let x_r be a basic variable forced to 0.

The new basis is obtained by having k enter and r leave.

$$\begin{array}{ll} \max & z = c_N^\top x_N + \bar{z} \\ \text{s.t.} & \\ & x_B + A_N x_N = b \\ & x \geq \mathbf{0} \end{array}$$

Canonical form for B

 $\bar{\boldsymbol{x}}$ basic solution

If $c_N \leq \mathbf{0}$, then STOP. The basic solution \bar{x} is optimal.

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If $c_N \leq \mathbf{0}$, then STOP. The basic solution \bar{x} is optimal.

Proof

$$\begin{array}{ll} \max & z = c_N^\top x_N + \bar{z} \\ \text{s.t.} & \\ & x_B + A_N x_N = b \\ & x \geq \mathbf{0} \end{array}$$

Canonical form for B

 $\bar{\boldsymbol{x}}$ basic solution

If $c_N \leq \mathbf{0}$, then STOP. The basic solution \bar{x} is optimal.

Proof

 $\bar{x}_B = b, \ \bar{x}_N = \mathbf{0}.$

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Proof

 $\bar{x}_B = b, \ \bar{x}_N = \mathbf{0}.$

 \bar{x} has value $z = c_N^\top \bar{x}_N + \bar{z} = \bar{z}$.

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Let x be a feasible solution.

$$z =$$

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$$\bar{x}$$
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Let x be a feasible solution.

$$z = \underbrace{c_N^\top}_{\leq \mathbf{0}} \underbrace{x_N}_{\geq \mathbf{0}} + \bar{z}$$

$$\begin{array}{ll} \max & z = c_N^\top x_N + \bar{z} \\ \text{s.t.} & \\ & x_B + A_N x_N = b \\ & x \geq \mathbf{0} \end{array}$$

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$$\begin{array}{ll} \max & z = c_N^\top x_N + \bar{z} \\ \text{s.t.} & \\ & x_B + A_N x_N = b \\ & x \geq \mathbf{0} \end{array}$$

Canonical form for B

 $\bar{\boldsymbol{x}}$ basic solution

If $A_k \leq \mathbf{0}$, then STOP. The LP is unbounded.

$$\begin{array}{ll} \max & z = c_N^\top x_N + \bar{z} \\ \text{s.t.} & \\ & x_B + A_N x_N = b \\ & x \geq \mathbf{0} \end{array}$$

Canonical form for B

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If $A_k \leq \mathbf{0}$, then STOP. The LP is unbounded.

Proof

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Let us see an example...

$$\begin{array}{ll} \max & (0 & 0 & 2 & 3)x \\ \text{s.t.} & \\ & \begin{pmatrix} 1 & 0 & 2 & -1 \\ 0 & 1 & 4 & 6 \end{pmatrix} x = \begin{pmatrix} 6 \\ 12 \end{pmatrix} \\ & x_1, x_2, x_3, x_4 \ge 0 \end{array}$$

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Pick $k \notin B$ such that $c_k > 0$ and set $x_k = t$:

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The NEW basis is $B = \{3, 4\}$.

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To do: Find a procedure to find a feasible basis.