

Module 2: Linear Programs (Certificates)

Recap and a Question

Fundamental Theorem of Linear Programming

For any linear program, **exactly one** of the following holds:

- It is infeasible.
- It has an optimal solution.
- It is unbounded.

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Consider a linear program.

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- If it is infeasible, how can we **prove** it?

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- If we have an optimal solution, how can we **prove** it is optimal?

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- If we have an optimal solution, how can we **prove** it is optimal?
- If it is unbounded, how can we **prove** it?

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Consider a linear program.

- If it is infeasible, how can we **prove** it?
- If we have an optimal solution, how can we **prove** it is optimal?
- If it is unbounded, how can we **prove** it?

This can be always be done!

Proving Infeasibility

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The following linear program is infeasible:

$$\max \quad (3, 4, -1, 2)^T x$$

s.t.

$$\begin{pmatrix} 3 & -2 & -6 & 7 \\ 2 & -1 & -2 & 4 \end{pmatrix} x = \begin{pmatrix} 6 \\ 2 \end{pmatrix}$$

$$x \geq 0$$

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Question

How can we prove this problem is, in fact, infeasible?

We **cannot** try all possible assignments of values to x_1, x_2, x_3 , and x_4 .

Claim

There is no solution to (1), (2) and $x \geq 0$ where

$$\begin{pmatrix} 3 & -2 & -6 & 7 \\ 2 & -1 & -2 & 4 \end{pmatrix} x = \begin{pmatrix} 6 \\ 2 \end{pmatrix} \quad \begin{matrix} (1) \\ (2) \end{matrix}$$

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Proof

Construct a new equation:

$$\begin{array}{r} -1 \times (1) : (-3 \quad 2 \quad 6 \quad -7)x = -6 \\ + \quad 2 \times (2) : (4 \quad -2 \quad -4 \quad 8)x = 4 \\ \hline (1 \quad 0 \quad 2 \quad 1)x = -2 \end{array} \quad (*)$$

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Contradiction!

Repeat using **matrix formulations**.

Proof

Suppose for a contradiction there is a solution \bar{x} to $x \geq 0$ and

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$$(-1 \ 2) \begin{pmatrix} 3 & -2 & -6 & 7 \\ 2 & -1 & -2 & 4 \end{pmatrix} x = (-1 \ 2) \begin{pmatrix} 6 \\ 2 \end{pmatrix}$$

$$(1 \ 0 \ 2 \ 1)x = -2 \quad (\star)$$

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$$\underbrace{\begin{pmatrix} 3 & -2 & -6 & 7 \\ 2 & -1 & -2 & 4 \end{pmatrix}}_A x = \underbrace{\begin{pmatrix} 6 \\ 2 \end{pmatrix}}_b$$

Construct a new equation:

$$\underbrace{\begin{pmatrix} -1 & 2 \end{pmatrix}}_{y^T} \begin{pmatrix} 3 & -2 & -6 & 7 \\ 2 & -1 & -2 & 4 \end{pmatrix} x = \underbrace{\begin{pmatrix} -1 & 2 \end{pmatrix}}_{y^T} \begin{pmatrix} 6 \\ 2 \end{pmatrix}$$

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$$\underbrace{\begin{pmatrix} 3 & -2 & -6 & 7 \\ 2 & -1 & -2 & 4 \end{pmatrix}}_A x = \underbrace{\begin{pmatrix} 6 \\ 2 \end{pmatrix}}_b \quad Ax = b$$

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Contradiction.

This suggests **the following result**...

Proposition

There is no solution to $Ax = b$, $x \geq 0$, if there exists y where

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Give a proof of this proposition.

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If no solution to $Ax = b$, $x \geq 0$ can we always prove it in that way?

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YES!!!!

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YES!!!!!!

Farkas' Lemma

If there is no solution to $Ax = b$, $x \geq 0$, then there exists y where

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Proving Optimality

$$\max \quad z(x) := (-1 \ -4 \ 0 \ 0)x + 4$$

s.t.

$$\begin{pmatrix} 1 & 3 & 1 & 0 \\ -2 & 6 & 0 & 1 \end{pmatrix} x = \begin{pmatrix} 4 \\ 5 \end{pmatrix}$$

$$x \geq 0$$

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Optimal solution:

$$\bar{x}_1 = 0$$

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We **cannot** try all possible feasible solutions.

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Claim

- \bar{x} is feasible solution of value 4.

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Claim

- \bar{x} is feasible solution of value 4. (easy)

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- 4 is an **upper bound**.

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Let x' be an **arbitrary** feasible solution.

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Let x' be an **arbitrary** feasible solution. Then

$$z(x') = (-1 \ -4 \ 0 \ 0)x' + 4$$

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$$z(x') = \underbrace{(-1 \ -4 \ 0 \ 0)x'}_{\leq 0} + 4 \leq 4.$$

Proving Unboundedness

$$\max \quad z := (-1 \ 0 \ 0 \ 1)x$$

s.t.

$$\begin{pmatrix} -1 & -1 & 1 & 0 \\ -2 & 1 & 0 & 1 \end{pmatrix} x = \begin{pmatrix} 2 \\ 1 \end{pmatrix}$$

$$x \geq 0$$

Proving Unboundedness

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Problem
is unbounded

Question

How can we prove that this problem is unbounded?

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Idea

Construct a family of **feasible solutions** $x(t)$ for all $t \geq 0$ and show that as t goes to infinity, the value of the objective function goes to infinity.

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$$x \geq 0$$

$$x(t) := \begin{pmatrix} 0 \\ 0 \\ 2 \\ 1 \end{pmatrix} + t \begin{pmatrix} 1 \\ 0 \\ 1 \\ 2 \end{pmatrix}$$

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Generalize and prove the following proposition.

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Proposition

The linear program,

$$\max\{c^T x : Ax = b, x \geq 0\}$$

is unbounded if we can find \bar{x} and r such that

$$\bar{x} \geq 0, \quad r \geq 0, \quad A\bar{x} = b, \quad Ar = 0 \quad \text{and} \quad c^T r > 0.$$

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Remark

We have not yet shown you how to **find** such proofs.